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1. Introduction

I am a Computer Science student, and I know from my own experience that it may be hard to grasp at some concepts and ideas. I have noticed that people find some topics in Computing very difficult to understand. When it came to choosing my A-level project I remembered that. Computing is quite a different subject from others. It describes a lot of dynamic processes unlike A-level Maths, and involves a lot of completely new concepts unlike, for example, English or A-level Physics. Practical demonstrations are a well-known teaching technique, and have proven to be very useful. But it is a lot easier to demonstrate, say, attraction between charged objects than a memory read operation. The reason seems to be that Physics has been taught for a lot longer than Computing, and so while most schools have enough equipment in their physics labs, hardly any schools have appropriate software equipment in their computing labs. They have word-processing and program development software, but apart from that only an occasional program to demonstrate a concept or a technique. So I though that I could contribute to that area and develop a system to assist students in learning Computing, and hopefully provide some useful hands-on experience in the field.

The most obious users for my system is one of my Computing teachers and his students. I assume that as an experienced teacher he will know exactly what he wants such a system to do. Although students will also use the system, they are hardly likely to have any idea as to what they want the system to do, apart from it being simple and easy to understand. So I will concentrate on my teacher as the main user.

2. Investigation

As the main user of my system will be my Computer Science teacher, I have had an interview with him. The goal of the interview was to find out the requirements for the new system. Here is an approximate log of the interview:

- I am interested in creating a computer system which would aid students in learning the principles of computers. I am sure there are a lot of things that could be done. Can you think of any programs you might be interested in?
- There are quite a few simulations that could be useful. Starting from showing what arrays are and how data can be sorted to simulating a whole virtual computer in work.
- Which ones do you think would be most useful to you?
- The more complicated the topic is, the more we need a good illustration of how things happen. Probably internal structure of a computer is one of the more complicated topics, at the same time being a vital part of the syllabus.
- Do you mean you need a program that would show the role of different parts in a computer and how they communicate with each other to do useful work?
- That's right.
- OK. Let me think about it. How interactive do you think the simulation should be? It could be just a set of video clips, or it could be based on some sort of initial data.
- In theory the best way I can think of doing this is to create an assembly language interpreter which would show every step required to get a program executed. In that way a student would be able, for instance, to investigate what happens to the buses, or 'feel' the von Neumann machine principle, etc.
- The amount of work required to create such a system is quite high. Do you think it's worth it?
- An assembly language interpreter by itself might not be worth it because although it is a profound topic in computer science, in the exam you usually get no more than 5% of the marks on assembly language. But if there were a system which would help students really understand what computers are made of and why – that would be very useful in my opinion.
- Another issue is how detailed you want the simulation to be. A very detailed system might resemble a real PC well but be too complicated to be of any use in teaching. On the other hand, if you have a very simple system you might not be able to show as much as you want.
- I agree with you. I would like to be able to teach both GCSE and A-level students, so it would be great if there was a way of switching complexity levels. These levels should contain as much information as the syllabuses do.
- *How do you teach assembly language at the moment?*
- I have to show everything on paper or on whiteboard. It is quite hard to show how a program works if you can't run it.
- What areas of assembly language cause most problems to the students?

- Different addressing modes definitely cause most problems, especially indirect addressing. It would be nice if some sort of animations were available. Other aspects, such as shifts, take some time for the students to grasp at.
- Let's think about the diagrams. Which ones do you want for GCSE and A-level?
- Well, probably a general diagram showing all computer components, and a diagram for the internal structure of the CPU for A-level. I understand we would be able to see everything animated on the diagram, e.g. how an instruction is fetched, or how an interrupt is processed?
- Yes, that is the idea. By the way, you would probably need more than two detail levels because, for example, you don't want to see anything but CPU, RAM and buses when you tell people about buses.
- That's true.
- Now let's concentrate on the assembly language. Examiners use different names for instructions and registers than those that are more common in Windows-targeted programming languages. I believe you would prefer examiners' names?
- Yes. And I think that anybody who understands assembly language would learn new names quickly if they need to.
- The instruction set would probably include arithmetic, logic and branching instructions. Anything else?
- Yes. I would like to have a kind of screen which would be able to output either text or graphics. Also primitive interrupt simulation, e.g. when a key is pressed, would be a good idea.
- How many registers would you like to have?
- Accumulator, plus general purpose registers. Four or five general purpose registers would do. I would prefer arithmetic commands to always have one of the operands being the accumulator. Also I would like to see special purpose registers such as Program Counter, Stack Pointer, Flags Register.
- Do you want to have different register sizes, e.g. ah, al being 8 bits each and eax being 32 bits, while still being the same register?
- I think that would confuse students without teaching them anything. It is not such an important part of assembly language to go into trouble of allowing that, and it is very CPU-specific.
- So what register size, or bus width, do you think you need?
- In the majority of things I will use your program for I will not need more than 8 bits address and 8 bits data registers. But if you find it easy you might allow for 16 bits address and data buses. What I definitely want is each memory location to be 8 bits in size.
- Do you want to have any segment registers? What addressing modes do you want to see in the program?
- Immediate, direct and indirect addressing. You might allow for indexed addressing, but it is of no major importance. In addition I would like to have relative jumps and calls. Segment registers... Well, as long as they would not make everything very confusing they would be OK. Probably I would prefer those as modifiers, not true

registers, e.g. when you access vs:[010h] then you always access location 10h in video memory segment. But I want to have a simple mode where everything would be in one big 'segment', including code, stack, data and video memory, to show the von Neumann principle.

- How are we going to store a program? We can either store and interpret it in text form, which would be the easiest way of doing it, where you have code separate from everything else. Alternatively we could write a program in a special window and then compile it into machine code. The latter is definitely more complicated to implement but otherwise I see no big difference.
- I think there is a big difference. The latter would give students an even better idea of how a byte can be interpreted in a lot of different ways, even as an instruction. Also that way we will have an assembler – showing how it works would be wonderful.
- OK. Is there anything else you have on your mind?
- Well, you must understand that good user interface is crucial. It should not look complicated because that would repel students. It is also very important that things are intuitive a student who can't do something would probably not persist and just go away.
- *I* will have that in mind.

3. Requirements

3.1. General requirements

This section describes general functionality which the system is expected to have.

- File operations: Load/Save project
- Writing a program: Program editor (with syntax highlighting if possible)
- Running a program: Run/Stop, Step
- Debugging: Breakpoints, watch variable/register values

Teaching assembly language should be one of requirements for the project, but not the only one. It should be possible to use this system to teach internal workings of a computer, and the CPU in particular, by showing how different components communicate with each other and letting the user see exactly what is going on.

3.2. Simulation requirements

Code

The user will develop the program in assembly language and compile it in machine codes. The virtual CPU will then be able to execute the program.

Memory

The most appropriate memory model will be flat – everything, including code, stack, data and video memory, in one long array at different offsets. This model is the best one to show von Neumann machine, and it is the simplest one to write code for. Memory should be at least 4 kilobytes long.

Storing 16-bit data

All 16-bit data will be stored in RAM in big-endian format, that is the most significant byte will be stored at the lower memory address.

Stack

Stack should start somewhere around the middle of RAM. It should grow towards the end of RAM in order to prevent it overwriting user code.

3.3. Detail Levels

As the system will be aimed at different user levels, it should be possible to change program accordingly. The system should allow the user to choose between Basic mode and A-level mode. For example, in the Basic mode it should be possible for the user to write and debug a program without ever knowing that it is compiled at all.

3.4. Instruction set

- Data movement move data between registers and memory, stack operations
- Arithmetic addition, subtraction, multiplication etc.

- Logic standard operations such as AND, OR, shifts, etc.
- Flags operations concerning the FLAGS register.
- Branching conditional/unconditional jumps, subroutine calls, halt.
- Input/output IO port operations, interrupt operations.

3.5. Registers

- General purpose used as temporary storage for data
- Special purpose registers with special meaning, e.g. stack pointer
- Internal registers cannot be used in programs, but can be viewed

3.6. Addressing modes

- Immediate target is a number
- Register target is a register
- Memory target is a memory cell

3.7. Extra functionality

- Animation of different addressing modes
- Animation on number representation
- Animation of logical and arithmetic shifts

3.8. Security

This system will not contain any sensitive data, and therefore no security is required.

4. Constraints

4.1. Hardware & Software

The college already has a network of computers set up, so it would be most costeffective if this system was able to utilise software and hardware that is already in place. Below is a summary of a typical machine set up:

Hardware:

- IBM-compatible machine
- AMD/PIII 400MHz
- 64 Mb RAM
- 500Mb+ free HDD space
- Video card, 1024x768+ resolution

The hardware requirements for this system will be up to 10Mb of disk space and a minimum 1024x768 resolution video card.

Software:

- Windows 2000 / Windows 98
- Visual Basic 6.0

The only software requirement for this system will be a Win32 operating system.

4.2. CS proficiency

Targeted users of this system will be students, and although they will be Computer Science students they will still need an easily understandable program. Therefore, using the program should be as simple as possible, that is, interface should be as intuitive as possible.

4.3. Feasibility

Taking all requirements and constraints into account, it seems that this project is feasible. It is technically feasible (technology exists to implement the solution), economically feasible (enough funds is available to implement it), legally feasible (a licensed copy of VB will be used, and all extra controls that may be used will be either freeware or licensed), operationally feasible (it will be relatively easy to incorporate new teaching methods using this system, and the users should only be happy to transfer to it).

The project should be completed in about 5 months. Because VB is new to me, I cannot estimate very well the amount of time necessary to implement the solution, so I will concentrate on core features at first and then implement optional functionality if any time is available. That way the project will be schedule feasible.

5. Objectives

The main objective of this project is to provide an interactive teaching tool which can be used in Computing lessons to demonstrate computer science concepts and give students some hands-on experience. The system will have to meet three distinct needs. One is to demonstrate operation of internal components of a computer, especially the CPU. Another objective is to provide a fully functional fool-proof and easy-to-use assembly language simulator. The third requirement is to provide a set of interactive simulations (e.g. sorting algorithms). This requirement is optional and depends on how much time will be available.

It is difficult to assess the system from its teaching potential point of view because it is hard to judge how much this program contributed towards a particular student's achievement. Therefore, project's success should be analysed based on its usability. This project should be considered successful if at least 8 out of 10 students will be able to write their first ever assembly language program on their first ever encounter with the system in a 90 minute session with a teacher available to give advice. The system should also withstand at least half an hour of intentional attempts to make it crash in order to make sure some students will not be able to crash it to avoid work.

6. Solution system

This system should be written in one of the high-level programming languages as its main characteristics are:

- Is not very speed-critical
- Intuitive user interface is crucial
- Design time is limited
- Funds are very limited
- Relatively advanced programming techniques are required

Several programming languages are available. Below is a summary of these languages with all their advantages and disadvantages:

Visual Basic 6.0

Advantages:

- Low cost (already available to programmer)
- Rapid application development
- Easy creation of advanced user interface features

Disadvantages:

- Low programming flexibility
- Low efficiency of compiled code
- Programmer has very little experience with it

Delphi 5.0

Advantages:

- Rapid application development
- Easy creation of advanced user interface features
- Programmer has a lot of experience with it
- High programming flexibility
- Good efficiency of compiled code

Disadvantages:

- High cost (license required)
- Executables tend to be big

Microsoft Visual C++

Advantages:

- High programming flexibility
- High efficiency of compiled code
- Comparatively small executables
- Programmer has some experience with it

Disadvantages:

- High cost (license required)
- Takes a lot of time to create good GUI

Borland C++ 5.02

Advantages:

- Very high programming flexibility
- High efficiency of compiled code
- Comparatively small executables

Disadvantages:

- High cost (license required)
- A lot of time required to develop applications
- Difficult to create good GUI
- Programmer has little experience with it

Conclusion

Below is a summary of language correspondence to key features:

Feature (Importance)	Best	Worst
– Is not very speed-critical (3)	VC, Borland	C, Delphi, VB
– Intuitive user interface is crucial (8)	VB, Delphi, '	VC, Borland C
– Design time is limited (10)	VB, Delphi, '	VC, Borland C
– Advanced programming is required (6)	VC, Borland	C, Delphi, VB
– Funds are very limited (7)	VB, VC, Del	phi, Borland C
– Programmer experience (5)	Delphi, VC, '	VB, Borland C

If every language was assigned points from 3 to 0 depending on how well they match a requirement, and multiplied by the weight (importance) of the requirement, here are the totals for the languages:

VB:	0x3 + 3x8 + 3x10 + 0x6 + 3x7 + 1x5	= 80
VC:	3x3 + 1x8 + 1x10 + 3x6 + 2x7 + 2x5	= 69
Delphi:	1x3 + 2x8 + 2x10 + 1x6 + 1x7 + 3x5	= 67
Borland C:	2x3 + 0x8 + 0x10 + 2x6 + 0x7 + 0x5	= 18

Therefore, Visual Basic is more suitable for developing this project than any other language considered.

Design

The Imaginary Computer

7. Introduction

This section describes the way the system should work. Every system designed for the end-user is a lot more complicated on the inside than it seems on the outside. The same will be true of CLab. Design section will describe the internal workings of a complex system, so the description will be technical and complicated, with a lot of subtle details.

I do not see how to avoid this complexity and detail, but moreover, I do not see the need to do so. The users will only be aware of a fraction of what will be described in this section.

Note that the system will have several complexity levels, and while at full level the users will be able to interact with a good deal of all this, in the basic mode they will only see a tiny fraction.

8. Central Processing Unit

The CPU in this system is *not* going to process commands. It will only be there to show the users how a real CPU executes a program. Therefore, it doesn't have to be fully functional from electronics point of view, but it should show different structural elements of the CPU and how they interact.

8.1. Architecture

The architecture of this CPU is very loosely based on that of the Z80 processor. The whole CPU has been designed from scratch with its teaching purpose always being the main guideline.

The major difference between this CPU and real modern CPU architectures will be that no steps are taken to optimize and speed up program execution in order to keep the structure as simple as possible.



Fig.6.1. CPU diagram

In this diagram, the Control Unit is the main section that does all the "clever" work of the CPU. Its internal structure is hidden because it is too advanced for the level this system is aimed for. Control Unit will run the fetch-decode-execute cycle, coordinating all components around it.

All components of the CPU are connected to the **Internal Data Bus** (IDB). Control Unit uses it to transfer data between itself and components such as registers or the ALU. It also uses it to interact with the world outside the CPU via MDR and Data Bus Buffer, including fetching instructions and port I/O operations.

Internal Address Bus (IAB) is connected to the Control Unit through the **Addressing** component. Addressing component is an adder which, in response to Control Unit's directions, adds together one or more pieces of data that are connected to it in order to obtain a full absolute address for operations with memory. It outputs the resulting address directly to the Internal Address Bus and indirectly (at Control Unit's discretion) to the Address Bus outside the CPU.

Internal Control Bus is not shown in the diagram as it would overload it. There is a network of signals going to each component in the CPU which tell components when and what to do. For example, there are such signals as Select, Read and Write going to the Registers component. They allow Control Unit to choose on which register to operate, and then whether that register should read the data from the IDB or write what it contains to the IDB. Another example would be a signal going to the Addressing component telling it that it should add, say, B register and value on the IDB together.

ALU stands for the Arithmetic Logic Unit, and, as the name suggests, it is responsible for all arithmetic and bitwise calculations. It can do following arithmetic operations: add, subtract, multiply, divide, take remainder from division, increment/decrement, change sign, arithmetic shift. Among its bitwise operations are: AND, OR, XOR, NOT, and logical shifts. If it takes two operands, one always has to be the Accumulator register.

The $\pm 1/2$ component is a separate adder which can add or subtract 1 or 2 from IP and SP registers. This is a very frequent operation, so it would not be wise to do it through the ALU.

The **MDR** and **MAR** registers (Memory Data Register and Memory Address Register) hold the data that has arrived from the respective internal or external bus, and can also output that data to either of the buses. These registers are used in operations with external buses. Some operations also use them as temporary storage.

The **LEA** section (Load Effective Address) is a link between the IAB and IDB. Sometimes, when an effective address has been calculated by the Addressing section, the address needs to be used rather than transferred to the external Address Bus. The LEA section is responsible for transferring the address from the IAB to IDB. As soon as the address is on IDB, it can be saved in any register or used otherwise.

8.1.1. Memory models

To simplify the project, only one memory model will be supported – flat memory model. Segmented memory model, although mentioned in A-level course, is not really studied at all, so it would be better to reduce program functionality with no negative effect on the users, while greatly increasing program simplicity, which is very important for the users. In flat model RAM is seen as a single long block of data. In this project RAM will be 64 Kb long. Anything can be stored anywhere. Though, some memory areas will have a special meaning. Program execution will begin at address 0, so the beginning of RAM will act like a code segment. Stack pointer will

be initialised to address 6000h and will grow upwards, so area of RAM from 6000h onwards will act like a stack segment. Video memory will be initialised to addresses E000h-EFFFh. And finally, the last 256 bytes of RAM will be used for interrupt vectors table. All this will be discussed in detail further in this document.

8.1.2. Addressing modes

The CPU will support three basic addressing modes, one of which will be split into four sub-modes. The modes are listed below:

- Immediate. The operand is a numerical constant, 8 or 16 bits.
- **Register**. The operand is a register.
- Memory. The operand is a memory cell. Is subdivided into following:
 - Direct (Immediate)
 - Indirect (via Register)
 - Indirect (via Immediate)
 - \circ Indexed

Direct (Immediate) mode points at a memory cell whose address is specified as a 16-bit immediate constant. The absolute address calculated is the value of the immediate constant.

Indirect (via Register) mode points at a memory cell whose address is held in a register. The absolute address calculated is the value held in the register.

Indirect (via Immediate) mode points at a memory cell whose address is held in another memory cell (intermediary cell). The address of the intermediary cell is specified as a 16-bit immediate constant. The CPU loads contents of the intermediary cell and uses it as the final absolute address.

Indexed mode allows the program to specify an expression to calculate the required address. The Addressing section of the CPU can add together B register (offset), a 16bit immediate constant (another offset) and a general purpose register multiplied by 1, 2 or 4 in order to calculate the effective address.

8.1.3. Registers

As shown in <u>fig.1</u>, there will be several registers which the CPU will utilise to execute a program. They are divided into sections by type, and their meaning is described below.

General purpose

These registers are used as a temporary storage for data. They are all 16 bits wide.

Name	Description
А	Accumulator. This register is involved in operations with the ALU (it has to be one of the
	operands in two-operand operations). Some operations are shorter when they use the
	accumulator.

В	General purpose register. Also used as base register in memory addressing.
C	General purpose register.
D	General purpose register.
Е	General purpose register.

Special purpose

These registers have a special meaning, e.g. to show where the next program instruction is. They can't be used as sources or destinations in most operations. Though they will be used by some operations indirectly. They are all 16 bits wide.

Name	Description
PC	Program counter. The next instruction to be executed starts at address stored in PC. Can be
	modified indirectly by jump instructions.
SP	Stack Pointer. The next empty cell in stack is at address stored in SP. Used indirectly by stack
	operations.
FLAGS	Flags register. Contains information about current CPU state or the result of some operations.
	Flags map can be found below.

Flags

In the FLAGS register, the low-order byte will contain the basic flags, and the highorder byte will contain auxiliary flags which will help students learn but will not be used by the system otherwise.

Low-order byte

High-order byte

1011	0100	10,00										-			0,00
0	1	2	3	4	5	6	7	8	9	Α	В	С	D	Е	F
Z	S	0	C	i				n	р						

z – zero flag. Is set every time the result of a calculation is zero.

- s sign flag. Is always equal to the high-order bit of the result.
- o overflow flag. Set when operation causes a carry into OR out of high-order bit.
- c carry flag. Set when operations cause a carry out of high-order bit.
- i interrupts flag. When this flag is cleared, processor ignores all interrupts
- n negative flag. Is set when result is negative. n = s and not z
- p positive flag. Is set when result is positive. p = not(s or z) = not s and not z

Internal registers

These registers cannot be used in programs, but can be viewed. They will be contained inside the CPU, and an assembly language programmer will not need to know that they exist at all. But they are crucial to understand how the CPU works.

All registers under the thick line are contained within the Control Unit, all others are outside the Control Unit but inside the CPU.

Name	Description
MAR	Memory Address Register. Acts as a link between the internal CPU address bus and the external address bus.
MDR	Memory Data Register. Acts as a link between the internal CPU data bus and the external data bus.
CIR	Current Instruction Buffer. Accumulates instruction to be decoded and executed over several fetch cycles (for multi-byte instructions).

IS	Interrupt Status. Contains a bit for every hardware interrupt, indicating whether that interrupt
	is pending or not.

8.2. Input/output

8.2.1. Interrupts

The CPU will support software and hardware interrupts. Each external device will be connected to an Interrupt Request line (IRQ) and to an Interrupt Acknowledge line (INTA), which are separate for every device. When a device wants to send an interrupt request, it will send a signal down its IRQ. If the CPU decides to process it, it will send an INTA signal back. The project will allow for up to 16 external devices, so there will be 16 IRQ lines, for interrupt request numbers 0 to 15.

Interrupt requests will have different priorities. IRQ0 will have the highest priority, and IRQ15 – the lowest priority.

Interrupt request number (IRQ) and actual interrupt number that the IRQ generates (INT) are two different concepts, and do not always match in real computers. For example, in IBM-compatible computers IRQ0 (timer) generates INT8, and IRQ1 (keyboard data ready) generates INT9. But because at A-level these concepts are not differentiated, they will be hidden from the user in this project, and IRQ0 to IRQ15 will always invoke interrupts 0 to 15 respectively.

There will be a special register in the CPU to store pending interrupts. The register will be called *IS* for Interrupt Status. It will have one bit for every one of the 16 hardware interrupts. If a bit is set, then a respective interrupt has been requested and accepted by the CPU.

When the CPU receives an interrupt request, it will first check whether interrupts are allowed. Programs will be able to allow or disallow hardware interrupts with special instructions. If interrupts are allowed and the corresponding bit in IS is cleared then the CPU will accept the interrupt request by sending INTA signal to the requesting device and set the corresponding bit in IS.

Having completed an instruction, the CPU will check whether any of the IS bits is set to zero. Interrupt 0 will have the highest priority, and interrupt 15 - the lowest. Therefore, the CPU will start checking bits from 0 to 15. If a bit is set, the CPU will initiate the interrupt.

To initiate a hardware interrupt, the CPU will first clear that interrupt's bit in IS. It will then push PC and FLAGS on stack, clear the Interrupts Flag in FLAGS register, and jump to a respective interrupt handler. Interrupt handlers will end with a special instruction, which will pop FLAGS and PC from stack.

There will be an Interrupt Vector Table of size 256 bytes in RAM which will store effective addresses of interrupt service procedures (ISPs) for respective interrupts. This table will start at address 0FF00h (for interrupt 0) and end at 0FFFEh (for interrupt 127) and contain 16-bit absolute addresses of respective ISPs.

Software interrupts, unlike hardware interrupts, will be called by programs themselves, and cannot be disabled. Programs will invoke them with a special instruction, specifying which interrupt number they want to invoke. It will be possible to invoke any interrupt, from 0 to 127. The procedure for invoking software interrupts will be the same, except for the fact that Interrupts Flag will be left intact by the CPU.

Clearing the "interrupts allowed" flag will *not* cancel pending interrupts. It will only forbid acceptance of new hardware interrupts. As said above, this flag will have no effect on software interrupts.

8.2.2. I/O ports

The CPU will support operations with input/output ports. There will be 256 I/O ports. All devices outside the CPU will communicate with the CPU using input/output ports.

If a program wants to receive data from a device, it will use a special instruction which will cause the CPU to put port number on the Address Bus and send a "port read" signal down the Control Bus. External devices, on receiving this signal, will test the Address Bus to see if their port number was specified. If they decide that they want to respond to this port read operation, they put one word of data on the Data Bus. That is the data that the destination of the instruction will receive.

If a program wants to send data to an external device, it will call another instruction which will cause the CPU to put port number on the Address Bus, the specified data on the Data Bus, and send a "port write" signal down the Control Bus. On receiving this signal, external devices will check if their port number was specified, and take the data from the Data Bus if they decide to.

The computer will be set up in such a way that no two devices use same port numbers. In this case this is not a major concern because the virtual "computer" will be set up by the system developer once, with all ports assigned to devices without clashes, and the user need not be aware of that. See (<u>Peripherals.Architecture</u>) section to read more about external devices and their ports.

8.3. Instruction set

Below is a list of instructions that the CPU will support, with their arguments (operands) and a description of what they do.

Operand types

- **m** memory (any addressing mode),
- **R** register (A, B, C, D or E),
- **Rn** register (B, C, D or E),
- A accumulator,
- **1** 16-bit immediate,
- **18** 8-bit immediate,
- **N** immediate as part of the machine code.

8.3.1. Data movement

These instructions move data between registers and memory. They also include stack operations. None of these modify the FLAGS register.

Name	Arguments	Description			
ld	dest, src	Copies contents of src to dest. Allowed dest/src combinations: R/R, R/M,			
		R/I, M/R.			
st	src, dest	Copies contents of src to dest. Allowed src/dest combinations: R/R, M/R,			
		I/R, R/M.			
push	src	Copies contents of src to [SP], then increments SP by 2. Src is type R or I.			
pop	dest	Decrements SP by 2, then copies contents of [SP] to dest. Dest is type R.			
pushpc	-	Pushes pc register onto stack, pointing to after the pushpc instruction.			
pushsp	-	Pushes SP register onto stack. SP value before this operation is pushed.			
pushfl	-	Pushes FLAGS register onto stack.			
popsp	-	Pops SP register from stack.			
popfl	-	Pops FLAGS register from stack.			
sp2b	-	Copies the contents of SP register into B register. Used to access parameters			
		that are passed on stack quickly.			
lea	dest, src	Load effective address. Allowed dest/src combinations: Rn/M. Loads the			
		address calculated for src into register dest.			
xchg	r1, r2	Swaps values in registers r1 and r2 so that value in r1 goes to r2 and vice			
		versa. r1 and r2 are type Rgn.			

8.3.2. Arithmetic

These instructions do addition, subtraction, multiplication etc. All of these set the FLAGS register (flags z, s, o, c; n, p) according with the result.

Name	Arguments	Description			
add	addto,	Adds addwhat to addto, saves result in addto. Allowed addto/addwhat			
	addwhat	combinations: A/I , A/M , A/R , R/A .			
sub	subfrom,	Subtracts subwhat from subfrom, saves result in subfrom. Allowed			
	subwhat	subfrom/subwhat combinations: A/I, A/M, A/R, R/A.			
adc	addto,	Adds addwhat, addto and carry, saves result in addto. Allowed addto/addwhat			
	addwhat	combinations: A/I, A/M, A/R, R/A.			

sbb	anhfrom	Subtractory is from it than subtracto commuting the result correct
SDD	subfrom,	Subtracts subwhat from subfrom, then subtracts carry from the result, saves
	subwhat	final result in subfrom. Allowed subfrom/subwhat combinations: A/I, A/M, A/R,
		R/A.
cmp	left, right	Compares left with right and sets flags so that conditional jumps work
		correctly. E.g. if left <right a="" do="" jl="" jump.="" opeartion="" subtracts<="" td="" the="" then="" will=""></right>
		right from left and discards the result. Allowed left/right combinations:
		A/I, A/M, A/R, R/A.
mul	arg1, arg2	Multiplies arg1 by arg2. Saves result in arg1. Treats values as unsigned
		integers. Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.
div	num, denom	Divides num by denom, saves the integer part of the result in num. Interprets num
		and denom as unsigned integers. Allowed num/denom combinations: A/I, A/M,
		A/R, R/A.
imul	arg1, arg2	Multiplies arg1 by arg2. Saves result in arg1. Treats values as signed integers.
		Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.
idiv	num, denom	Divides num by denom, saves the integer part of the result in num. Interprets num
		and denom as signed integers. Allowed num/denom combinations: A/I, A/M, A/R,
		R/A.
mod	num, denom	Divides num by denom, saves the remainder part of the result in num. Interprets
		num and denom as unsigned integers. Allowed num/denom combinations: A/I,
		A/M, A/R, R/A.
inc	arg	Increments arg, that is adds 1 to it. Arg is type R.
dec	arg	Decrements arg, that is subtracts 1 from it. Arg is type R.
neg	arg	Reverses the sign of arg. This is equivalent to not arg; inc arg; but occupies
	-	only one byte. Arg is type R.

8.3.3. Bitwise

Bitwise operations such as AND, OR, shifts, etc. All of them modify the FLAGS register (flags z, s, c; n, p) according with the result.

Name	Arguments	Description			
not	arg	Bitwise NOT – inverts all bits in arg. Arg is type R.			
and	arg1, arg2	Bitwise AND. Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.			
or	arg1, arg2	Bitwise OR. Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.			
xor	arg1, arg2	Bitwise XOR. Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.			
test	left, right	Performs a bitwise AND operation on left and right and sets the flags			
		according with the result. Result itself is discarded. A/I, A/M, A/R, R/A.			
lshr	arg, num	Shifts ³ bits in arg by num to the right. Low-order bit goes to carry, high-order			
		bit becomes zero. Allowed arg/num combinations: A/Rn, Rn/N.			
lshl	arg, num	Shifts ³ bits in arg by num to the left. High-order bit goes to carry, low-order bit			
		becomes zero. Allowed arg/num combinations: A/Rn, Rn/N.			
ashr	arg, num	Shifts ³ bits in arg by num to the right. Low-order bit goes to carry, high-order			
		bit stays the same. Allowed arg/num combinations: A/Rn, Rn/N.			
ashl	arg, num	Entirely equivalent to lsh1.			
ror	arg, num	Rotates ³ bits in arg by num to the left. Low-order bit goes to high-order bit and			
		carry. Allowed arg/num combinations: A/Rn, Rn/N.			
rol	arg, num	Rotates ³ bits in arg by num to the left. High-order bit goes to low-order bit and			
		carry. Allowed arg/num combinations: A/Rn, Rn/N.			
rcr	arg, num	Rotates ³ bits in arg by num to the left through carry. Carry goes to high-order			
		bit and low-order bit goes to carry. Allowed arg/num combinations: A/Rn, Rn/N.			
rcl	arg, num	Rotates ³ bits in arg by num to the left through carry. Carry goes to low-order bit			
		and high-order bit goes to carry. Allowed arg/num combinations: A/Rn, Rn/N.			
bswp	src	Swaps bytes in src so that the high-order byte becomes the low-order byte and			
		vice versa. Src is type R.			

8.3.4. Flags

These operations are used to modify the FLAGS register.

Name	Arguments	Description		
stz	-	Sets zero flag.		
clz	-	Clears zero flag.		
stc	-	Sets carry flag.		
clc	-	Clears carry flag.		
sto	-	Sets overflow flag.		
clo	-	Clears overflow flag.		
sts	-	Sets sign flag.		
cls	-	Clears sign flag.		
sti	-	Sets interrupt flag.		
cli	-	Clears interrupt flag.		

8.3.5. Branching

These are all operations that change execution order. They change IP register (and CS where applicable), so the next instruction to be executed changes as well.

Name	Arguments	Description
jg,	addr	Jumps to addr if $z = 0$ and $s = 0$. Addr is an absolute address of type M.
jnle		
jl,	addr	Jumps to addr if $s \ll 0$. Addr is an absolute address of type M.
jnge		
jge,	addr	Jumps to addr if $s = o$. Addr is an absolute address of type M.
jnl	. 11.	
jle,	addr	Jumps to addr if $z = 1$ and $s \iff 0$. Addr is an absolute address of type M.
jng jz,	addr	Jumps to a data if = -1
je,	auui	Jumps to addr if $z = 1$. Addr is an absolute address of type M.
jnz,	addr	Jumps to addr if $z = 0$. Addr is an absolute address of type M.
jne	uuu	
jc	addr	Jumps to addr if $c = 1$. Addr is an absolute address of type M.
jnc	addr	Jumps to addr if $c = 0$. Addr is an absolute address of type M.
jo	addr	Jumps to addr if $o = 1$. Addr is an absolute address of type M.
jno	addr	Jumps to addr if $o = 0$. Addr is an absolute address of type M.
js	addr	Jumps to addr if $s = 1$. Addr is an absolute address of type M.
jns	addr	Jumps to addr if $s = 0$. Addr is an absolute address of type M.
jmp	addr	Unconditional jump. Addr is an absolute address of type M.
call	addr	Pushes IP registers onto stack; then jumps to addr. Addr is an absolute
		address of type M.
ret	-	Pops data from stack to IP (i.e. does a jump to address on stack)
int	num	Initiates software interrupt num. Num is type 18.
iret	-	Return from interrupts handlers. Pops FLAGS and IP from stack.
halt	-	Brings processor to a halt. In this project this instruction will stop simulation.

8.3.6. Input/output

This section contains operations that send and receive data via input/output ports.

Name	Arguments	Description
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in	dest, prt	Reads data from port prt and places it to dest. Dest/prt can be following
		combinations: R/R, R/18.
out	prt, src	Writes data src to port prt. Allowed prt/src combinations: R/R, R/I, I8/R,
		18/1.

8.3.7. Other

Name	Arguments	Description
nop	-	No operation. The CPU goes on to fetch next instruction after fetching this one.

8.3.8. Notes

³ Shifts/rotations with num greater than 1 are equivalent to several shifts/rotations by 1. Only 4 low-order bits matter in num operand. Therefore, the maximum number of shifts/rotates in one operation is 15 and the minimum is 0.

8.4. Machine codes

When an assembly language program is assembled, it will be converted into machine codes. The CPU will be able to understand only those codes. This section describes the format of machine codes which this CPU will use, including instruction codes, operand formats etc.

8.4.1. Conventions

All instructions will be listed in a table containing the following fields: Instruction, Binary, Hexadecimal, Length and Operands.

Instruction field

This field will contain assembly language instruction with operands. Part of the *Instruction* field will be in bold dark red font – that is the part coded in the first instruction byte. Bright red font will show the part described in extra bytes.

Binary and Operands field

The operands will be shown in bright red font in the *Operands* field. Each bit of the machine code (including the first byte) may be shown as 1, 0 or a lower-case letter. Letters will be used to show that several different options are available and their meaning is described in the *Operands* field. Also, a combination of upper-case letters can be used to describe a whole byte in Operands. Underscore _ will separate bits in the same byte, forward slash / will separate bytes.

Operand types

Operands shown in the *Instruction* field consist of a mnemonic followed by a list of operand types. Operand types can consist of one to three letters, optionally followed by a number when there are two operands of the same type. Allowed types are:

- M memory (any addressing mode),
- **R** register (A, B, C, D or E),
- **Rn** register (B, C, D or E),
- A accumulator,
- **1** 16-bit immediate,
- **18** 8-bit immediate,
- **N** immediate as part of the machine code.

Memory addressing

Wherever an operand is a pointer to a memory cell, one or more bytes will be added *to the end* of the whole instruction. They can describe any addressing mode available, and will be shown as type M in the *Instruction* field. Nothing will be said about it in the *Operands* field other than to show their presence with the MEM word. The structure of those bytes is described in section (Memory Addressing Bytes) below.

8.4.2. Registers

All general purpose registers except the will have a binary ID code associated with them. This code will be used as part of instructions to show that a specific register should be used. The IDs are as follows:

Name	Binary ID
В	000
С	001
D	010
Е	011

The accumulator doesn't have an ID associated with it. This is due to the architecture of the CPU – as can be seen in <u>fig.1</u>, the accumulator stands very separate from the other registers – inside the ALU block. Either a separate instruction or a separate bit is required to specify operations on the accumulator.

8.4.3. Memory addressing bytes

In all instructions operating on memory addresses at least one byte will be added to the end of the machine code describing the way the address should be calculated. The CPU will use this first byte to decide which addressing mode it is dealing with. The format of this byte and other bytes if any is described below. In all format descriptions question mark ? will represent an unused bit, which does not matter and can be set to anything. Square brackets [] will enclose optional bytes. For a description of memory addressing modes, refer to (CPU.Architecture) section.

Direct (via Immediate)

The length will always be 3 bytes. The format is as follows: 00_?????_0 /Y1 /Y2

Bytes Y1Y2 are the offset constant.

Indirect (via Immediate)

The length will always be 3 bytes. The format is as follows: $00_{????_1}$ /Y1 /Y2 Bytes Y1Y2 are the offset constant.

Indirect (via Register)

The length will always be 1 byte. The format is as follows: 01_????_rr RR is the binary ID code for the register used.

Indexed

The length will vary from 1 to 3 bytes. The format is as follows: 1_b_p_?_mm_rr [/Y1 /Y2]

B specifies whether to use the base register (1 means use). P specifies if any offset bytes are present (1 means they are). RR specifies the binary ID for the offset register. MM is the scaling factor -00 if RR should not be taken into account, 01 if it is to be multiplied by 1, 10 - by 2, 11 - by 4. Y1Y2 is the offset constant.

8.4.4. Instructions

Instruction	Binary	Hex	Len	Operands
adc A, I	1000 0111	87	3	y1/y2. y1y2 is the constant I.
adc A,M	1000 1000	88	2-5	
adc A,R	1000 0110	86	2 3	MEM.
adc R,A	1000 0110	86	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.
add A,I	1000 0001	81	3	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
add A,M	1000 0001		2-5	Y1/Y2. Y1Y2 is the constant I.
-		82		MEM.
add A,R	1000 0000	80	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.
add R,A	1000 0000	80	2	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
and A,I	1011 0001	В1	3	Y1/Y2. Y1Y2 is the constant I.
and A,M	1011 0010	В2	2-5	MEM.
and A,R	1011 0000	в0	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.
and R,A	1011 0000	в0	2	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
ashl A,Rn	1100 0010	C2	2	0_xx_?_???. xx is ID for Rn.
ashl Rn,N	1100 0010	C2	2	1_xx_a_nnnn. xx is ID for Rn. If a=1 then Rn is accumulator.
				nnnn is the constant N.
ashr A,Rn	1100 0011	C3	2	0_xx_?_????. xx is ID for Rn.
ashr Rn,N	1100 0011	C3	2	1_xx_a_nnnn. xx is ID for Rn. If a=1 then Rn is accumulator.
				nnnn is the constant N.
bswp A	1001 1111	9F	1	-
bswp Rn	1101 11 xx	DC-DF	1	xx is ID for Rn.
call M	0111 0001	71	2-4	MEM.
clc	0101 0111	57	1	-
cli	0111 0111	77	1	-
clo	0110 0101	65	1	-
cls	0110 0111	67	1	-
clz	0101 0101	55	1	-
cmp A,I	1000 1101	8D	3	$\frac{y1}{y2}$. $\frac{y1y2}{y1y2}$ is the constant I.
cmp A,M	1000 1110	8E	2-5	MEM.
cmp A,R	1000 1100	8C	2	????_0_a_xx. xx is ID for R. If $a=1$ then R is accumulator.
cmp R,A	1000 1100	8C	2	????_1_a_xx. xx is ID for R. If $a=1$ then R is accumulator.
dec A	1010 1101	AD	1	
dec Rn	1010 01 xx	A4-A7	1	xx is ID for Rn
div A,I	1001 0100	94	3	$\frac{1}{1}$ $\frac{1}$
div A,M	1001 0101	95	2-5	
div A,R	1001 0011	93	2	MEM. ????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.
div R,A	1001 0011	93	2	
halt	0111 0101	75	1	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
idiv A,I	1001 1010	9A	3	-
idiv A,M	1001 1010	9B	2-5	Y1/Y2. Y1Y2 is the constant I.
idiv A,R	1001 1011	9B 99	2-5	MEM.
idiv R,A	1001 1001	99	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.
idiv R,A imul A,I	1001 1001 1001			????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
		97	3	y_1/y_2 . y_1y_2 is the constant I.
imul A,M	1001 1000	98	2-5	MEM.
imul A,R	1001 0110	96	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.
imul R,A	1001 0110	96	2	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
in A,18	1110 0101	E5	2	Y. Y is the constant 18.
<pre>in Rg1,Rg2</pre>	1110 0100	E4	2	00_a1_a2_xx_yy. xx is ID for Rg1, yy is ID for Rg2. a1=1
				makes $Rg1$ accumulator, $a2=1 - Rg2$.
in Rgn,18	1101 10 xx	D8-DB	2	Y. Y is the constant 18. xx is ID for Rgn
inc A	1010 1100	AC	1	-
inc Rn	1010 00 xx	A0-A3	1	xx is ID for Rn
int 18	0111 0100	74	2	Y. Y is the constant 18.
				•

Instruction	Binary	Hex	Len	Operands
iret	0111 0011	73	1	-
јс М	0101 0010	52	2-4	MEM.
jg M	0100 0000	40	2-4	MEM.
jge M	0100 0011	43	2-4	MEM.
jl M	0100 0010	42	2-4	MEM.
jle M	0100 0001	41	2-4	MEM.
jmp M	0111 0000	70	2-4	MEM.
jnc M	0101 0011	53	2-4	MEM.
jno M	0110 0001	61	2-4	MEM.
jns M	0110 0011	63	2-4	MEM.
jnz M	0101 0001	51	2-4	MEM.
јо М	0110 0000	60	2-4	MEM.
js M	0110 0010	62	2-4	MEM.
jz M	0101 0000	50	2-4	MEM.
ld A,I	0010 0100	24	3	Y1/Y2. Y1Y2 is the constant I.
ld M,R	0010 0111	27	3-5	?????_axx/MEM. xx is ID for R. A makes R accumulator.
ld R,M	0010 0110	26	3-5	?????_axx/MEM. xx is ID for R. A makes R accumulator.
1d R1,R2	0010 0101	25	2	a1_a2_?xx_?yy. xx is binary ID for R1, yy – for R2, a1 makes
				R1 accumulator, a2 – R2.
ld Rn,I	0010 00 xx	20-23	3	Y1/Y2. xx is binary ID for Rn, Y1Y2 is the constant I.
lea A,M	0000 1100	00	2-4	MEM.
lea Rn,M	0000 10 xx	08-0B	2-4	MEM. xx is ID for Rgn.
lshl A,Rn	1100 0000	C0	2	0_xx_?_???. xx is ID for Rn.
lshl Rn,N	1100 0000	C0	2	$1_{xx_a_nnnn. xx}$ is ID for Rn. If a=1 then Rn is accumulator.
lshr A,Rn	1100 0001	C1	2	0_xx_?_???. xx is ID for Rn.
lshr Rn,N	1100 0001	C1	2	$1_{xx_a,nnnn}$, xx is ID for Rn. If a=1 then Rn is accumulator.
				nnnn is the constant N.
mod A,I	1001 1101	9D	3	$\frac{y1}{y2}$. $y1y2$ is the constant I.
mod A,M	1001 1110	9E	2-5	MEM.
mod A,R	1001 1100	9C	2	<pre>????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.</pre>
mod R,A	1001 1100	9C	2	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
mul A,I	1001 0001	91	3	<u>y1/y2</u> . y1y2 is the constant I.
mul A,M	1001 0010	92	2-5	MEM.
mul A,R	1001 0000	90	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.
mul R,A	1001 0000	90	2	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
neg A	1010 1110	AE	1	-
neg Rn	1010 10 xx	A8-AB	1	xx is ID for Rn
nop	1000 1111	8F	1	-
not A	1010 1111	AF DC DE	1	-
not Rn	1011 11 xx	BC-BF	1	xx is ID for Rn
or A,I	1011 0100	B4	3	Y1/Y2. Y1Y2 is the constant I.
or A,M	1011 0101	B5	2-5	MEM.
or A,R	1011 0011	B3	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.
or R,A	1011 0011	B3	2	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.
out A,I out 18,A	1101 0101 1101 0110	D5 D6	3	Y1/Y2. Y1Y2 is the constant I.
out 18,A	1101 0110	D6 D7	2	Y. Y is the constant 18.
out 18,1 out 18,Rgn	1101 0111 1101 00 xx	D7 D0-D3	4	$y_1/y_2/y_3$. y_1 is the constant 18. y_2y_3 is the constant 1.
out 10, Rgi	1101 0100	D0-D3 D4	2	Y. Y is the constant 18. xx is ID for Rgn
Rg1,Rg2	TTOT 0100	FU	4	00_a1_a2_xx_yy. xx is ID for Rg1, yy is ID for Rg2. a1=1 makes Rg1 accumulator, a2=1 - Rg2.
out Rgn,I	1110 00 xx	E0-E3	3	Y1/Y2. Y1Y2 is the constant I. xx is ID for Rgn
pop A	0001 0001	11	1	-
pop Rn	0000 01 xx	04-07	1	xx is ID for Rn.
popfl	0001 0111	17	1	-

Instruction	Binary	Hex	Len	Operands		
popsp	0001 0110	16	1	-		
push A	0001 0000	10	1	-		
push I	0001 0010	12	3	Y. Y is the constant I.		
push Rn	0000 00 xx	00-03	1	xx is ID for Rn.		
pushfl	0001 0101	15	1	-		
pushpc	0001 0011	13	1	-		
pushsp	0001 0100	14	1	-		
rcl A,Rn	1100 0110	CG	2	0_xx_?_????. xx is ID for Rn.		
rcl Rn,N	1100 0110	CG	2	$1_xx_a_nnnn$. xx is ID for Rn. If a=1 then Rn is accumulator.		
				nnnn is the constant N.		
rcr A,Rn	1100 0111	C7	2	0_xx_?_????. xx is ID for Rn.		
rcr Rn,N	1100 0111	C7	2	$1_xx_a_nnnn$. xx is ID for Rn. If a=1 then Rn is accumulator.		
				nnnn is the constant N.		
ret	0111 0010	72	1	-		
rol A,Rn	1100 0100	C4	2	0_xx_?_????. xx is ID for Rn.		
rol Rn,N	1100 0100	C4	2	$1_xx_a_nnnn$. xx is ID for Rn. If a=1 then Rn is accumulator.		
				nnnn is the constant N.		
ror A,Rn	1100 0101	C5	2	0_xx_?_????. xx is ID for Rn.		
ror Rn,N	1100 0101	C5	2	$1_xx_a_nnnn$. xx is ID for Rn. If a=1 then Rn is accumulator.		
				nnnn is the constant N.		
sbb A,I	1000 1010	8A	3	$\frac{1}{1}$		
sbb A,M	1000 1011	8B	2-5	MEM.		
sbb A,R	1000 1001	89	2	????_0_a_xx. xx is ID for R. If $a=1$ then R is accumulator.		
sbb R,A	1000 1001	89	2	$???_1_a_xx. xx$ is ID for R. If a=1 then R is accumulator.		
sp2b	0000 1111	OF	1	$\frac{1}{2} = \frac{1}{2} = \frac{1}$		
st I,A	0011 0100	34	3	$\frac{1}{1/1}$ $\frac{1}{1}$ $$		
st I,Rn	0011 00 xx	30-33	3			
st M,R	0011 0110	36	3-5	Y1/Y2. xx is binary ID for Rn, Y1Y2 is the constant I.		
st R,M	0011 0111	37	3-5	????_axx/MEM. xx is ID for R. A makes R accumulator.		
st R2,R1	0011 0101	35	2	?????_axx/MEM. xx is ID for R. A makes R accumulator.		
DC 112,111	0011 0101	55	2	a1_a2_?xx_?yy. xx is binary ID for R1, yy – for R2, a1 makes		
stc	0101 0110	56	1	R1 accumulator, a2 – R2.		
sti	0111 0110	76	1	-		
sto	0110 0100	64	1	-		
sts	0110 0110	66	1	-		
stz	0101 0100	54	1	-		
sub A,I	1000 0100	84	3	-		
sub A,M	1000 0100	85	2-5	Y1/Y2. Y1Y2 is the constant I.		
sub A,R	1000 0101	83	2-5	MEM.		
sub R,A	1000 0011	83	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.		
test A,I	1000 0011	BA	3	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.		
test A,M	1011 1010	BB	3 2-5	Y1/Y2. Y1Y2 is the constant I.		
test A,M				MEM.		
	1011 1001	B9	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.		
test R,A	1011 1001	B9	2	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.		
xchg A,Rn	1111 00 xx	F0-F3	1	xx is ID for Rgn.		
xchg b,c	1111 0101	F5	1	-		
xchg b,d	1111 0110	F6	1	-		
xchg b,e	1111 0111	F7	1	-		
xchg c,d	1110 0110	E6	1	-		
xchg c,e	1110 0111	E7	1	-		
xchg d,e	1111 0100	F4	1	-		
xor A,I	1011 0111	B7	3	Y1/Y2. Y1Y2 is the constant I.		
xor A,M	1011 1000	В8	2-5	MEM.		
xor A,R	1011 0110	Вб	2	????_0_a_xx. xx is ID for R. If a=1 then R is accumulator.		
xor R,A	1011 0110	вб	2	????_1_a_xx. xx is ID for R. If a=1 then R is accumulator.		

8.4.5. Instructions allocation

Below is a table showing allocation of instruction codes to different instructions.

	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F
00	push b	push c	push d	push e	pop b	pop c	pop d	pop e	lea b,M	lea c,M	lea d,M	lea e,M	lea a,M			sp2b
10	push a	pop a	push I	pushpc	pushsp	pushfl	popsp	popfl								
20	ld b,I	ld c,I	ld d,I	ld e,I	ld a,I	ld R,R	ld R,M	ld M,R								
30	st I,b	st I,c	st I,d	st I,e	st I,a	st R,R	st M,R	st R,M								
40	jg M	jle M	jl M	jge M												
50	jz M	jnz M	јс М	jnc M	stz	clz	stc	clc								
60	јо М	jno M	js M	jns M	sto	clo	sts	cls								
70	jmp M	call M	ret	iret	int I8	halt	sti	cli								
80	add a,RN add RN,a	add a,I		sub a,RN sub RN,a	sub a,I	sub a,M	adc a,RN adc RN,a	adc a,I		sbb a,RN sbb RN,a	sbb a,I	sbb a,M	cmp a,RN cmp RN,a	cmp a,I	cmp a,M	nop
90	mul a,RN mul RN,a	mul a,I		div a,RN div RN,a	div a,I	div a,M	imul a,RN imul RN,a			idiv a,RN idiv RN,a		idiv a,M	mod a,RN mod RN,a	mod a,I	mod a,M	bswp a
A0	inc b	inc c	inc d	inc e	dec b	dec c	dec d	dec e	neg b	neg c	neg d	neg e	inc a	dec a	neg a	not a
в0	and a,RN and RN,a	and a,I		or a,RN or RN,a	or a,I	or a,M	xor a,RN xor RN,a	xor a,I	xor a,M	test a,RN test RN,a	test a,I	test a,M	not b	not c	not d	not e
			ashl a,RN ashl RN,N			ror a,RN ror RN,N	rcl a,RN rcl RN,N	rcr a,RN rcr RN,N								
D0	out I8,b	out I8,c	out I8,d	out I8,e	out RN,RN	out a,I	out I8,a	out I8,I	in b,I8	in c,I8	in d,I8	in e,I8	bswp b	bswp c	bswp d	bswp e
ΕO	out b,I	out c,I	out d,I	out e,I	in RN,RN	in a,I8	xchg c,d xchg d,c	xchg c,e xchg e,c								
FO	xchg a,b	xchg a,c	xchg a,d	xchg a,e		xchg b,c xchg c,b	xchg b,d xchg d,b	xchg b,e xchg e,b								

Legend:

Param	Description
R	Any register (B, C, D, E)
RN	Any register (A, B, C, D, E)
Ι	16-bit immediate constant
I8	8-bit immediate constant
М	Memory addressing

Color	Category	Color	Category	Color	Category
	Data movement		Flags		Other
	Arithmetic		Branching		
	Bitwise		Input/output		

9. Peripherals

A computer consists of the CPU at its centre, main memory and peripheral devices surrounding them. As the purpose of this system is not only teaching people assembly language but also what different components of a computer do, the system should show some devices and how they interact with the CPU.

9.1. Computer structure

The structure of the computer being simulated is *very* simplified. No system devices are shown at all, and the ones that are shown are those that a typical user would be aware of plus a component linking them to the CPU known as *I/O controller*.

The structure of the computer is shown in $\underline{fig.2}$.



Fig.7.1. Computer structure

Components shown are: CPU, buses, RAM, video card and screen, keyboard controller and keyboard, speaker controller and speaker.
9.2. CPU

The CPU in the diagram is the Central Processing Unit. A whole section (<u>CPU</u>) was devoted to describing how it works, so there is nothing else to discuss here.

9.3. Buses

The imaginary computer will use buses as communication lines between the CPU, RAM and external devices. There will be three buses, just as in a generic computer:

- Data Bus
- Address Bus
- Control Bus

Buses are relatively simple devices, and because this is technical documentation, it is assumed that the way buses work need not be described. So buses won't be discussed any further here.

9.4. RAM

RAM (Random Access Memory) is the place where the program being executed and its data are stored. Memory organisation is discussed in detail in the (<u>CPU.Architecture</u>) section.

Memory is connected to data, address and control buses. The sole purpose of the RAM "device" is to provide the CPU with a place to store the data it needs. Therefore, RAM supports only two functions: memory read and memory write.

Read

When the CPU wants to read data from memory, it will put the address of the first byte to be read on the address bus, and then send a Memory Read signal down the Control Bus. On receiving this signal, the RAM device will read the address from the address bus and put the two bytes starting at the specified address on the Data Bus.

Write

When the CPU wants to write data to memory, it will put the address of the memory location where the data is to be written on the address bus, the actual data to be written on the data bus, and then send a Memory Write signal. On receiving this signal, the RAM device will read both buses and write the data from the data bus to the specified address.

9.5. Video controller

Video card is the component linking the monitor and the CPU. Video card will be assigned several input/output ports to show how the CPU would normally work with external devices.

The image currently displayed on the screen will be stored in RAM in an area 4096 bytes long. Video controller will have an internal register to store the pointer to the

first byte of video memory, which will be E000h by default. The user will be able to change this register easily, thus allowing for such techniques as page switching.

To give user control as to when a picture is formed and ready to be displayed, video controller will have a mode in which it does not reflect changes to RAM until explicitly told to do so. See below for further detail.

Screen modes

There will be several screen modes supported by the video card. The only limitation imposed on it is that the whole video memory should fit into 4096 bytes. The following screen modes will be available:

- 01h: Monochrome text; 1 byte per char; 40x15 characters Every byte represents one character's ASCII code.
- 02h: Color text; 2 bytes per char; 40x15 characters

Every two bytes represent one character's ASCII code and color. The first byte in the pair is the character's ASCII code, the second one – its colour. The color byte format is: LRGB lrgb, where R, G and B are Red, Green and Blue components respectively, L is a brightness bit, uppercase means background color, lowercase – text color.

- 03h: Monochrome graphics; 1 bit per pixel; 208x156 pixels Every byte describes eight pixels. If a bit is set, color seen will be white; otherwise – black.
- 04h: 16 color graphics; 4 bits per pixel; 104x78 pixels Every byte describes two pixels. The format is: LRGB, where R, G and B are Red, Green and Blue components respectively, L is a brightness bit.
- 05h: 256 color graphics; 8 bits per pixel; 74x55 pixels; paletted Every byte describes one pixel. The color that is seen on screen will be taken from a palette array inside the video controller memory which is 256x3 bytes long. That is the palette memory, which stores three bytes (RGB) for every color in this mode.
- 07h: 16M color graphics; 24 bits per pixel; 42x32 pixels Every three bytes describe one pixel. The format is, RGB where R, G and B are *bytes* describing respective colors.

To switch between different modes the programmer will write a word with mode number to a specific port, described below.

Input/Output ports

Screen mode port 50h

Writing screen mode number to this port will cause the video controller to switch screen modes. If it receives any other word apart from valid screen mode numbers, it will ignore it. The changes will be reflected immediately, even in manual refresh mode.

Reading from this port will cause the video controller to return its current screen mode.

Palette port 51h

To change an entry in the palette array, programs will write two words to this port. The first one will contain palette entry number in the low-order byte and the red component in the high-order byte. The second word will contain green and blue components in low- and high-order bytes respectively. Note that once sent to the video controller, palette cannot be read from it. Also, palette only influences images in screen mode 05h.

Memory port 52h

Writing to this port will change the offset to video memory buffer in RAM. The changes will be reflected immediately. That is, even in manual refresh mode the screen will be updated to reflect changes to video memory.

If a program reads from this port, it will receive current pointer to video memory.

Refresh port 54h

Writing 0 to this port will disable auto screen refresh, so changes to video RAM will only be reflected when the programmer wants to. Writing 1 will enable auto screen refresh, so the screen will be updated every once in a while. Writing anything else will force the screen to be refreshed.

Reading from this port will return either 0 or 1 to indicate the state of auto refresh.

9.6. Keyboard controller

Keyboard controller will operate through I/O ports. The user will be able to get pending keys via an input port. Keyboard controller will also send a specific interrupt every time a key is pressed.

Key In port 60h

The programmer will read from this port to get the pending key code. If no key is pending, keyboard controller will return OFFFFh. As soon as a pending key is read once, keyboard controller will forget about that key. Only one key can be pressed at one time, and only Key Down events will be recognised. There will be no way of determining whether a key is still down. This should not be a problem unless someone decides to write a game in this system. Considering that the interpreter will most probably be not fast enough for a game, advanced keyboard features should not be required.

Every time a key is pressed, the keyboard controller will request interupt 1 (IRQ1). If accepted by the CPU, this IRQ will cause interrupt 1. If not accepted by the CPU, controller will keep sending requests and ignoring all key inputs until the request is successful.

Keyboard scancodes

A scancode is the code returned by the keyboard controller when a given key is pressed. The table of scancodes for this keyboard controller is given below.

Scancode	Hex	Key
0	00	A
1	01	В
2	02	С
3	03	D
4	04	Е
5	05	F
6	06	G
7	07	Н
8	08	I
9	09	J
10	0A	K
11	0B	L
12	0C	М
13	0D	N
14	0E	0
15	OF	Р
16	10	Q
17	11	R

Scancode	Hex	Key
18	12	S
19	13	Т
20	14	U
21	15	V
22	16	W
23	17	Х
24	18	Y
25	19	Z
26	1A	
27	1B	Enter
28	1C	Spacebar
29	1D	=
30	1E	0
31	1F	1
32	20	2
33	21	3
34	22	4
35	23	5

Scancode	Hex	Key
36	24	6
37	25	7
38	26	8
39	27	9
40	28	Numpad .
41	29	/
42	2A	*
43	2B	-
44	2C	+
45	2D	Left
46	2E	Right
47	2F	Up
48	30	Down
49	31	Circle
50	32	Square
51	33	Triangle

9.7. Speaker

Speaker can be used as a means of giving signals to the user easily. It will be much easier and faster to switch speaker state than output something on the screen. Speaker will be operated through a single port -80h.

Speaker port 80h

Writing 0 to this port will set speaker to the low state. Writing 1 will set speaker to the high state. Writing any other number will set speaker frequency to 20/65536*W where W is the word sent to this port.

Reading from this port will return the last word written to this port.

10. Assembly language

Generally, the syntax of assembly language in this project should be as similar to the one used in the exams as possible. But because not one examining board is consistent even with its own past papers, this syntax will be only approximately like that in exams or textbooks.

10.1. Statements

A *statement* is the smallest unit of division of programs which can be taken out of context and still have a meaning. The whole statement has to be written on one line, and there can only be one statement on every line.

The structure of a typical statement is shown below:

[label:] [data | command] [;comment]

Square brackets [] enclose elements that are optional, vertical bar | indicates that there are two possibilities, and only one can be present. Only spaces or tabs can separate the elements, but there can be as many of those as needed. Empty lines are allowed since there is no element which is not optional. No part of assembly language syntax is case sensitive.

10.2. Labels

A label is a pointer to a part of code which enables the programmer to reference that part in instructions, letting the compiler do all the calculations. Labels can be used to reference code to use with jump and call instructions, or they can be used to reference data in data movement instructions. See (<u>Referencing</u>) for information about how to reference labels in operands.

A label has to start with a letter and can contain letters, numbers and underscores. Every label has to end with a colon : followed by at least one space or tab character, and there can be no whitespace between the name of the label and the colon.

10.3. Data declarations

When a program is compiled, compiler generates code for every operation. To tell compiler to insert specific bytes in compiled code programmers can use keywords db, dw or ds. The purpose of these keywords is to reserve some space in the machine code for data. If the programmer declares a label pointing to that data, the data effectively becomes a variable. To see how to use variables as operands see sections (Offset macro) and (Referencing). db stands for "declare byte" and reserves 1 byte of space, dw – "declare word" and reserves 2 bytes, ds – "declare string", the number of bytes depends on the length of the string, between 0 and 255.

Every data declaration keyword should be followed by initialisation sequence to tell the compiler the initial contents of the data. db should be followed by a numerical constant that fits into 8 bits, dw – by a numerical constant that fits into 16 bits, ds – by a string literal, described below. Alternatively, each of these keywords may be

followed by a question mark ? to indicate uninitialised variable. Compiler will then initialise db and dw with zeroes, and ds – an empty string, reserving no space for it in the code.

String literals are constants, but unlike numerical constants, string literals can be used in only one case – to initialise ds data declaration. String literals must be enclosed by quotation marks ". Everything between the marks is the value of string literal. Double quotation marks can be used to include quotation marks as part of string literal's value. When compiled, every byte of string literal's value will be copied directly to machine code, as is. Note that, although string literals cannot be used where numerical constants can, variables declared as string constants are completely identical to those declared as numerical constants.

Because the data will be stored together with the code, the programmer will have to make sure that data is not executed accidentally.

10.4. Commands

Commands are basic instructions which the CPU can process, written in a form readable by humans. Every command has the following syntax:

```
opcode [operand1 [, operand2 ] ]
```

Opcode is a symbolic form of writing a CPU instruction; it should be one of the opcodes listed in section (<u>CPU.Instruction Set</u>). Operands should be separated by a comma, and there has to be at least one space or tab character between opcode and operand1. More tabs or spaces can be used between the elements if needed. Operands are described in detail in the next section.

10.5. Operands

Operands can be of the following types:

- Register
- Immediate
- Memory

Register

This operand type corresponds to the Register addressing mode. It can be A, B, C, D, E. Sometimes only specific registers can be used – it depends on operand. If register is source, data passed to the instruction is the contents of the register. If register is destination, result is written to the register.

Registers can appear as separate operands, or can be used as part of memory operands described below.

Immediate

Immediate operands are numerical constants. Assembly language will support decimal, hexadecimal and binary numbers. It will have to discern whether a constant is a byte or word. The syntax is as follows:

integer[h|b]

Integer is any combination of digits from 0 to 9 and letters from A to F, but it has to start with a digit. If the number is followed by letter h, integer will be interpreted as a hexadecimal number. If it is followed by b, integer will be interpreted as a binary number, and integer should consist of 0's and 1's only. If it is not followed by either h or b, integer is interpreted as a decimal number and should contain only digits from 0 to 9. If the resulting constant is greater than OFFh, it is always interpreted as a 16-bit constant. If it isn't, interpretation depends on the opcode.

Negative immediate constants are allowed. They will be stored in two's complement format, and the way numbers with the highest bit set are interpreted will depend on the opcode.

Immediate constants can stand as separate operands, or can be used as part of memory operands described below.

Memory

Memory type operands are used for indexed, direct and indirect (register/immediate) addressing. The syntax of each one is described below.

Memory Direct

offset

The whole construction should be enclosed in square brackets [], and no spaces are allowed between any sections of the construction. Offset is a 16-bit immediate constant which specifies absolute memory address. Variable name as an operand has Memory Direct operand type.

Memory Indirect Register

register

The whole construction should be enclosed in square brackets [], and no spaces are allowed between any sections of the construction. Register specifies the register holding absolute memory address, and can be B, C, D or E.

Memory Indirect Immediate

offset

The whole construction should be enclosed in double square brackets [[]], and no spaces are allowed between any sections of the construction. Offset is a 16-bit immediate constant which specifies absolute memory address. The CPU will first read the two bytes at [offset] and then use the value read as the final absolute memory address.

Memory Indexed

[base+]register[*scale][+offset]

The whole construction should be enclosed in square brackets [], and no spaces are allowed between any sections of the construction. Register specifies the register holding offset, and can be B, C, D or E. This offset is multiplied by scale, and offset (16-bit immediate constant) can be added to act as base offset. This parameter is often specified with the (Offset macro). If this operand is a source, its value is the contents of the two bytes of memory at the address obtained by calculating the expression. If it is a destination, the result is saved in the two bytes of memory at the address obtained by calculating the expression.

10.6. Offset macro

A *macro* will be supported by the assembly language. Programmers will not be able to define new macros, like in C. The built-in macro will be called <code>offset</code>. It will take one parameter – a variable name – and return its absolute address as if an immediate constant was specified. For example, assuming that my_var is a variable at offset 200h, the following code:

```
ld a,offset(my_var)
```

is equivalent to

ld a,200h

The advantage of using macros is that the address at which a variable is stored depends on the length of all the code preceding that variable. If a programmer uses a constant to specify its address, he will have to manually recalculate the address every time it changes, whereas offset macro will do all calculations for the programmer. The net effect of the substitution is that offset(my_var) is converted to a 16-bit immediate constant type operand.

10.7. References

A *reference* is a special operand type which is used to refer to labels. Only absolute references will be supported in this assembly language for simplicity.

The idea behind references is the same as that behind offset macro – let the compiler calculate the address. Wherever a memory operand is required, variable name can be used to specify the address. Compiler will replace that with [offset(varname)] and compile as described above. The net effect after the substitution is that the variable name is converted to a direct memory addressing operand.

10.8. Comments

In every line of code, everything that follows a semi-colon *;* is completely ignored by the compiler together with the semi-colon. Every line which is empty or consists of only spaces or tabs is ignored as well.

Design

The Real Program

11. Interface

This section describes the interface of the program – the part that the user will see and interact with.

It is probably a good idea to design the whole system in such a way that the user sees the system as a computer, rather than as a program which simulates a computer. This will be taken into account at all stages of interface design.

11.1. Conventions

Numbers

Whenever a number is shown on screen, the user should be able to change its representation between binary, decimal signed/unsigned and hexadecimal. All binary numbers will be followed by a lower-case "b" letter. All hexadecimal numbers will be followed by a lower-case "b" letter and have an extra zero in front of them if they start with a letter.

Window types

There will be three basic window types in the system. To indicate that, every picture of a window in this document will have one of the following icons in the top-left corner:

 \square – windows that make the imaginary computer look like a usual PC does.

 \mathbf{x} – windows that represent internal computer hardware.

■ – windows that provide operating system and software facilities, such as assembly language Integrated Development Environment (IDE), or stack viewer.

11.2. Main window

Main window will be positioned at the top of the screen. It will provide a menu to load and save the project, run the simulation and change global settings. It will also provide access to every other window in the system.

Main window will contain a toolbar, giving the user quick access to the most frequently used functions, such as open/save project, run/pause/stop simulation, and quick access to system's windows. Minimizing the main window will minimise all windows of the system, and closing the main window will shut down the system. The main window will look something like this:

ASMP - noname

_ 8 ×

File View Simulation Options

Fig.9.2. Main window

The menu will have approximately the following structure:

- File
 - o New project
 - Open project
 - o Save project
 - o Exit
- View
 - A list of all windows in the project, separated into submenus by category
- Simulation
 - o Start
 - o Pause
 - o Stop
 - o Reset
 - o Simulation speed
- Options
 - o Complexity level
 - o Number format

11.3. Computer window

This window will simulate the end-user computer. The user will be able to see the computer as if they were sitting in front of a real PC, with a keyboard and a monitor in front of them. They will be able to start or reset a program, and to interface with a running program.



Fig.9.3. Computer window

The user will be able to press any key on the keyboard, which will be sent to the program in an appropriate way. When the user presses a key, that key will become red for a fraction of a second, so that the user knows the key has been pressed.

11.4. Monitor window

This window will duplicate the monitor on the computer window. Apart from the fact that it can be moved around more easily, it will provide scaling facilities and some manual monitor operations such as manual mode switching. The window will look something like this:



Fig.9.4. The Monitor window

The caption of this window will contain information about the current screen mode. The window will be sizable, so that the user will be able to change the scale if necessary. Right-clicking on the window will bring up a popup menu with the following structure:

- Screen mode
 - \circ 1 Monochrome text
 - \circ 2 Color text
 - 3 Monochrome graphics
 - \circ 4 16 color graphics
 - \circ 5 256 color graphics (paletted)
 - \circ 6 65k color graphics
 - \circ 7 16M color graphics
- Scale
 - o 200%
 - o 400%
 - o 800%
- View video RAM (brings up the RAM window)
- View controller (brings up the Video Controller window)

11.5. Keyboard window

This window will duplicate the keyboard in the computer window. The main advantage of having this window separately from the main window is that in this way the keyboard can be placed wherever the user wants to see it.

A possible keyboard window layout is shown below:

⊒ к	eyb	oard	d – "	'J″ (28h)										X
Q	W	Ε	R	Т	Y	U	Ι	0	Р	/	*	-	-	\bigcirc		
Α	S	D	F	G	Η	J	Κ	L		7	8	9				
Ζ	Χ	С	V	В	Ν	М		4	<u>'</u>	4	5	6	+			1
										1	2	3			T	
										()		=	-	↓	-

Fig	9.5.	Keyboard
	/	iii y boui u

The caption of the window will contain information about currently pressed key. That key will also be highlighted with red color. The user will be able to press keys with the left mouse button. Only one key can be pressed at a time.

11.6. Speaker window

This window will just be a tiny window to indicate current state of the speaker. There is a chance that producing real sound through PC speaker will not be possible under VB, so the speaker will flash rapidly to indicate sound.

11.7. RAM window

The RAM window will allow to view and edit the Random Access Memory. It will have to provide such facilities as finding a specific part of memory, jumping to specific addresses (such as current Stack Pointer), show selected byte(s) as different data types (numbers in all representations, disassembled instruction).

The RAM window will consist of two main parts – the memory table and the interpretation panel.

🛠 RAM																>	۲
	8 bit			16	6 bit			Disassembled:									
Hex	B5h			в5	C5h			or a	, [ċ	ls:3	80h+	c*2]					-
Dec	181			46	533			[ds:	3A61	ı] is	52F	Fh					
SDec	-75			-1	9002			A	В	D	В	С	В				
Bin	1011	010)1b	11	00 0	1011	5	C)	C)	C)				
	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F	
0000	EF	02	48	89	44	24	24	75	CA	8B	7C	24	3C	8B	54	24	
0010	38	8B	5C	24	10	8B	74	24	18	8B	4C	24	14	8B	44	24	
0020	7F	85	81	C3	02	C0	00	00	BD	00	00	00	80	66	85	DB	
0030	89	5C	в5	C5	A 0	03	80	6C	24	30	75	29	8D	54	24	28	
0040	52	E8	бA	ED	\mathbf{FF}	\mathbf{FF}	00	00	00	00	00	00	00	00	00	00	
0050	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
0060	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
0070	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
0080	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	

Fig.9.7. RAM window

The memory table will display a part of the RAM contents. It will have a fixed column displaying row address and a fixed row to display byte offset within the row.

The main part will show the byte at the corresponding address. The user will be able to move a cursor around the table. The cell with the cursor will be called the *current* cell and filled with blue. The bytes that would make an instruction starting from the current byte will be shown in blue font. Bytes pointed at by PC and SP register values will be highlighted with red and green background respectively. The beginning of every new segment will be highlighted with an aqua colored row indicating segment number.

The interpretation panel is the top part of the RAM window. The user will be able to choose to show it or hide it. It will show the interpretation of the current byte (and the one that follows) in different number bases and formats for 8 bit values (16 bit values). It will also disassemble and display the instruction that the CPU would execute if IP was pointing at the current byte.

When the user right-clicks anywhere in the window, a popup menu will appear with the following structure:

- Show interpretation panel
- Jump to...
 - ...Current PC
 - o ...Current SP
 - o Stack
 - o Video Memory
- Set...
 - ... PC to current byte
 - o ... SP to current byte
- Fill
- Cut
- Copy
- Paste
- Open stack
- Open variables
- Open disassembler

11.8. Buses

The buses window is intended to show at all times what data is currently going through the buses. It will show the actual data going through the buses and its textual interpretation (where applicable). By right-clicking in the window the user will be able to change number representation (binary, decimal signed/unsigned, hex), and show/hide a special panel which would show a small circle for every wire on every bus, red if a wire is in high logical state and blue for low state. This should help the user to understand the bus concept easily.

When many windows representing actual components are displayed on the screen, they need to be connected with each other through system buses. Rather than showing buses as thick lines going through the screen, which would be very awkward, the buses will not be shown at all. Instead, every window will have a "connector" – a set

of three circles labelled "AB", "DB" and "CB" for Address Bus, Data Bus and Control Bus respectively. It will be implied that all AB circles are connected with each other, and so will be DB and CB circles. Every time data is sent across a bus from one device to another, that data will be shown as a tiny window with the data printed in it flying from the sender connector to the connector on the Buses window and then to all destinations it may go to. The data will skip the Buses window if it is closed (it will fly straight to the destination(s)).

🛠 Sy	buses X	
AB	0	89F5h
DB	0	0000h
CB	0	Memory write
		-

Fig.9.8. Buses window

The left part of the window is the "connector". The right part will display the same text as that in the flying windows. Changing the numerical representation will also affect the representation for the flying windows. The text in this window will not be updated to the new value put by a device until a respective window "lands" on this window's connector.

11.9. CPU window

The CPU window will show internal structure of the CPU and what happens when it executes a program. The window will show the components the CPU consists of and the values of all registers. Fig.9.9 below shows a possible layout of the form.



Fig.9.9. CPU window

The purpose of all components that will be shown is described in detail in section <u>CPU.Architecture</u>. It will also be described in user manual.

There are a lot of different buses in the diagram. Some of them are major buses such as the Internal Data and Address Buses, some are minor and have no special name. In any case, the data flowing through them will have to be shown somehow. The most visual way to do that is to show tiny windows with data flowing along the buses, very similar to the way described in section <u>Interface.Buses</u>.

The contents of all registers, including the internal registers and the decoded contents of the FLAGS register, will be displayed in the diagram to make it easier to see how the CPU operates. Whenever a register's value changes, it will be displayed in red for one clock cycle. Whenever the Control Unit chooses a register for a read/write operation, that register will be displayed in a thick blue frame.

The top right part of the window will display a "connector", as described in <u>Interface.Buses</u>. In this diagram, more than anywhere else, the user will be able to see how data that goes through connectors is used.

There are a lot of components in the diagram, and if the user sees the diagram of the inside of the CPU for the first time, they will be repelled by its apparent complexity. It is therefore absolutely vital that it is possible to choose which components are to be displayed. By default the simplest view will be selected. The three available views will be as follows:

Basic view:

SP and $\pm 1/2$ will be invisible LEA with arrows leading to it will be invisible Base register going to addressing will be invisible Insides of the Control Unit will be invisible Flags register will be invisible MDR and MAR registers will be invisible

A-level view:

LEA with arrows leading to it will be invisible Base register going to addressing will be invisible

Full view:

Everything will be visible

Independent of whether a specific component is turned on or off, the data flow will remain unaffected (with one exception – signals controlling invisible devices will also be invisible). The user will be aware that there is something in the white space where the data comes from, but they will not be able to see that component until they decide they are more familiar with the CPU structure.

If the user right-clicks in the window, a following popup menu will appear:

- Detail level
 - o Basic

- o A-level
- o Full
- Number representation
 - Decimal signed
 - Decimal unsigned
 - o Hexadecimal
- Open registers
- Open flags
- Open control unit
- Open ALU

11.10. Control Unit window

This window will display the status of the control unit to give an interested student a slight idea of how it works. Showing decoded instruction will not be a complicated programming task as the system will have to decode instructions into microinstructions anyway in order to show signals flowing around the CPU.

🛠 Control Unit	X
Cycle: Execute To fetch: 0 To exec: 1	Interrupt Status: FEDCBA98 76543210 00000000 00000000
CIR: 23 10 00	Disasm: ld e, 16
Decoded: reg_se, idb_i	m(0010h), reg_w

Fig.9.10. Control Unit window

At the top of the window the system will show some internal registers, such as the Current Instruction Register (CIR), and how many fetch or execute cycles are left in the current instruction. It will also show the disassembled instruction to make it easier for the user to understand what the control unit is doing. The Decoded part will show a decoded microprogram for current instruction. It will contain one or more microinstructions, consisting of one or more signals for internal control bus.

There will be no complexity level setting for the control unit window. In the Basic mode the user will not be able to open this window at all, whereas in the A-level mode the user will be warned that CPU windows contains all the required information, and Control Unit window is more advanced than required at A-level.

11.11. ALU window

The Arithmetic Logic Unit (ALU) window will teach students how the ALU works by showing what it does, and in some cases animating the operations. Right-clicking in the window will bring up a menu with an option to change number representation. If an animation for the operation is available, the user will be able to run it from the menu.

	×
208Ah	
0003h	
multiply	
inp1, inp2	
619Eh	
	0003h multiply inp1, inp2

Fig.9.11. ALU window

11.12. Video controller

The video controller window will show details about video controller's state, as well as give user some reminder regarding how to work with the controller. This should help the user understand how devices in general and the video controller in particular work.



Fig.9.12. Video controller window

As usual for controller windows, this window will have a bus "connector" on it (see section (<u>Interface. Buses</u>)). To the right of the "connector" the window will display the status of its ports. The bottom part will decode current screen mode. To remind the user what each of the ports does, a hint will pop up every time the user hovers over a port number. Port 50h hint will also display a list of all screen modes.

11.13. Keyboard controller window

The keyboard controller window will show the user current keyboard controller status. This should help the user understand how devices in general and the keyboard controller in particular work.

🛠 Keyboard controller 🛛 🗙									
Keyboard controller Port 60h, IRQ 1									
AB O DB O CB O	Pending key: "I" Last key: "L" IRQ: pending								

Fig.9.13. Keyboard controller window

The window will contain the system buses "connectors" (see section (<u>Interface.</u> <u>Buses</u>)) as well as status information. The window will tell the user whether interrupt request for last key pressed has been accepted and whether the last key has been sent to the program.

11.14. Speaker controller window

The speaker window will be a small window showing current state of the system speaker and its controller. This window will be the simplest of all device controller windows, and so may be used as a good aid in introducing such concepts as external devices and I/O ports.



Fig.9.14. Speaker controller window

Speaker state (high or low) will be indicated in a field at the bottom on the window. Also, when speaker is high the picture of the speaker will become red. This way, a buzzing speaker will be blinking rapidly. Right-clicking anywhere in the window will bring up a menu with two options – switch to high state and switch to low state. The AB, DB and CB labels are "connectors" for the respective system buses (see section (Interface.Buses)).

11.15. Code window

The user will usually write and debug programs in the code window. Using this window the user will be able to load and save their programs on disk, compile them into machine codes, load them into memory and run them, and debug them.

🖽 c	Code –	mu	ltiply.a	sm *		×				
Fil	e E	dit	Run	Tools	Settings Help					
;	A mu	lti	ply pr	ogram d	demo					
		.mo	del		flat					
		.st	ack		downward					
		ld	A, <mark>0</mark>		;A holds current result					
		ld	В,М1		;B holds one multiplier					
		ld	C,M2		;C holds the other multiplier					
⇒ 1	.oop:	tes	t C,1		;Test the low-order bit					
		jz	skip		;Add only if bit is set					
		add	A,B		;Update current result					
		lsh	r C,1		;User next bit in one multiplier					
		lsh	1 B,1		;Increase the other multiplier					
Erro	or (8): II	lega	l combir	nation of	opcode/operand: "test" and "C"					
					Modified Line 3 I	nsert				

Fig.9.15.1. Code window

The main part of the window will be the code editor. To the left of it will be the socalled "gutter", where different types of markers will be displayed (e.g. error, breakpoint, current execution point). The bottom part of the window will display messages, such as error and warning messages, generated by syntax check or the compiler. Double-clicking on a message will show the offending line. In the bottom right corner the user will be able to see file status (modified/saved), current line in the editor and the editing mode (insert/overwrite). The code window will support syntax highlighting, which will greatly assist writing programs.

The integrated help system will show the user help on current instruction in the editor if the user presses F1. If the user selects an error/warning message and presses F1, the system will display a relevant help topic. Full help will be available through the window menu.

There will be a menu at the top of the screen, which will have the structure displayed in <u>fig. 9.15.2</u>.

*	 File New program Open program Save program Save as Open workspace Save workspace 		s Trace table Interface Computer Monitor Keyboard Speaker	*	Help ≻ Full help ≻ Keyword
*	Edit Cut Copy Paste Replace Find Find next		Hardware Speaker Keyboard controller Video controller System buses RAM CPU Control Unit		
*	RunSyntax checkCompileRunPauseStopStepToggle breakpoint	~	 ALU OS/Debug Stack Variables Disassembler Registers Flags 		

Fig.9.15.2. Menu structure for Code window

11.16. Registers subwindow

This window will show all the registers that there are in the CPU. The user will be able to choose whether to show segment and internal registers. The user will also be able to change number representation for the registers.



Vertical layout

Right-clicking in the registers window will bring up a popup menu which will allow to change layout and number representation. If the user right-clicks on a register, the menu will also allow to change its value. Only registers A-F registers can be edited in this way.

11.17. Variables subwindow

The variables window will display the variables declared in the source code. It will show variable address, variable name and its value. The user will be able to choose number representation and variable size.

11.18. Stack subwindow

The stack window will display stack contents, decoded into separate stack elements (unlike in the RAM window where everything is just an array of bytes). The user will see element addresses and data stored in those elements. The user will be able to choose different number representation. The element to which SP points at a given moment will be highlighted with aqua background.

Uar-s (signed word dec)			
Adr	Name	Value	
028	ml	0020h	
030	m2	0846h	
04A	product	000Ah	

Fig.9.17. Variables window

🖽 RAM – stack (word hex) 🛛 🗙				
	/	Stack	element	
	206	38Efh		
	208	7F89h		
	20A	0052h		
	20C	8620h		
	20E	0000h		

Fig.9.18. Stack window

12. Modules

This section describes the modules of which the program will consist. A module is a logically and physically separate unit of a program which either contains a set of specific operations (a procedural module) or code related to window interface (a form module).

12.1. Form modules

Interface windows

- fiMain Main window
- fiComp Computer window
- fiDisplay Display window
- fiKeyboard Keyboard window
- fiSpeaker Speaker window

Hardware windows

- fhCPU CPU window
- fhCU Control Unit window
- fhALU ALU window
- fhRAM RAM window
- fhBus Buses window
- fhVideo Video controller window
- fhKeyboard Keyboard controller window
- fhSpeaker Speaker controller window

OS/Debug windows

- fsCode Code window
- fsRegs Registers window
- fsVars Variables window
- fsStack Stack window

12.2. Procedural modules

- pWinAPI all necessary declarations to use WinAPI functions
- pUtils commonly used procedures not available in standard libraries
- pGlobals global variable and data type declarations
- pExec procedures that are responsible for execution of an instruction
- pCompile assemble a program, syntax check, plus assemble/disassemble a given string procedures
- pIO input/output procedures to simplify interfacing with device modules, providing such functions as PortRead, PortWrite, RequestInterrupt etc.

13. Data structures and globals

It is very hard to think about data structures in advance because they will be severely affected by exact implementation of algorithms. At this stage many points are not clear enough yet, and will only be decided upon in the process of implementation. What makes laying down data structures even harder is that I am not familiar with Visual Basic, and I have never coded anything more advanced than a simple single-windowed program in it. I could use my experience in Delphi programming to compile a list of all global data structures that I would expect, but most probably they will not be entirely applicable to Visual Basic. I am planning to develop data structures while learning about VB coding principles during implementation stage.

I would expect to declare all global variables and structures in a module called pGlobals. I would have a structured variable to hold all information about current project, another one - to hold execution (simulation) state, and probably one for application-related variables.

14. Assembly process

When an assembly language program is compiled into machine code, a specific algorithm is at work. This section describes how the algorithm will work.

14.1. Conventions and terms

Tokens

A token is the smallest unit of program that is meaningful to the compiler. In this assembler, a token will be a string containing no spaces after pass 1 (see below). Therefore, a token can be:

- Label
- Variable declaration type
- Variable initialisation sequence
- Opcode
- Operand

Note that an operand token can not be divided any further. That is, even a complete indexed memory addressing will be parsed into a single token.

14.2. Passes overview

To assemble a program the algorithm works through it several times; each time is called a *pass*. To fully compile a program, three passes are required.

Pass 1. Tokenize

At this stage the source code is converted into a set of tokens which, unlike such languages as C and Pascal, will still be separated into lines. Internally a tokenized program will be stored as an array of token lines, each being an array of strings containing one token each.

Pass 2. Code generation

Each token line is converted into respective machine codes which are added up and saved in a special byte array. Where a reference is used, the program will remember reference name and backpatch its address in pass 3.

Pass 3. Backpatching

All references are replaced by physical addresses.

14.3. Pass 1. Tokenize

This is a relatively simple pass. The following procedures are carried out on every line of code:

- Remove all comments, unnecessary spaces and empty lines
- Split every line into tokens
- Determine token types

- Analyse token patterns (see Token Patterns)
- Prepare references for pass 2.

14.4. Pass 2. Code generation

At this stage machine codes can be generated. The algorithm will go through all token lines and generate respective machine codes. The following procedures will be carried out:

Generate code for correct instructions Issue errors for incorrect instructions Build a list of labels with their physical addresses Build a list of reference backpatch requests.

14.5. Token patterns

Only specific combinations of tokens will be valid. To simplify referring to token types, the following abbreviations will be used:

- Label label
- Variable declaration type vardecl
- Variable initialisation sequence varinit
- Opcode opcode
- Operand operand

Allowed token patterns

label opcode opcode operand opcode operand operand vardecl varinit

Patterns that can be corrected

If a pattern contains label tokens that follow some non-label tokens then a warning is issued and the label tokens are moved to the beginning of the line.

If a pattern contains a label token (or several label tokens) then the token line is split into two or more token lines, with a token line for every label and the last token line being what remains of the original token line. If it has no other tokens, it is removed completely. This procedure generates no warnings.

If a pattern contains only a vardecl token then a varinit token with no initialization will be added, generating a warning.

All other patterns will generate an error. The error message will say that a specific token combination is invalid.

15. Execution process

This section describes the algorithm that executes a program. The algorithm will effectively simulate a clock tick. Below is a flowchart for the algorithm.



Fig.14.1. Execute instruction algorithm

To invoke an interrupt the system will create a microprogram which will load ISP address from the interrupt vector table, save flags and return address on stack and jump to ISP address. As soon as the microprogram is in the buffer, all that will necessary is just to continue running the clock tick procedure.

16. Sample scenario

The purpose of this section is to analyse whether design is acceptable by going through all stages of development of a simple assembly language program in an A-level computing lesson, and thus clearly showing what this project is going to be like.

Suppose that the goal of the lesson is to write a program which will add the values of two variables, v1 and v2 and store the result in a third variable, R. Each student is issued with a copy of user manual for A-level students. The teacher has already taught students some theory about assembly language before, so students have a general idea of what they will be doing.

The students start the system. By default they are in GCSE mode. The teacher tells students to switch to A-level mode. At the top of the screen each student sees a window with a menu and many buttons. The teacher tells them to click on the one saying, "Write a program". A window with a text editor appears. Students can start to code. Let's concentrate on one student, named James.

James knows that there is a special instruction "add" that will instruct the computer to add two numbers. James also knows that normally the result of this operation will be stored in the accumulator. So James types:

```
add V1, V2
st a, R
```

When James clicks "Run", he gets an error message, with the first line highlighted in red, saying "Syntax error in operand OR opcode and operand incompatible. Offending operand: V1". James has no idea what this means, so he asks the teacher. The teacher notices that many people have the same problem and reminds everybody that one can only add together two registers or a register and a variable, but not two variables. So the values of the variables should be loaded into registers first. James remembers that this can be done with the "1d" instruction. So he modifies his code and gets the following:

```
ld b, V1
ld c, V2
add b, c
st a, R
```

James reasons that he should not load anything into the accumulator because the result will be stored there. So he loads the variables into two other registers. When he tries to run this code, he sees that the third line is highlighted, and he gets a similar error message: "Syntax error in operand OR opcode and operand incompatible. Offending operand: c". James seeks for some help from the teacher, but teacher is busy helping someone else and tells James to consult the manual regarding the ADD instruction and the error message James is getting. James reads the description of the ADD instruction, which, among other things, mentions that if one tries to add two operands apart from those that are allowed they will get exactly this error message. So James realises that again, he is trying to add something he is not allowed to add. He notices in an example that two registers will generate an error unless one of them is the

accumulator. James decides that he could load v_2 into accumulator. He modifies his code accordingly:

ld b, V1 ld a, V2 add b, a st a, R

When James tries to run this program, he gets an error message – for the third time. The message is pointing at the first line and saying, "*Undeclared reference: V1*". He gets quite annoyed, but soon remembers that he has to declare the variables v1 and v2 in his code. So he adds to the end of the code:

V1: dw 14 V2: dw 8 R: dw 0

This time the program actually runs. James can see the instruction pointer slowly moving down. But when it reaches the end of the program, it doesn't stop but tries to execute variable declaration, generating an error message saying *"Failed to decode instruction starting with 0Eh"*. James asks the teacher how to make the program stop. The teacher tells him to put a "halt" instruction where he wants the program to end. James updates the program, which by now looks like this:

```
ld b, V1
ld a, V2
add b, a
st a, R
halt
V1: dw 14
V2: dw 8
R: dw 0
```

The program runs successfully and ends by displaying a message, "*CPU halted*". James now wants to see the results of his work. So he clicks on a button in the main window saying, "Variables". He sees that v1 contains 14 and v2 contains 8, just as they should, but R contains 8 instead of 22 as expected. He asks the teacher what to do. The teacher advices James to open the "Registers" window and carefully go through every line of code by pressing F8 and looking at what happens. James does so. He sees that after the first two instructions the registers contain what they should. But after he executes the third instruction he notices that contents of register b changes to red (indicating that value changed), and it is indeed the sum of two numbers. James realises that result goes into b and not the accumulator for some reason, so he changes the fourth instruction to store register b in R. James runs the program again and sees that this time everything works perfectly.

17. File formats

There will be only one file type – assembly language program. The file will have no special format – it will simply store all the code that the user writes in text form, "as is". Among the advantages of this approach are the facts that it will be very simple to implement, and the users will be able to edit their code without having CLab.

18. Security and integrity

This system will store no sensitive data and therefore will require no security measures to be taken.

All data that may suffer loss or corruption is programs saved on disk or the program under development. It is not crucial to make sure that programs that are loaded from disk are error-free, but if any time is available then an integrity check such as a checksum may be implemented. To maintain integrity of the program under development error traps should be used in implementation so that even if something goes wrong the user will still be able to save their work on disk.

19. Design confirmation

Having completed the design of the system, it should be confirmed with the user(s) the system is developed for. I have discussed this design section in detail with my end-user, and below is a list of all modifications that should be made.

- Disassembler window and disassembled instructions in Control Unit, although useful, will not be worth the implementation time necessary to have them
- The user should be able to edit all register values in the Registers and CPU window, as well as variables in the Variables window. Editing stack will not be crucial but could be useful if it won't be too time-consuming to implement it.
- Microprogram in the Control Unit window should be hidden in A-level mode, and only showed in the Full mode. The warning about complexity of the window should not be shown.
- The user should be able to turn off the "flying data windows" if they are not necessary.
- The Buses window should be extended to show an overview of the system unit, showing the CPU, RAM, buses and the three controllers video, keyboard and speaker and show the "flying windows" to display data flow between them.
- All animated examples should be available through a menu on the Main window. It should be possible to run them without having anything to do with the rest of the system. Apart from ALU operation animations already mentioned, the following animations should be added provided there is enough time to develop them:
 - o Differences between different addressing modes
 - Arrays uses, especially in loops
 - Sorting algorithms
 - Binary trees explained, binary trees and searching.
- Window captions should not contain any detailed information
- Each window should have a What's This button in the caption, and a short note should popup describing any window element should the user use this button.
- RAM window: It is not necessary to show what the instruction would be if PC was to point at the selected cell; rather, the bytes being executed should be highlighted in some way. Also, all the code generated by the compiler should be highlighted. Interrupt vector table should be highlighted with a special color. Stack values should be highlighted. All non-empty memory cells should be highlighted. All other memory cells should be dimmed. Block operations such as copy/paste are not required.
- Whenever the user is not running a program, all debug windows should clearly indicate that so that the user does not accidentally try to use them.

20. Testing strategy

Fully testing a system as big as this is extremely time-consuming. Therefore having a testing strategy is vital.

In this document, I will only discuss alpha testing. Beta testing takes a lot of time and requires many people to use the program for a while. Unfortunately, no time is available for beta testing. Note that I *will* let my end-user use the system and get some feedback from him – that will be discussed in Appraisal.

Alpha testing will be mostly black box testing. I have two reasons for choosing black box as opposed to white box testing. One is that preparing for white box testing is a lot more time consuming. The other is that I tested the procedures while implementing them, trying out every possibility, so doing white box testing again may not be very efficient.

Alpha testing will be split into the following parts:

- Testing assembly language this will involve writing different instructions with different opcodes etc. and making sure they are executed correctly.
- Testing windows this will be going through all windows and make sure they function the way they are supposed to.
- Overall testing this will involve developing several programs entirely in CLab. The purpose of this is to make sure it is in fact possible to develop a program in CLab (which neither of the previous two tests can prove).

Implementation

21. Plan

Developing a system of this complexity is a serious and time consuming task. To minimise time losses in case something goes wrong the system will be developed incrementally, in steps, so that at the end of each step the system can be run to see the results. An alternative way of developing it would be to write everything from the beginning to the end and only then run it for the first time. The advantage of the latter approach is that no time is spent on making an intermediary state work, but the disadvantage is that in case of a major design flaw a lot of code will have to be changed.

The outline of the implementation process is laid down below.

- Main form with stubs on most events; most frequent procedures in pUtils and pWinAPI; main types and variables in pGlobals.
- RAM window with basic functionality; pSynHigh unit with the required procedures doing everything in a single color
- CPU window with basic functionality; pExec executing simple instructions, do not show any data flow yet.
- pExec executes most instructions (except those that need anything not yet implemented)
- Code window with open/save facilities; program compilation
- Implement hardware windows; hardware interacts with code
- Interface windows
- OS/Debug windows
- Finish syntax highlighting etc.
- Any extra facilities if time left, such as a primitive Basic to assembly compiler

Note that at every stage some extra functionality may be implemented (this is especially true of pWinAPI, pUtils and pGlobals); the further the plans go the harder it becomes to predict precisely what will seem reasonable to develop next. This is why the further into the implementation the more general the points become.

At this point the actual coding begins; the next section will list the code after everything will have been written.

22. Listings

22.1. pGlobals

Option Explicit
' SUBSTRUCTURES' ' SUBSTRUCTURES'
'Video system state structure Public Type TpVideo
Mode As Integer 'Mode number character autoUpdate As Boolean 'Whether screens are refreshed automatically
mdResX As Integer 'number of characters/pixels horizontally mdResY As Integer 'number of characters/pixels vertically mdType As Integer '0=text, 1=graphics direct, 2=graphics paletted
mdColors As Integer'0=monochrome, 1=16, 2=256, 3=65536,4=16777216'one char width in screen pixelsmdFntX As Integer'one char width in screen pixels'one char height in screen pixels
MemOff As Long 'Offset to video memory in RAM PalMem(O To 255) As Long 'Video controller palette memory vDC As VirtualDC 'Virtual DC - bitmap End Type
' APP'
Private Type TApp 'Is set to true while unloading forms when app shuts down Termingting Ag Replan
Terminating As Boolean 'Previous window procedure pointer for fiMain PrevWndProc As Long
'OS version RunningOnWinXP As Boolean End Type Bublic Approved Type
Public Appp As TApp
' PROJ ' '' Private Type TProj
Modified As Boolean 'Variable for confirm queries
' Global settings' Complexity As Integer '0-basic, 1-alv1, 2-full NmbRep As Integer '0hex 1bin 2decU 3decS
' Program' P As TpPrg
' Execution'
Running As Boolean Paused As Boolean
Halted As Boolean TickCount As Long
CPU As TpCPU 'fhCPU, fhCU, fhALU RAM() As Byte 'fhRAM
Video As TpVideo 'fhVideo End Type
Public Proj As TProj
Sub Main()
' DESCRIPTION: Defines application entry point '
<pre>NOTES: all this is rather unusual and unnatural in VB but it gives me much more control and it seems to work fine.</pre>
Sub Main()
'We have started Appp.Terminating = False
'XP controls if running on XP, no harm otherwise InitCommonControls

	r.dwOSVersionInfoSize = 148 'according to Delphi's SizeOf Microsoft, what would I do without Delphi? Buy VC?
	GetVersionEx(osver)
	<pre>osver.dwMajorVersion > 4) And (osver.dwMinorVersion > 0) ' unningOnWinXP = True Else Appp.RunningOnWinXP = False</pre>
F F • • • •	
	play startup form
	lash.Show lash.Refresh
-1	
	d all forms but keep them invisible yboard.Hide
	eaker.Hide
	deo.Hide
	J.Hide .Hide
	M.Hide
	mp.Hide
	splay.Hide
	yboard.Hide in.Hide
fsCo	de.Hide
	gs.Hide
	ack.Hide rs.Hide
	tialise Proj
	.Modified = False .Complexity = 0
	NmbRep = 0
	Running = False
	.Paused = <mark>False</mark> tialize Proj.P
	n Proj.P.Ref(-1 To -1)
	n Proj.P.TknLine(-1 To -1)
	<pre>m Proj.P.Backpatch(-1 To -1) m Proj.P.ErrL.lError(-1 To -1)</pre>
	n Proj.P.ErrL.lWarning(-1 To -1)
	n Proj.P.ErrL.nError(-1 To -1)
	<pre>m Proj.P.ErrL.nWarning(-1 To -1) m Proj.P.ErrL.sError(-1 To -1)</pre>
	n Proj.P.ErrL.sWarning(-1 To -1)
	n Proj.P.Code_O2L(-1 To -1)
	<pre>m Proj.P.Code_L2O(-1 To -1) m Proj.P.Vars(-1 To -1)</pre>
	.P.Code = ""
	.P.CompileNeeded = True
	complexity J.SetComplexity
	tialize p* modules
	pCompile.cmpInit pExec.exeInit
	tialise hardware modules
	fhCPU.Reset fhCU.Init
	fhRAM.Init
	tialise devices
	devInit
	tialize fs* and fi* modules fsRegs.Init
Call	fiMain.Init
	fsCode.Init
	fsStack.Init fsVars.Init
	w main window
	in.Show w computer
	mp.Show
'Dest	troy splash form
Unloa	ad fiSplash
' Th	is is the end of the Main procedure, but '
	the application will keep running until '



22.2. pWinAPI

Option Explicit

'---- GDI ----'

--- Text ----

Public Declare Function SelectObject Lib "gdi32" (ByVal hdc As Long, ByVal hObject As Long) As Long

Public Declare Function DeleteObject Lib "gdi32" (ByVal hObject As Long) As Long

Public Declare Function GetStockObject Lib "gdi32" (ByVal nIndex As Long) As Long

Public Declare Function SetBkMode Lib "gdi32" (ByVal hdc As Long, ByVal nBkMode As Long) As Long

Public Declare Function SetTextColor Lib "gdi32" (ByVal hdc As Long, ByVal crColor As Long) As Long

Public Declare Function CreateFont Lib "gdi32" Alias "CreateFontA" (ByVal h As Long, ByVal W As Long, ByVal E As Long, ByVal o As Long, ByVal W As Long, ByVal i As Long, ByVal u As Long, ByVal s As Long, ByVal c As Long, ByVal OP As Long, ByVal CP As Long, ByVal Q As Long, ByVal PAF As Long, ByVal f As String) As Long

Public Declare Function CreateBrushIndirect Lib "gdi32" (lpLogBrush As LOGBRUSH) As Long

Public Type LOGBRUSH

lbStyle As Long lbColor As Long lbHatch As Long

End Type

Public Declare Function GetTextExtentPoint32 Lib "gdi32" Alias "GetTextExtentPoint32A" (ByVal hdc As Long, ByVal lpsz As String, ByVal cbString As Long, lpSize As Size) As Long Public Declare Function TextOut Lib "gdi32" Alias "TextOutA" (ByVal hdc As Long, ByVal x As Long, ByVal y As Long, ByVal lpString As String, ByVal nCount As Long) As Long Public Declare Function DrawText Lib "user32" Alias "DrawTextA"

Public Declare Function DrawText Lib "user32" Alias "DrawTextA" (ByVal hdc As Long, ByVal lpStr As String, ByVal nCount As Long, lpRect As RECT, ByVal wFormat As Long) As Long

'--- Pictures ---'

Public Declare Function BitBlt Lib "gdi32" (ByVal hDestDC As Long, ByVal x As Long, ByVal y As Long, ByVal nWidth As Long, ByVal nHeight As Long, ByVal hSrcDC As Long, ByVal XSrc As Long, ByVal YSrc As Long, ByVal dwRop As Long) As Long Public Declare Function StretchBlt Lib "gdi32" (ByVal hdc As Long,

ByVal x As Long, ByVal y As Long, ByVal nWidth As Long, ByVal nHeight As Long, ByVal hSrcDC As Long, ByVal XSrc As Long, ByVal YSrc As Long, ByVal nSrcWidth As Long, ByVal nSrcHeight As Long, ByVal dwRop As Long) As Long

Public Declare Function MoveToEx Lib "gdi32" (ByVal hdc As Long, ByVal x As Long, ByVal y As Long, lpPoint As POINTAPI) As Long Public Declare Function LineTo Lib "gdi32" (ByVal hdc As Long, ByVal x As Long, ByVal y As Long) As Long

Public Declare Function Rectangle Lib "gdi32" (ByVal hdc As Long, ByVal X1 As Long, ByVal Y1 As Long, ByVal X2 As Long, ByVal Y2 As Long) As Long

 $\label{eq:public beclare Function SetPixelV Lib "gdi32" (ByVal hdc As Long, ByVal x As Long, ByVal y As Long, ByVal crColor As Long) As Long (ByVal crColor As L$

Public Declare Function InvalidateRect Lib "user32" (ByVal hwnd As Long, lpRect As RECT, ByVal bErase As Long) As Long Public Declare Function GetClientRect Lib "user32" (ByVal hwnd As

Long, lpRect As RECT) As Long Public Const TRANSPARENT = 1

Public Type POINTAPI

NOTES:
I wish someone knew how much I hate VB. This is one of hundreds of other things which cause it.
VB does not allow to use AddressOf on any procs
which are declared in a Form module, so even a
function so closely related to the form (much closer than ANY other function) cannot be in the same module. It is very clumsy having it here.

Function WindowProc(ByVal hw As Long, ByVal uMsg As Long, ByVal
wParam As Long, ByVal lParam As Long) As Long
WindowProc = fiMain.WindowProc(hw, uMsg, wParam, lParam)
End Function



End Type Public Const SRCCOPY = &HCC0020 Public Type Size cx As Long cy As Long End Type

System colors

Public Declare Function GetSysColor Lib "user32" (ByVal nIndex As Long) As Long Public Const COLOR_ACTIVEBORDER = 10 Public Const COLOR_ACTIVECAPTION = 2 Public Const COLOR_APPWORKSPACE = 12 Public Const COLOR_BACKGROUND = 1 Public Const COLOR_BINFACE = 15 Public Const COLOR_BINFIGHLIGHT = 20

Public	Const	COLOR_ADJ_MAX = 100	
Public	Const	$COLOR_ADJ_MIN = -100$	
Public	Const	$COLOR_BTNSHADOW = 16$	
Public	Const	COLOR_BTNTEXT = 18	
Public	Const	COLOR_CAPTIONTEXT = 9	
Public	Const	COLOR_GRAYTEXT = 17	
Public	Const	COLOR_HIGHLIGHT = 13	
Public	Const	COLOR_HIGHLIGHTTEXT = 14	
Public	Const	COLOR_INACTIVEBORDER = 11	
Public	Const	$COLOR_INACTIVECAPTION = 3$	
Public	Const	COLOR_INACTIVECAPTIONTEXT = 19	
Public	Const	$COLOR_MENU = 4$	
Public	Const	COLOR_MENUTEXT = 7	
Public	Const	COLOR_SCROLLBAR = 0	
Public	Const	COLOR_WINDOW = 5	
Public	Const	$COLOR_WINDOWFRAME = 6$	
Public	Const	COLOR_WINDOWTEXT = 8	

Window procedures & messaging

Public Declare Function CallWindowProc Lib "user32" Alias "CallWindowProcA" (ByVal lpPrevWndFunc As Long, ByVal hwnd As Long, ByVal Msg As Long, ByVal wParam As Long, ByVal lParam As Long) As Long Public Declare Function SetWindowLong Lib "user32" Alias "SetWindowLongA" (ByVal hwnd As Long, ByVal nIndex As Long, ByVal dwNewLong As Long) As Long Public Declare Function DefWindowProc Lib "user32" Alias "DefWindowProcA" (ByVal hwnd As Long, ByVal wMsg As Long, ByVal wParam As Long, ByVal lParam As Long) As Long

Public Const GWL_WNDPROC = -4

Public Const WM_SYSCOMMAND = &H112 Public Const SC_MINIMIZE = &HF020& Public Const SC_RESTORE = &HF120& Public Const WM_NCLBUTTONDOWN = &HA1 Public Const WM_SIZING = 532 Public Const WM_SIZING = 532 Public Const WM_SIZE = &H5 Public Type RECT Left As Long Top As Long Right As Long



22.3. pUtils

Option Explicit

	' error icon and an OK button. '
Public declarations in this module:	
Public declarations in this module:	PARAMETERS:
PROCEDURES:	' msg - the message to be displayed. '
FROCEDORES.	Public Sub Errr(Msg As String)
Tally	Call MsgBox(Msg, vb0KOnly Or vbExclamation, "Error")
FieldStr	End Sub
InStrBack	
Hex2Dec	tt
Dec2Hex	' Public Function Tally(where As String, what As '
Bin2Dec '	' String)
Dec2Chr	· · ·
Chr2Dec	' DESCRIPTION: Counts the number of occurences of '
Str2Chr	' What in Where. What should be one character '
TestCharset	' long only, or the function will return 0 '
StringIsInt '	''
StringIsLong	Public Function Tally(where As String, what As String)
·(Dim i As Integer, A As Integer
	A = 0
Font for easier custom-drawn text	For i = 1 To Len(where)
Public Type TFnt	If Mid(where, i, 1) = what Then $A = A + 1$
'Parameters	Next
ForeColor As Long	Tally = A
BackColor As Long	End Function
FaceName As String	
Size As Long	۱۱
Weight As Long	' Public Function FieldStr(row As String, index '
'Associated GDI objects - call CreateFnt to init	' As Integer, separator As String)
fntFont As Long	1
fntBrush As Long	' DESCRIPTION: Returns element number Index from '
'Internal params	' Row in which elements are separated by char '
fntCreated As Boolean 'to free it safely	' Separator. Separator should be 1 char long! '
and Type	··
	Public Function FieldStr(row As String, Index As Integer,
<u>(</u>	separator As String)
Public Sub Errr(msg As String)	Dim i As Integer
	Dim s As String

Displays an error messsage box with a red


Function TestCharset(testwhat As String, charset As String) As Boolean Dim i As Integer For i = 1 To Len(testwhat) If InStr(charset, Mid(testwhat, i, 1)) = 0 Then TestCharset = False GoTo tested End If Next TestCharset = True End Function Function StringIsInt(s As String) As Boolean DESCRIPTION: Returns true if string can be Function StringIsInt(s As String) As Boolean On Error GoTo strNotInt Dim i As Integer i = CInt(s) StringIsInt = True Exit Function StringIsInt = False End Function Function StringIsLong(s As String) As Boolean DESCRIPTION: Returns true if string can be converted to type Long. Function StringIsLong(s As String) As Boolean On Error GoTo strNotLong Dim i As Long i = CLng(s)StringIsLong = True Exit Function StringIsLong = False End Function Function GetFilename(Save As Boolean, ByRef FileName As String, _ InitDir As String, Filter As String, DefExt As String, Title As String) As Boolean 'Fill open structure Dim o As OPENFILENAME o.lStructSize = 88 'Delphi's SizeOf(TOpenFilen. o.hInstance = 0o.lpstrFilter = Replace(Filter, "|", Chr(0)) + Chr(0) + Chr(0) o.nFilterIndex = 0 o.nMaxFile = 260o.lpstrFile = FileName + String(262 - Len(FileName), Chr(0)) o.lpstrInitialDir = InitDir o.lpstrDefExt = DefExt o.lpfnHook = 0& o.hwndOwner = 0 If Save Then o.FLAGS = OFN_PATHMUSTEXIST + OFN_OVERWRITEPROMPT Else o.FLAGS = OFN_FILEMUSTEXIST End If 'Return result If Save Then GetFilename = (GetSaveFileName(o) <> 0) Else GetFilename = (GetOpenFileName(o) <> 0) End If 'Return file nam If InStr(0.lpstrFile, Chr(0)) > 0 Then o.lpstrFile = Left(o.lpstrFile, InStr(o.lpstrFile, Chr(0)) - 1) FileName = 0.lpstrFile End Function Public Function AppDir() As String AppDir = App.Path + IIf(Right(App.Path, 1) = "\", "", "\") End Function Public Function Dec2Fmt16(num As Long, fmt As Integer) As String If fmt = 0 Then 'hex Dec2Fmt16 = Dec2Hex(num, 4) + "h" ElseIf fmt = 1 Then 'bin

Dec2Fmt16 = "binary not supported yet"
ElseIf fmt = 2 Then 'decU Dec2Fmt16 = CStr(num) ElseIf fmt = 3 Then 'decs Dec2Fmt16 = CStr(IIf(num >= 32768, -65536 + num, num)) Else Dec2Fmt16 = "Invalid format number" End If End Function Public Function IsFmt16(num As String) As Boolean Public Function IsFmt16(num As String) As Boolean On Error GoTo IsNot Dim s As String, testval As Long, minus As Boolean If Len(num) = 0 Then GoTo IsNot s = UCase(num) minus = False If Left(s, 1) = "-" Then If Len(s) = 1 Then GoTo IsNot s = Mid(s, 2)minus = True End If If Right(s, 1) = "H" Or Right(s, 1) = "B" Then If Len(s) = 1Then GoTo IsNot nvert (overflow will be trapped) 'Check charset and try to If Right(s, 1) = "H" Then s = Left(s, Len(s) - 1)If Not TestCharset(s, "0123456789ABCDEF") Then GoTo IsNot testval = Hex2Dec(s) ElseIf Right(s, 1) = "B" Then
s = Left(s, Len(s) - 1) If Not TestCharset(s, "01") Then GoTo IsNot testval = Bin2Dec(s)Else If Not TestCharset(s, "0123456789") Then GoTo IsNot testval = CLng(s) End If 'Check range If minus Then testval = -testval If testval < -32768 Or testval > 65535 Then GoTo IsNot 'Everything is fine TsFmt16 = True Exit Function IsFmt16 = False End Function Public Function Fmt2Dec16(c As String) As Long trap and hope ked C before calling this On Error GoTo HoustonWeVeGotAProblem Dim s As String, minus As Boolean, n As Long s = UCase(c)minus = False If Left(s, 1) = "-" Then minus = True s = Mid(s, 2)End If If Right(s, 1) = "H" Then n = Hex2Dec(Left(s, Len(s) - 1))ElseIf Right(s. 1) = "B" Then n = Bin2Dec(Left(s, Len(s) - 1))Else n = CLng(s)End If 'Deal with minus sign If minus Then n = -n 'Return Fmt2Dec16 = Exit Function HoustonWeVeCo Fmt2Dec16 = -1 'should not happen unless IsFmt16 not called before End Function

Listings



22.4. pCompile

Option Explicit

Option Explicit	
	'Token structure - describes one token
··································	Private Type TpToken
' Public declarations in this module: '	Text As String 'eg MOV
1	Type As Integer 'eg Operand (see constants above)
' TYPES: '	End Type
' TPrg - assembly language program '	**
1 1	'Token line - array of tokens in one token line
' PROCEDURES:	Private Type TpTokenLine
' cmpInit - initialises this module '	Token() As TpToken
	CodeLine As Integer
' PrgCompile - compiles a program '	
PrgLoad - loads a program into RAM	CodeOffset As Integer
CreatePrg - inits TPrg structure	End Type
' VARIABLES:	'Structure to store backpatch requests
' Proj.P - program being worked on '	Private Type TpBackpatch
	Name As String 'which variable's address needed
' CONSTANTS: '	Addr As Long 'where to write the address
' CharsetLabel - charset for labels '	IsDW As Boolean 'true for absolute address
(<u></u>)	RelTo As Long 'to calculate relative address
	CodeLine As Integer where requested
۱۱	End Type
' Local declarations'	11
''	'Stores all references & their addresses
	Private Type TpRef
'Token type constants	Name As String
Private Const tkUnknown = -1 'Used in the process of	Addr As Long
tokenization	
Private Const tkLabel = 0	
	End Type
Private Const tkVarDec1 = 3	
Private Const tkVarInit = 4	'Error log
Private Const tkOpcode = 5	Private Type TpErrLog
Private Const tkOperand = 6	sError() As String 'Message
	lError() As Integer 'Line number
'Token descriptions	nError() As String 'Error number
Private tkName(-1 To 6) As String	sWarning() As String 'Message
	lWarning() As Integer 'Line number
'Opcode-to-machinecode reference structure	nWarning() As String 'Warning number
'Arrays of these structures are grouped by similarity in	End Type
compilation	
Private Type TOpMCode	'Data about all declared variables
Opcode As String	Private Type TpVars
Code As Byte	Name As String
Code2 As Byte	Addr As Long
End Type	End Type
'List of type 0 opcodes (all operandless ones)	'Assembly language program
Private cdType0() As TOpMCode	Public Type TpPrg
'List op type 1 opcodes (bitwise shifts)	AsmLine() As String
Private cdType1() As TOpMCode	TknLine() As TpTokenLine
'List op type 2 opcodes (4-way arithmetic & bitwise)	Code As String
Private cdType2() As TOpMCode	Code_02L() As Integer 'convert offset to source code line
'List op type 3 opcodes (inc,dec,neg,not,bswp)	Code_L2O() As Integer 'convert source code line to offset
Private cdType3() As TOpMCode	
'List op type 4 opcodes (jmp, jxx, call)	Ref() As TpRef
Private cdType4() As TOpMCode	Backpatch() As TpBackpatch
'Charset for label names	Vars() As TpVars 'Stores info about db/dw/ds
Public CharsetLabel As String	ErrL As TpErrLog

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CompileNeeded As Boolean 'True if Compile called after editing
End Type
'' ' Public Sub cmpInit() '
Initializes this module
Public Sub cmpInit()
'Initialise charsets
CharsetLabel =
"abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ1234567890" 'Init token names
tkName(tkUnknown) = "unknown token"
tkName(tkLabel) = "label"
tkName(tkVarDecl) = "variable declaration"
<pre>tkName(tkVarInit) = "variable initialisation" tkName(tkOpcode) = "opcode"</pre>
tkName(tkOperand) = "operand"
'Init opcode-to-machinecode array
ReDim cdType0(0 To 19) cdType0(0).Opcode = "pushpc" 'MUST BE LOWERCASE!
cdType0(0).Code = & H13
cdType0(1).Opcode = "pushsp"
cdType0(1).Code = &H14
cdType0(2).Opcode = "pushfl"
cdType0(2).Code = &H15 cdType0(3).Opcode = "popsp"
cdType0(3).Code = &H16
cdType0(4).Opcode = "popfl"
cdType0(4).Code = &H17
cdType0(5).Opcode = "sp2b" cdType0(5).Code = &HF
cdType0(6).Opcode = "stz"
cdType0(6).Code = &H54
cdType0(7).Opcode = "clz" cdType0(7).Code = &H55
cdType0(8).0pcode = "stc"
cdType0(8).Code = &H56
cdType0(9).Opcode = "clc"
cdType0(9).Code = &H57 cdType0(10).Opcode = "sto"
cdType0(10).Code = &H64
cdType0(11).Opcode = "clo"
cdType0(11).Code = &H65
cdType0(12).Opcode = "sts" cdType0(12).Code = &H66
cdType0(13).Opcode = "cls"
cdType0(13).Code = &H67
cdType0(14).Opcode = "sti" cdType0(14).Code = &H76
cdType0(14).Code = "cli"
cdType0(15).Code = &H77
cdType0(16).Opcode = "ret"
cdType0(16).Code = &H72 cdType0(17).Opcode = "iret"
cdType0(17).Code = &H73
cdType0(18).Opcode = "halt"
cdType0(18).Code = &H75
cdType0(19).Opcode = "nop" cdType0(19).Code = &H8F
' <u>Init</u> opcode- <u>to-machi</u> necode array
ReDim cdType1(0 To 7)
cdType1(0).Opcode = "lshl" cdType1(0).Code = &HCO
cdType1(1).Opcode = "lshr"
cdType1(1).Code = &HC1
cdType1(2).Opcode = "ashl"
<pre>cdType1(2).Code = &HC2 cdType1(3).Opcode = "ashr"</pre>
cdType1(3).Code = &HC3
cdType1(4).Opcode = "rol"
cdType1(4).Code = &HC4 cdType1(5).Opcode = "ror"
cdType1(5).Code = &HC5
cdType1(6).Opcode = "rcl"
cdType1(6).Code = &HC6
cdType1(7).Opcode = "rcr" cdType1(7).Code = &HC7
'Init opcode-to-machinecode array
ReDim cdType2(0 To 13)
cdType2(0).Opcode = "add"
cdType2(0).Code = &H80 cdType2(1).Opcode = "sub"
cdType2(1).Code = &H83
cdType2(2).Opcode = "adc"

cdType2(2).Code = &H86 cdType2(3).Opcode = "sbb"
cdType2(3).Opcode = SDDcdType2(3).Code = &H89
cdType2(4).Opcode = "cmp"
cdType2(4).Code = &H8C
cdType2(5).Opcode = "mul"
cdType2(5).Code = &H90
cdType2(6).Opcode = "div"
cdType2(6).Code = &H93
cdType2(7).Opcode = "imul"
cdType2(7).Code = &H96
cdType2(8).Opcode = "idiv"
cdType2(8).Code = &H99
cdType2(9).Opcode = "mod"
cdType2(9).Code = &H9C
cdType2(10).Opcode = "and"
cdType2(10).Code = &HB0
cdType2(11).Opcode = "or"
cdType2(11).Code = &HB3
cdType2(12).Opcode = "xor"
cdType2(12).Code = &HB6
cdType2(13).Opcode = "test"
cdType2(13).Code = &HB9
<u>'Init</u> opcode-to-machinecode array ReDim cdType3(0 To 4)
cdType3(0).Opcode = "inc"
cdType3(0).Code = &HA0
cdType3(0).Code2 = &HAC
cdType3(1).Opcode = "dec"
cdType3(1).Code = &HA4
cdType3(1).Code2 = &HAD
cdType3(2).Opcode = "neg"
cdType3(2).Code = &HA8
cdType3(2).Code2 = &HAE
cdType3(3).Opcode = "not"
cdType3(3).Code = &HBC
cdType3(3). $Code2 = &HAF$
cdType3(4).0pcode = "bswp"
cdType3(4).Code = &HDC
cdTwpe3(4) $Code2 = &H9F$
cdType3(4).Code2 = &H9F 'Init opcode-to-machinecode array
'Init opcode-to-machinecode array
'Init opcode-to-machinecode array ReDim cdType4(0 To 13)
'Init opcode-to-machinecode array
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Code = %H70 cdType4(1).Opcode = "jg"</pre>
<pre>'Init_opcode-to-machinecode array ReDim_cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Code = &H70 cdType4(1).Opcode = "jg" cdType4(1).Code = &H40</pre>
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Code = &H70 cdType4(1).Opcode = "jg" cdType4(1).Code = &H40 cdType4(2).Opcode = "j1"</pre>
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Code = &HT0 cdType4(1).Opcode = "jg" cdType4(1).Code = &H40 cdType4(2).Opcode = "j1" cdType4(2).Code = &H42</pre>
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<pre>'Init_opcode-to-machinecode array ReDim_cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Code = &H70 cdType4(1).Opcode = "jg" cdType4(1).Code = &H40 cdType4(2).Opcode = "j1" cdType4(2).Opcode = "j1" cdType4(3).Opcode = "jge" cdType4(3).Code = &H43 cdType4(4).Opcode = "jle"</pre>
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<pre>'Init_opcode-to-machinecode array ReDim_cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Opcode = %H70 cdType4(1).Opcode = %H70 cdType4(1).Opcode = "jg" cdType4(1).Opcode = "j1" cdType4(2).Opcode = "j1" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(4).Opcode = "j1" cdType4(5).Opcode = "j1" cdType4(5).Opcode = "j1" cdType4(5).Opcode = "j1" cdType4(5).Opcode = "j1" cdType4(6).Opcode = "j1" cdType4(6).Opcode = "j1" cdType4(6).Opcode = "j1" cdType4(7).Opcode = "j1" cdType4(7).Opcode = "j1" cdType4(8).Opcode = "j0" cdType4(8).Opcode = "j0" cdType4(9).Opcode = "j0" cdType4(9).Opcode = "j0"</pre>
<pre>'Init_opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(1).Opcode = "jm" cdType4(1).Opcode = "jg" cdType4(1).Code = &H40 cdType4(2).Opcode = "j1" cdType4(2).Code = &H42 cdType4(3).Opcode = "jge" cdType4(3).Opcode = "jge" cdType4(3).Code = &H43 cdType4(4).Opcode = "jle" cdType4(5).Opcode = "j1" cdType4(5).Opcode = "j1" cdType4(5).Code = &H50 cdType4(5).Code = &H50 cdType4(6).Opcode = "jc" cdType4(6).Opcode = "jc" cdType4(6).Code = &H51 cdType4(7).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(10).Code = &H60 cdType4(10).Code = &H61</pre>
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<pre>'Init_opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(1).Opcode = "jm" cdType4(1).Opcode = "jg" cdType4(1).Code = &H40 cdType4(2).Opcode = "j1" cdType4(2).Code = &H42 cdType4(3).Opcode = "jge" cdType4(3).Opcode = "jge" cdType4(3).Code = &H43 cdType4(3).Opcode = "jle" cdType4(4).Code = &H41 cdType4(5).Opcode = "j1" cdType4(5).Opcode = "j1" cdType4(5).Opcode = "j1" cdType4(6).Code = &H50 cdType4(6).Code = &H51 cdType4(6).Code = &H51 cdType4(7).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(9).Code = &H53 cdType4(9).Opcode = "jo" cdType4(10).Opcode = "jo" cdType4(11).Opcode = "js" cdType4(11).Opcode = "js" cdType4(12).Opcode = "jns" cdType4(12).Code = &H63</pre>
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<pre>'Init_opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Opcode = %H70 cdType4(1).Opcode = %H70 cdType4(1).Opcode = %jg" cdType4(2).Opcode = "jg" cdType4(2).Opcode = "jl" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(4).Opcode = "je" cdType4(5).Opcode = "je" cdType4(5).Opcode = "je" cdType4(5).Opcode = "je" cdType4(6).Opcode = "je" cdType4(6).Opcode = "je" cdType4(6).Opcode = "je" cdType4(7).Opcode = "je" cdType4(8).Opcode = "je" cdType4(8).Opcode = "jo" cdType4(8).Opcode = "jo" cdType4(8).Opcode = "jo" cdType4(9).Opcode = "jo" cdType4(9).Opcode = "jo" cdType4(9).Opcode = "jo" cdType4(10).Opcode = "js" cdType4(11).Opcode = "js" cdType4(11).Opcode = "js" cdType4(12).Opcode = "jns" cdType4(13).Opcode = "call" cdType4(13).Code = &H71 End Sub</pre>
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Opcode = "jmp" cdType4(1).Opcode = "jg" cdType4(1).Opcode = "jg" cdType4(2).Opcode = "j1" cdType4(2).Opcode = "j1" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(4).Opcode = "je" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(7).Opcode = "jc" cdType4(7).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(9).Opcode = "jr" cdType4(10).Opcode = "js" cdType4(11).Opcode = "js" cdType4(12).Opcode = "js" cdType4(12).Opcode = "call" cdType4(13).Opcode = %H71 End Sub</pre>
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Opcode = %H70 cdType4(1).Opcode = %H70 cdType4(1).Opcode = %J9" cdType4(2).Opcode = "jg" cdType4(2).Opcode = "j1" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(4).Opcode = "je" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(7).Opcode = "jr" cdType4(8).Opcode = "jr" cdType4(8).Opcode = "jr" cdType4(8).Opcode = "jr" cdType4(8).Opcode = "jr" cdType4(9).Opcode = "jr" cdType4(9).Opcode = "jr" cdType4(9).Opcode = "jr" cdType4(10).Opcode = "js" cdType4(11).Opcode = "js" cdType4(11).Opcode = "js" cdType4(11).Opcode = "js" cdType4(112.Opcode = "call" cdType4(13).Opcode = "call" cdType4(13).Opcode = %Call" cdType4(13).Opcode = %Call" cdType4(14).Opcode = %Call" cdType4(15).Opcode = %Call" cdType4(15).Opcode</pre>
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(1).Opcode = "jmp" cdType4(1).Opcode = "jg" cdType4(1).Code = &H40 cdType4(2).Code = &H42 cdType4(2).Code = &H42 cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(3).Code = &H43 cdType4(4).Code = &H41 cdType4(5).Opcode = "je" cdType4(5).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(6).Code = &H50 cdType4(6).Code = &H51 cdType4(6).Opcode = "jc" cdType4(7).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jr" cdType4(9).Code = &H52 cdType4(9).Opcode = "jo" cdType4(9).Opcode = "jo" cdType4(9).Opcode = "jo" cdType4(10).Opcode = "js" cdType4(11).Opcode = "js" cdType4(11).Opcode = "js" cdType4(12).Opcode = "js" cdType4(13).Opcode = "call" cdType4(13).Opcode = "call" cdType4(13).Code = &H71 End Sub</pre>
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Opcode = %H70 cdType4(1).Opcode = %H70 cdType4(1).Opcode = %J9" cdType4(2).Opcode = "jg" cdType4(2).Opcode = "j1" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(4).Opcode = "je" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(6).Opcode = "jr" cdType4(7).Opcode = "jr" cdType4(8).Opcode = "jr" cdType4(8).Opcode = "jr" cdType4(8).Opcode = "jr" cdType4(8).Opcode = "jr" cdType4(9).Opcode = "jr" cdType4(9).Opcode = "jr" cdType4(9).Opcode = "jr" cdType4(10).Opcode = "js" cdType4(11).Opcode = "js" cdType4(11).Opcode = "js" cdType4(11).Opcode = "js" cdType4(112.Opcode = "call" cdType4(13).Opcode = "call" cdType4(13).Opcode = %Call" cdType4(13).Opcode = %Call" cdType4(14).Opcode = %Call" cdType4(15).Opcode = %Call" cdType4(15).Opcode</pre>
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Ocode = "jmp" cdType4(0).Code = &H70 cdType4(1).Ocode = &H70 cdType4(1).Ocode = %H70 cdType4(2).Ocode = %J9" cdType4(2).Ocode = %J1" cdType4(2).Code = &H42 cdType4(3).Ocode = "je" cdType4(3).Ocode = "je" cdType4(4).Code = &H43 cdType4(5).Ocode = "je" cdType4(5).Ocode = "je" cdType4(5).Ocode = "je" cdType4(5).Ocode = "je" cdType4(6).Ocode = #H51 cdType4(6).Ocode = "jc" cdType4(6).Ocode = "jc" cdType4(7).Ocode = "jc" cdType4(8).Ocode = "jc" cdType4(8).Ocode = "jo" cdType4(8).Ocode = "jo" cdType4(8).Ocode = "jo" cdType4(9).Ocode = "jo" cdType4(9).Ocode = "jo" cdType4(9).Ocode = "jo" cdType4(10).Ocode = "jo" cdType4(10).Ocode = "js" cdType4(11).Ocode = %H61 cdType4(11).Ocode = %H62 cdType4(12).Ocode = "jns" cdType4(13).Ocode = "call" cdType4(13).Ocode = %code ' Private Sub ReadCodeIntoProj() 'Initialise program</pre>
<pre>'Init opcode-to-machinecode array ReDim cdType4(0 To 13) cdType4(0).Opcode = "jmp" cdType4(0).Code = & HTO cdType4(1).Opcode = "jg" cdType4(1).Opcode = "jg" cdType4(2).Opcode = %J1" cdType4(2).Opcode = %J1" cdType4(2).Code = & H42 cdType4(3).Opcode = "jge" cdType4(3).Opcode = "je" cdType4(3).Opcode = "je" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(5).Opcode = "jr" cdType4(6).Code = & H51 cdType4(7).Opcode = "jc" cdType4(7).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(8).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(9).Opcode = "jc" cdType4(9).Opcode = "jn" cdType4(10).Opcode = "jn" cdType4(10).Opcode = "jns" cdType4(11).Opcode = "js" cdType4(12).Opcode = "js" cdType4(12).Opcode = "js" cdType4(13).Opcode = "call" cdType4(13).Opcode = & H63 cdType4(13).Opcode = & H71 End Sub</pre>

KeDIm Proj.P.AsmLine(-1 TO -1)
'Copy source text
Dim i As Integer
ReDim Proj.P.AsmLine(-1 TO fsCode.RTB.Lines.Count - 1)
For i = 0 To UBound(Proj.P.AsmLine)
Proj.P.AsmLine(i) = fsCode.RTB.Lines.Item(i)

'If Right(Proj.P.AsmLine(i), 1) = Chr(13) Then Proj.P.AsmLine(i) = Mid(Proj.P.AsmLine(i), 1,	'Copy program to work on it (and do some processing in the
Len(Proj.P.AsmLine(i)) - 1) Next End Sub	meantime) Dim sc_line() As Integer ReDim sc_line(-1 To UBound(.AsmLine))
······································	<pre>'nl = 0 For i = 0 To UBound(.AsmLine)</pre>
' Public Sub PrgCompile()	Copy s = .AsmLine(i)
' Compiles program in Proj, returning ' ' machine code in p.Code '	<pre>"Remove comment A = InStr(s, ";")</pre>
Public Sub PrgCompile()	If A > 0 Then s = Left(s, A - 1) 'Clean spaces
'Load program into Proj ReadCodeIntoProj	s = CleanSpaces(s)
'Empty everything but AsmLine ReDim Proj.P.Ref(-1 To -1)	s = Trim(s) 'Save string
ReDim Proj. P. TknLine(-1 To -1) ReDim Proj. P. Backpatch(-1 To -1)	.AsmLine(i) = s 'Remember line number
ReDim Proj.P.ErrL.lError(-1 To -1) ReDim Proj.P.ErrL.lWarning(-1 To -1) ReDim Proj.P.ErrL.nError(-1 To -1)	<pre>sc_line(i) = i 'Increase nl to point to next line 'nl = nl + 1</pre>
ReDim Proj.P. ErrL. nWarning(-1 To -1) ReDim Proj.P. ErrL. sError(-1 To -1)	Next
ReDim Proj.P. ErrL. SWarning(-1 To -1) ReDim Proj.P. Code_02L(-1 To -1)	'' ' Pure tokenization'
ReDim Proj. P. Code_L20(-1 To -1) ReDim Proj. P. Vars(-1 To -1)	· · · · · · · · · · · · · · · · · · ·
Proj. P. Code = "" Proj. P. CompileNeeded = True	'Prepare Tokenized ReDim .TknLine(-1 To UBound(.AsmLine))
Compile Proj.P.CompileNeeded = False	'Tokenize
Call CompilePass1 If UBound(Proj.P.ErrL.sError) = -1 Then Call CompilePass2	nl = 0 Dim sptr As Long, tmps As String, bs As Integer
If UBound(Proj.P.ErrL.sError) = -1 Then Call CompilePass3 'Display errors	<pre>For i = 0 To UBound(.AsmLine) If .AsmLine(i) <> "" Then</pre>
Dim i As Integer fsCode.LErr.Clear	'Prep token line ReDim .TknLine(nl).Token(-1 To -1)
<pre>For i = 0 To UBound(Proj.P.ErrL.SError) Call fsCode.LErr.AddItem("Error (" +</pre>	.TknLine(nl).CodeLine = sc_line(i) 'Get the tokens
CStr(Proj.P.ErrL.lError(i) + 1) + "): " + Proj.P.ErrL.sError(i) + "(" + Proj.P.ErrL.nError(i) + ").")	<pre>sptr = 1 tmps = .AsmLine(i)</pre>
Next 'Display warnings Display warnings	Do While sptr <= Len(tmps) <u>Skip all spaces</u> Do Wid(tmps ant 1)
<pre>For i = 0 To UBound(Proj.P.ErrL.sWarning) Call fsCode.LErr.AddItem("Warning (" + Cstr(Proj.P.ErrL.lWarning(i) + 1) + "): " +</pre>	<pre>Do While (Mid(tmps, sptr, 1) = " " Or Mid(tmps, sptr, 1) = Chr(9)) And sptr <= Len(tmps) sptr = sptr + 1</pre>
Proj.P.ErrL.sWarning(i) + " (" + Proj.P.ErrL.nWarning(i) + ").") Next	Loop 'Prep token
fiMain.UpdateAll True	ReDim Preserve .TknLine(nl).Token(-1 To UBound(.TknLine(nl).Token) + 1)
End Sub	.TknLine(nl).Token(UBound(.TknLine(nl).Token)).Text = "" .TknLine(nl).Token(UBound(.TknLine(nl).Token)).Type =
'' ' Public Sub PrgLoad '	tkUnknown 'Is_it_a_string?
' ' ' Loads the program in Proj into '	<pre>If Mid(tmps, sptr, 1) = """" Then .TknLine(n1).Token(UBound(.TknLine(n1).Token)).Text =</pre>
'RAM at offset zero. '	sptr = sptr + 1
Public Sub PrgLoad() Dim i As Integer	<pre>If sptr > Len(tmps) Then GoTo dneS 'I know it's bad programming but bs = 0</pre>
<pre>For i = 1 To Len(Proj.P.Code) Proj.RAM(i - 1) = Asc(Mid(Proj.P.Code, i, 1)) Next</pre>	Do Do If Mid(tmps, sptr, 1) = """" Then
fhRAM.Update End Sub	bs = bs + 1 .TknLine(n1).Token(UBound(.TknLine(n1).Token)).Text
· · · · · · · · · · · · · · · · · · ·	<pre>= .TknLine(nl).Token(UBound(.TknLine(nl).Token)).Text + """" Else</pre>
' Private Sub CompilePass1() '	bs = 0 .TknLine(nl).Token(UBound(.TknLine(nl).Token)).Text
' DESCRIPTION: Compilation pass 1: tokenize program ' ' 1. Clean the source code '	<pre>= .TknLine(nl).Token(UBound(.TknLine(nl).Token)).Text + Mid(tmps, sptr, 1)</pre>
' 2. Split every line into tokens ' ' 3. Determine token types '	End If sptr = sptr + 1
' 4. Token pattern analysis ' ' 5. Preparations for variable compilation '	Loop Until ((bs \ 2 <> bs / 2) And (Mid(tmps, sptr, 1) = """")) Or (sptr > Len(tmps))
' OUTPUT: Proj.P.Tkn* containing tokenized program. '	<pre>dneS: .TknLine(nl).Token(UBound(.TknLine(nl).Token)).Text =</pre>
Proj will be ready for CompilePass2	.TknLine(nl).Token(UBound(.TknLine(nl).Token)).Text + """" Else
Private Sub CompilePass1() Dim i As Integer, A As Integer, nl As Integer	'There is no need for any advanced token splitting 'in a language as simple as this. In an expression
Dim s As String With Proj.P	'like [b*c+4356h] we don't care that * and + should 'be separate tokens - we don't need that. The only 'reason to setue an alocation like this which has the
' Prepare to tokenize'	'reason to setup an algorithm like this which has the 'power to tokenize like that is to recoginze strings. 'Otherwise all we need is to separate tokens by the
' Prepare to tokenize'	'UtherWise all we need is to separate tokens by the 'fact that there is a whitespace inbetween them.



''	···
۲۲	'After this stage every OFFSET(varX) will be replaced with varX 'and every varX with [varX]
b = False If UBound(tkl(nl).Token) + 1 = 1 Then If tkl(nl) Taken(0) Theo = tkl(bbl (n tkl(nl) Taken(0) Theo	Dim acc As String, doing As Boolean For i = 0 To UBound(.TknLine)
<pre>If tkl(nl).Token(0).Type = tkLabel Or tkl(nl).Token(0).Type = tkOpcode Then</pre>	For A = 0 To UBound(.TknLine(i).Token)
Else	<pre>If .TknLine(i).Token(A).Type = tkOperand Then</pre>
<pre>If tkl(nl).Token(0).Type = tkVarDecl Then s = "a variable declaration without initialisation"</pre>	<pre>s = .TknLine(i).Token(A).Text 'Go through every symbol, accumulating segments separated</pre>
<pre>'this is pretty much impossible ElseIf tkl(nl).Token(0).Type = tkVarInit Then</pre>	'by one of []+* into acc acc = ""
<pre>s = "variable initialisation" 'same as above ElseIf tkl(nl).Token(0).Type = tkOperand Then</pre>	<pre>doing = Not TestCharset(Right(s, 1), "[]+*") For nl = Len(s) To 1 Step -1</pre>
s = "an operand" 'same as above Else	<pre>If TestCharset(Mid(s, nl, 1), "[]+*") Then</pre>
s = "an unknown token" 'same as above End If	I f doing Then GoSub DoltGoSub doing = <mark>True</mark>
'Because the tokenization procedure should not allow for any	acc = "" Else
'of these errors, this is an internal error Call Errr("pCompile.TokenizeCode: Internal Error (" +	'Just keep accumulating acc If doing Then acc = Mid(s, nl, 1) + acc
<pre>CStr(tkl(nl).CodeLine) + "): a line cannot start with " + s + ". Contact the author.")</pre>	End If Next
End If ElseIf UBound(tkl(nl).Token) + 1 = 2 Then	'Run it once at the end in case s doesn't end with any of []+*
<pre>If tkl(nl).Token(0).Type = tkOpcode And tkl(nl).Token(1).Type = tkOperand Then</pre>	Gosub DoltGoSub GoTo EndOfGoSub
<pre>b = True ElseIf tkl(nl).Token(0).Type = tkVarDecl And</pre>	DoltGoSub: 'Analyse
<pre>tkl(nl).Token(1).Type = tkVarInit Then b = True</pre>	<pre>If UCase(Left(acc, 7)) = "OFFSET(" And Right(acc, 1) = ")" Then</pre>
Else Call AddErr("Invalid token combination: """ +	<pre>'Have an offset. Remove OFFSET() completely s = Left(s, nl) + Mid(acc, 8, Len(acc) - 8) + Mid(s, nl</pre>
tkName(tkl(nl).Token(0).Type) + """ and """ + tkName(tkl(nl).Token(1).Type) + """", tkl(nl).CodeLine, "EC1002")	+ Len(acc) + 1) ElseIf Not OperandIsRg(acc) Then
End If ElseIf UBound(tkl(nl).Token) + 1 = 3 Then	'Have either a number or a variable name If Not TestCharset(Left(acc, 1), "-0123456789") Then
<pre>If tkl(nl).Token(0).Type = tkOpcode And tkl(nl).Token(1).Type = tkOperand And tkl(nl).Token(2).Type =</pre>	<pre>'Definitely not a number. Enclose in [] s = Left(s, nl) + "[" + acc + "]" + Mid(s, nl +</pre>
tkOperand Then b = True	Len(acc) + 1) End If
Else Call AddErr("Invalid token combination: """ +	End If Return
<pre>tkName(tkl(nl).Token(0).Type) + """, """ + tkName(tkl(nl).Token(1).Type) + """ and """ +</pre>	EndOfGoSub: Save the analysed string back
<pre>tkName(tkl(nl).Token(1).Type) + """", tkl(nl).CodeLine, "EC1003") End If</pre>	.TknLine(i).Token(A).Text = s End If
<pre>ElseIf UBound(tkl(nl).Token) + 1 > 3 Then Call AddErr("A line cannot contain more than three tokens.</pre>	Next Next
This line contains " + CStr(UBound(tkl(nl).Token) + 1) + " tokens.", tkl(nl).CodeLine, "EC1004")	End With
Else Call Errr("pCompile.CompilePass1: token line with less than	End Sub
1 token encountered. Contact the author.") End If	'' ' Private Sub CompilePass2()
'Skip next stage if token line is not valid If Not b Then nl = nl + 1: GoTo nnext	DESCRIPTION: Compilation pass 2: code generation
nl = nl + 1	 1. Generate code for all correct instructions 2. Generate errors/warnings for incorrect instructions
nnext: Next	 3. Generate a label-to-address list used in backpatching 4. Generate a "variable requested" list for backpatching
··	' PARAMETERS:
' Return tkn'	' p - should contain tokenized program (Tkn*)
'Check if nl is what we thought it would be	' OUTPUT: p.Code containing compiled machine code.
<pre>If nl <> UBound(tkl) + 1 Then Call Errr("pCompile.TokenizeCode: Internal Error (N/A):</pre>	Private Sub CompilePass2() Dim tli As Integer 'token line loop var
predicted token line count is not equal to actual token line count. Contact the author.")	Dim tl As TpTokenLine 'current token line Dim ctl As String 'compiled token line
End If	Dim i As Integer, 11 As Long, 12 As Long, 13 As Long
ReDim .TknLine(-1 To nl - 1) For i = 0 To UBound(.TknLine)	Dim t As String, s As String
<pre>ReDim .TknLine(i).Token(-1 To UBound(tkl(i).Token)) .TknLine(i).CodeLine = tkl(i).CodeLine</pre>	With Proj.P
For A = 0 TO UBound (.TknLine(i).Token) .TknLine(i).Token(A).Text = tkl(i).Token(A).Text	'Initialise structures .Code = ""
.TknLine(i).Token(A).Type = tkl(i).Token(A).Type Next	ReDim .Ref(-1 To -1)
Next	<pre>'Loop through all token lines For tli = 0 To UBound(.TknLine)</pre>
' Deal with variables'	tl.CodeLine = .TknLine(tli).CodeLine tl.CodeOffset = .TknLine(tli).CodeOffset

tl.Token = .TknLine(tli).Token ctl = ""	GoTo NextTokenLine End If
Store offset of the beginning of the token line .TknLine(tli).CodeOffset = Len(.Code)	''
۱ <u></u> ۱	<u>'</u>
' <u></u> ' ' <u></u> '	<pre>If tl.Token(0).Type <> tkOpcode Then Errr ("pCompile.CompilePass2: first token in token line is</pre>
''	opcode. Contact the author.") Exit Sub
<pre>'Validity check If UBound(tl.Token) < 0 Then</pre>	End If
Call Errr("pCompile.CompilePass2: empty token line at	'O <u>p</u> code name
<pre>pmpilation stage 2. Contact the author.") Exit Sub</pre>	<pre>t = LCase(tl.Token(0).Text)</pre>
End If	' Type 0 opcodes'
··	··
'=== LABEL TOKEN ==='	<pre>For i = 0 To UBound(cdType0) If t = cdType0(i).Opcode Then</pre>
<pre>If tl.Token(0).Type = tkLabel Then 'Check label name</pre>	<pre>If UBound(tl.Token) + 1 > 1 Then Call AddErr("Opcode takes 0 operands, not " +</pre>
<pre>If Not TestCharset(Left(tl.Token(0).Text,</pre>	CStr(UBound(tl.Token) + 1) + ".", tl.CodeLine, "EC2001")
<pre>en(tl.Token(0).Text) - 1), CharsetLabel) Then Call AddErr("Invalid label name - '" +</pre>	ctl = "" Else
<pre>eft(tl.Token(0).Text, Len(tl.Token(0).Text) - 1) + "'",</pre>	ctl = Chr(cdType0(i).Code)
L.CodeLine, "EC2008") GoTo NextTokenLine	End If GoTo NextTokenLine
<pre>End If If UCase(tl.Token(0).Text) = "A" Or UCase(tl.Token(0).Text) =</pre>	End If Next
3" Or UCase(tl.Token(0).Text) = "C" Or UCase(tl.Token(0).Text) =	Next
<pre>O" Or UCase(tl.Token(0).Text) = "E" Then Call AddErr("Label name cannot be same as register name - '"</pre>	'' ' Type 1 opcodes'
<pre>Left(tl.Token(0).Text, Len(tl.Token(0).Text) - 1) + "'",</pre>	··
L.CodeLine, "EC2009") GoTo NextTokenLine	<pre>For i = 0 To UBound(cdType1) If t = cdType1(i).0pcode Then '(lshr,ashl,rol,rcl etc)</pre>
End If 'Add reference	'Check param count If UBound(tl.Token) + 1 <> 3 Then
ReDim Preserve .Ref(-1 To UBound(.Ref) + 1)	Call AddErr("Opcode takes 2 operands, not " +
.Ref(UBound (.Ref)).Addr = Len (.Code) .Ref(UBound (.Ref)). Name = Left (tl.Token(<mark>0</mark>).Text,	CStr(UBound(tl.Token) + 1) + ".", tl.CodeLine, "EC2001") ctl = ""
<pre>an(tl.Token(0).Text) - 1) .Ref(UBound(.Ref)).CodeLine = tl.CodeLine</pre>	GOTO NextTokenLine End If
GoTo NextTokenLine	'Compile depending on type
End If	<pre>If UCase(tl.Token(1).Text) = "A" And OperandIsRgn(tl.Token(2).Text) Then</pre>
'' '==== VARDECL TOKEN ==='	<pre>[A/Rg] 11 = (Asc(UCase(t1.Token(2).Text)) - 66) * 32</pre>
<u>'</u> '	ctl = Chr(cdTypel(i).Code) + Chr(l1)
<pre>If tl.Token(0).Type = tkVarDecl Then If UCase(tl.Token(0).Text) = "DB" Then 'Byte variable</pre>	ElseIf OperandIsRg(tl.Token(1).Text) Then 'Rg/N
<pre>If OperandIsIm8(tl.Token(1).Text) Then</pre>	<pre>If Not StringIsInt(tl.Token(2).Text) Then</pre>
<pre>ctl = Chr(CIm8(tl.Token(1).Text, Len(.Code), Len(.Code), L.CodeLine))</pre>	Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(2).Text
<pre>ElseIf tl.Token(1).Text = "?" Then ctl = Chr(0)</pre>	"'.", tl.CodeLine, "EC2005") ctl = ""
Else	GoTo NextTokenLine
Call AddErr("Variable initialisation sequence is neither ' nor a valid constant.", tl.CodeLine, "EC2002")	<pre>End If 11 = CInt(tl.Token(2).Text)</pre>
<pre>End If ElseIf UCase(tl.Token(0).Text) = "DW" Then 'Word variable</pre>	<pre>If l1 < 0 Or l1 > 15 Then Call AddErr("The number of shift cycles must be betw</pre>
'Store reference	0 and 15.", tl.CodeLine, "EC2006")
ReDim Preserve .Vars(-1 To UBound(.Vars) + 1) .Vars(UBound(.Vars)).Addr = Len(.Code)	ctl = "" GoTo NextTokenLine
<pre>'Compile If OperandIsIm16(tl.Token(1).Text) Then</pre>	End If 12 = 128 + 11
'Initialisation can be a variable offset.We don't mind	<pre>13 = (Asc(UCase(tl.Token(1).Text)) - 66) * 32</pre>
<pre>'if it is - CIml6 would just add a ref for a backpatch ctl = Dec2Chr(CIml6(tl.Token(1).Text, Len(.Code),</pre>	<pre>If UCase(tl.Token(1).Text) = "A" Then 12 = 12 + 16 Els = 12 + 13</pre>
<pre>L.CodeLine), 2) ElseIf tl.Token(1).Text = "?" Then</pre>	<pre>ctl = Chr(cdTypel(i).Code) + Chr(l2) Else</pre>
ctl = Chr(0) + Chr(0)	'Error
Else Call AddErr("Variable initialisation sequence is neither	<pre>If UCase(tl.Token(1).Text) = "A" Then Call AddErr("Syntax error in operand OR opcode and</pre>
<pre>?' nor a valid constant.", tl.CodeLine, "EC2002")</pre>	operand incompatible. Offending operand: '" + tl.Token(2).Text
End If Else 'String literal	"'.", tl.CodeLine, "EC2005") Else
<pre>If Left(tl.Token(1).Text, 1) = """" And ight(tl.Token(1).Text, 1) = """" Then</pre>	Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text
<pre>ctl = Replace(Mid(tl.Token(1).Text, 2,</pre>	"'.", tl.CodeLine, "EC2005")
<pre>en(tl.Token(1).Text) - 2), """""", """") ElseIf tl.Token(1).Text = "?" Then</pre>	End If ctl = ""
ctl = ""	GoTo NextTokenLine
Else Call AddErr("DS variable should be initialised with either	End If GoTo NextTokenLine
or a string literal enclosed with """".", tl.CodeLine, "EC2019") End If	End If Next
End If	NCAL

Type 2 opcodes -For i = 0 To UBound(cdType2) If t = cdType2(i).Opcode Then '(and.or.add.sub etc) Check par If UBound(tl.Token) + 1 <> 3 Then Call AddErr("Opcode takes 2 operands, not " + CStr(UBound(tl.Token) + 1) + ".", tl.CodeLine, "EC2001") ctl = GoTo NextTokenLine End If 11 = 0 If (OperandIsRg(tl.Token(2).Text) And UCase(tl.Token(1).Text) = "A") Then 'A/Rq If UCase(tl.Token(2).Text) = "A" Then l1 = 4 Else l1 = (Asc(UCase(tl.Token(2).Text)) - 66) ctl = Chr(cdType2(i).Code) + Chr(l1) Elself (OperandIsRg(tl.Token(1).Text) And UCase(t1.Token(2).Text) = "A") Then 'Rq/A If UCase(tl.Token(1).Text) = "A" Then l1 = 4 Else l1 = (Asc(UCase(tl.Token(1).Text)) - 66) 11 = 11 + 8ctl = Chr(cdType2(i).Code) + Chr(l1) ElseIf UCase(tl.Token(1).Text) = "A" And OperandIsIm16(tl.Token(2).Text) Then A/I 11 = CIm16(tl.Token(2).Text, Len(ctl) + 1, tl.CodeLine) ctl = Chr(cdType2(i).Code + 1) + Dec2Chr(l1, 2) Else Call AddErr("16 bit immediate constant is out of range.", tl.CodeLine, "EC2007") ctl = "" End If ElseIf UCase(tl.Token(1).Text) = "A" And OperandIsMem(tl.Token(2).Text) Then 'A/M ctl = CompileMemoryAddressing(tl.Token(2).Text, Len(.Code) + 1, tl.CodeLine) If ctl <> "" Then
 ctl = Chr(cdType2(i).Code + 2) + ctl End If Else If UCase(tl.Token(1).Text) = "A" Or $\texttt{OperandIsRg(tl.Token(1).Text)} \ \textbf{Then}$ Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(2).Text + ".", tl.CodeLine, "EC2005") Else **Call** AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + '.", tl.CodeLine, "EC2005") End If ctl = " End If GoTo NextTokenLine End If Next For i = 0 To UBound(cdType3) If t = cdType3(i).Opcode Then '(inc,dec,neg,not,bswp) If UBound(tl.Token) + 1 <> 2 Then Call AddErr("Opcode takes 1 operand, not " + CStr(UBound(tl.Token) + 1) + ".", tl.CodeLine, "EC2001") ctl = GoTo NextTokenLine End If If OperandIsRgn(tl.Token(1).Text) Then ctl = Chr(cdType3(i).Code + (Asc(UCase(tl.Token(1).Text)) - 66)) ElseIf UCase(tl.Token(1).Text) = "A" Then ctl = Chr(cdType3(i).Code2) Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + "'. tl.CodeLine, "EC2005") ctl = "" End If GoTo NextTokenLine End If Next

For i = 0 To UBound(cdType4) If t = cdType4(i).Opcode Then '(jmp, If UBound(tl.Token) + 1 <> 2 Then Call AddErr("Opcode takes 1 operand, not " + CStr(UBound(tl.Token) + 1) + ".", tl.CodeLine, "EC2001") ctl = GoTo NextTokenLine End If s = tl.Token(1).Text If OperandIsMem(s) Then ctl = Chr(cdType4(i).Code) + CompileMemoryAddressing(tl.Token(1).Text, Len(.Code) + 1, tl.CodeLine) Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + "'. tl.CodeLine, "EC2005") ctl = End If GoTo NextTokenLine End If Next If t = "ld" Then If OperandIsIm16(tl.Token(2).Text) Then If OperandIsRgn(tl.Token(1).Text) Then 'Rn/Im16 = Chr(&H20 + (Asc(UCase(tl.Token(1).Text)) - 66)) + ctl Dec2Chr(CIm16(tl.Token(2).Text, Len(.Code) + 1, tl.CodeLine), 2) ElseIf UCase(tl.Token(1).Text) = "A" Then 'A/Iml6 ctl = Chr(&H24) + Dec2Chr(CIm16(tl.Token(2).Text, Len(.Code) + 1, tl.CodeLine), 2) ElseIf OperandIsMem(tl.Token(1).Text) Then Call AddErr("Cannot load a constant into a memory cell directly.", tl.CodeLine, "EC2014") ct1 = "" Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + "'. tl.CodeLine, "EC2005") ctl = "" End If ElseIf OperandIsRq(tl.Token(1).Text) And OperandIsRg(tl.Token(2).Text) Then 'R/R 11 = 0If UCase(tl.Token(1).Text) = "A" Then 11 = 11 + 128 Else 11 = 11 + (Asc(UCase(t1.Token(1).Text)) - 66) * 8 If UCase(tl.Token(2).Text) = "A" Then 11 = 11 + 64 Else 11 = 11 + (Asc(UCase(tl.Token(2).Text)) - 66) ctl = Chr(&H25) + Chr(11)ElseIf OperandIsRg(tl.Token(1).Text) And OperandIsMem(tl.Token(2).Text) Then 'R/Me em If UCase(tl.Token(1).Text) = "A" Then l1 = 4 Else l1 = (Asc(UCase(tl.Token(1).Text)) - 66) ctl = Chr(&H26) + Chr(l1) +CompileMemoryAddressing(tl.Token(2).Text, Len(.Code) + 2, tl.CodeLine) ElseIf OperandIsRq(tl.Token(2).Text) And OperandIsMem(tl.Token(1).Text) Then 'Mem/R If UCase(tl.Token(2).Text) = "A" Then l1 = 4 Else l1 = (Asc(UCase(tl.Token(2).Text)) - 66) ctl = Chr(&H27) + Chr(ll) +CompileMemoryAddressing(tl.Token(1).Text, Len(.Code) + 2, tl.CodeLine) Else If OperandIsIm16(tl.Token(1).Text) Then Call AddErr("Cannot load into a constant (first operand cannot be a constant).", tl.CodeLine, "EC2013") Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: check both.", tl.CodeLine, "EC2005" End If ctl = End If GoTo NextTokenLine

End If

If UCase(tl.Token(1).Text) = "A" Then

If t = "st" Then If OperandIsIm16(tl.Token(1).Text) Then If OperandIsRgn(tl.Token(2).Text) Then 'Im16/Rn ctl = Chr(&H30 + (Asc(UCase(tl, Token(2), Text)) - 66)) +Dec2Chr(CIm16(tl.Token(1).Text, Len(.Code) + 1, tl.CodeLine), 2)
ElseIf UCase(tl.Token(2).Text) = "A" Then 'Im16/A ctl = Chr(&H34) + Dec2Chr(CIm16(tl.Token(1).Text, Len(.Code) + 1, tl.CodeLine), 2)
ElseIf OperandIsMem(tl.Token(2).Text) Then Call AddErr("Cannot store a constant in a memory cell directly.", tl.CodeLine, "EC2016") ctl = "' Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(2).Text + "'.", tl.CodeLine, "EC2005") ctl = "" End If ElseIf OperandIsRg(tl.Token(2).Text) And OperandIsRg(tl.Token(1).Text) Then idls. <u>'R/R</u> ' = 0 If UCase(tl.Token(2).Text) = "A" Then 11 = 11 + 128 Else 11 = 11 + (Asc(UCase(tl.Token(2).Text)) - 66) * 8 If UCase(tl.Token(1).Text) = "A" Then 11 = 11 + 64 Else 11 =
11 + (Asc(UCase(tl.Token(1).Text)) - 66) ctl = Chr(&H35) + Chr(11)ElseIf OperandIsRg(tl.Token(2).Text) And OperandIsMem(tl.Token(1).Text) Then 'Mem/R If UCase(tl.Token(2).Text) = "A" Then l1 = 4 Else l1 = (Asc(UCase(tl.Token(2).Text)) - 66) ctl = Chr(&H36) + Chr(11) + CompileMemoryAddressing(tl.Token(1).Text, Len(.Code) + 2, tl.CodeLine) ${\tt Elseif} \ {\tt OperandIsRg(tl.Token(1).Text)}$ And OperandIsMem(tl.Token(2).Text) Then 'R/Mem If UCase(tl.Token(1).Text) = "A" Then l1 = 4 Else l1 = (Asc(UCase(tl.Token(1).Text)) - 66) ctl = Chr(&H37) + Chr(l1) +CompileMemoryAddressing(tl.Token(2).Text, Len(.Code) + 2, tl.CodeLine) Else If OperandIsIm16(tl.Token(2).Text) Then Call AddErr("Cannot store in a constant (second operand cannot be a constant).", tl.CodeLine, "EC2015") Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: check both.", tl.CodeLine, "EC2005" End If ctl = " End If GoTo NextTokenLine End If '`` PUSH ```'
If t = "push" Then If UCase(tl.Token(1).Text) = "A" Then 'push A
ctl = Chr(&H10) ElseIf OperandIsRgn(tl.Token(1).Text) Then 'push Rn ct] = Chr(0 + Asc(UCase(t], Token(1), Text)) - 66) ElseIf OperandIsIm16(tl.Token(1).Text) Then 'push I ctl = Chr(&H12) + Dec2Chr(CIm16(tl.Token(1).Text, Len(.Code) + 1, tl.CodeLine), 2) Else **Call** AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + "'.", tl.CodeLine, "EC2005") ctl = "" End If GoTo NextTokenLine End If If t = "pop" Then

'pop A ctl = Chr(&H11)ElseIf OperandIsRgn(tl.Token(1).Text) Then 'pop Rr ctl = Chr(4 + Asc(UCase(tl.Token(1).Text)) - 66) Else **Call** AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + "'.", tl.CodeLine, "EC2005") ctl = "" End If GoTo NextTokenLine End If LEA "lea" Then If t If UCase(tl.Token(1).Text) = "A" Then ctl = Chr(&H2F) + CompileMemoryAddressing(tl.Token(2).Text, Len(.Code) + 1, tl.CodeLine) ElseIf OperandIsRgn(tl.Token(1).Text) Then 'lea Rn,M ctl = Chr(&H8 + Asc(UCase(tl.Token(1).Text)) - 66) + CompileMemoryAddressing(tl.Token(2).Text, Len(.Code) + 1, tl.CodeLine) Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + "'.", tl.CodeLine, "EC2005") ctl = "" End If GoTo NextTokenLine End If If t = "xchg" Then
If UCase(tl.Token(1).Text) = "A" Then 'xchg A.Rn
ctl = Chr(&HF0 + Asc(UCase(tl.Token(2).Text)) - 66) ElseIf UCase(tl.Token(2).Text) = "A" Then ctl = Chr(&HF0 + Asc(UCase(tl.Token(1).Text)) - 66) "C") And (UCase(tl.Token(1).Text) = "D")) Then ng c,d ctl = Chr(&HE6)ElseIf ((UCase(tl.Token(1).Text) = "C") And(UCase(tl.Token(2).Text) = "E")) Or ((UCase(tl.Token(2).Text) = "C") And (UCase(tl.Token(1).Text) = "E")) Then 'xchg c,e ctl = Chr(&HE7) Elself ((UCase(tl,Token(1),Text) = "D") And(UCase(tl.Token(2).Text) = "E")) Or ((UCase(tl.Token(2).Text) ="D") And (UCase(tl.Token(1).Text) = "E")) Then 'xchg d,e ctl = Chr(&HF4) ElseIf ((UCase(tl.Token(1).Text) = "B") And (UCase(tl.Token(2).Text) = "C")) Or ((UCase(tl.Token(2).Text) = "B") And (UCase(tl.Token(1).Text) = "C")) Then 'xchg b,c = Chr(&HF5) ct1 ElseIf ((UCase(tl.Token(1).Text) = "B") And $\begin{array}{l} (\texttt{UCase(tl.Token(2).Text)} = \texttt{"D"})) \texttt{ Or } ((\texttt{UCase(tl.Token(2).Text)} = \texttt{"B"}) \texttt{ And } (\texttt{UCase(tl.Token(1).Text)} = \texttt{"D"})) \texttt{ Then} \end{array}$ 'xchg b,d ctl = Chr(&HF6) ElseIf ((UCase(tl.Token(1).Text) = "B") And (UCase(tl.Token(2).Text) = "E")) Or ((UCase(tl.Token(2).Text) = "B") And (UCase(tl.Token(1).Text) = "E")) Then 'xchg b,e ctl = Chr(&HF7)Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: check both.", tl.CodeLine, "EC2005") ctl = "" End If GoTo NextTokenLine End If INT ` If t = "int" Then If OperandIsIm8(tl.Token(1).Text) Then $\label{eq:linear} \begin{array}{l} \mbox{int I8} \\ \mbox{ctl} = \mbox{Chr}(\&H74) \ + \ \mbox{Clm8(tl.Token(1).Text}, \ \mbox{Len(.Code)} \ + \ \mbox{1}, \end{array}$ Len(.Code) + 2, tl.CodeLine)

Else ElseIf OperandIsIm8(tl.Token(1).Text) Then Call AddErr("Operand for INT must be an 8 bit immediate If UCase(tl.Token(2).Text) = "A" Then constant.", tl.CodeLine, "EC2017")
 ctl = "" 'out I8,A ct] = Chr(&HD6) + Chr(CIm8(t], Token(1), Text, Len(, Code) + End If 1. -1. tl.CodeLine)) GoTo NextTokenLine ElseIf OperandIsRgn(tl.Token(2).Text) Then End If 'out I8,R ctl = Chr(&HD0 + Asc(UCase(tl.Token(2).Text)) - 66) + Chr(CIm8(tl.Token(1).Text, Len(.Code) + 1, -1, tl.CodeLine)) Else If OperandIsIm8(tl.Token(2).Text) Then Call AddErr("Syntax error in operand OR opcode and operand If UCase(tl.Token(1).Text) = "A" Then incompatible. Offending operand: '" + tl.Token(2).Text + "'.", tl.CodeLine, "EC2005") 'in A,18 ctl = "" ctl = Chr(&HE5) + Chr(CIm8(tl.Token(2).Text, Len(.Code) + 1, -1, tl.CodeLine)) End If ElseIf OperandIsRgn(tl.Token(1).Text) Then **ElseIf** OperandIsIm16(t].Token(1).Text) Then Call AddErr("Port address has to be an 8 bit immediate 'in Rn,I8 ctl = Chr(&HD8 + Asc(UCase(tl.Token(1).Text)) - 66) + constant (0 to 255).", tl.CodeLine, "EC2018") Chr(CIm8(tl.Token(2).Text, Len(.Code) + 1, -1, tl.CodeLine)) ctl = "" Else Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + "'.", Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: check both.", tl.CodeLine, tl.CodeLine, "EC2005") "EC2005") ctl = "" ctl = End If End If ElseIf OperandIsRg(tl.Token(2).Text) Then GoTo NextTokenLine End If If OperandIsRg(tl.Token(1).Text) Then 'in R1,R2 If UCase(tl.Token(1).Text) = "A" Then 11 = 32 Else 11 = 'All instruction processing should end with a GoTo (Asc(UCase(tl.Token(1).Text)) - 66) * 4extTokenLine. If UCase(tl.Token(2).Text) = "A" Then 11 = 11 + 16 Else 11 'So if we arrive here, opcode was not found. Add error messa = 11 + (Asc(UCase(tl.Token(2).Text)) - 66) Call AddErr("Opcode not recognized: '" + t + "'. Check ctl = Chr(&HE4) + Chr(11)spelling.", tl.CodeLine, "EC2010") Else Call AddErr("Syntax error in operand OR opcode and operand NextTokenLine: incompatible. Offending operand: '" + tl.Token(1).Text + "'.", tl.CodeLine, "EC2005") ctl = "" End If ElseIf OperandIsIm16(tl.Token(2).Text) Then Call AddErr("Port address must be an 8 bit immediate .Code = .Code + ctl constant (0 to 255).", tl.CodeLine, "EC2018") Next ctl = " Else 'Calculate offset to code line conversion Call AddErr("Syntax error in operand OR opcode and operand ReDim .Code O2L(-1 To Len(.Code) - 1) incompatible. Offending operand: '" + tl.Token(2).Text + For i = Len(.Code) - 1 To 0 Step -1 For tli = UBound(.TknLine) To 0 Step -1 tl.CodeLine, "EC2005") ctl = If .TknLine(tli).CodeOffset <= i Then</pre> End If .Code O2L(i) = .TknLine(tli).CodeLine GoTo NextTokenLine GoTo takeNext End If End If Next .Code_02L(i) = 0 If t = "out" Then take If OperandIsRg(tl.Token(1).Text) And Next Calculate source code line to offset conversion ReDim .Code_L20(-1 To UBound(.AsmLine)) OperandIsRg(t], Token(2), Text) Then 'out R1,R2 If UCase(tl.Token(1).Text) = "A" Then l1 = 32 Else l1 = For i = UBound(.AsmLine) To 0 Step -1 (Asc(UCase(tl.Token(1).Text)) - 66) * 4For tli = 0 To UBound(.TknLine) If UCase(tl.Token(2).Text) = "A" Then 11 = 11 + 16 Else 11 = r tli = UBound(.TknLine) To 0 Step If .TknLine(tli).CodeLine = i Then .Code L20(i) = .TknLine(tli).CodeOffset ElseIf OperandIsIm8(tl.Token(1).Text) And GoTo takeNext2 OperandIsIm16(tl.Token(2).Text) Then End If 'out I8,I Next ctl = Chr(&HD7) + Chr(CIm8(tl.Token(1).Text, Len(.Code) + 1, If i = UBound(.AsmLine) Then .Code_L2O(i) = Len(.Code) Else -1, tl.CodeLine)) + Dec2Chr(CIm16(tl.Token(2).Text, Len(.Code) + Code L20(i) = .Code L20(i + 1)2, tl.CodeLine), 2) takeNext2: ElseIf OperandIsIm16(tl.Token(1).Text) And Next OperandIsIml6(tl.Token(2).Text) Then ReDim Preserve .Code_L20(-1 To UBound(.Code_L20) + 1) Call AddErr("Port address has to be an 8 bit immediate .Code_L20(UBound(.Code_L20)) = Len(.Code) + 1 'point to last constant (0 to 255).", tl.CodeLine, "EC2018") Halt ctl = " ElseIf OperandIsIm16(tl.Token(2).Text) Then End With If UCase(tl.Token(1).Text) = "A" Then End Sub 'out A,I ctl = Chr(&HD5) + Dec2Chr(CIm16(tl.Token(2).Text, Len(.Code) + 1, tl.CodeLine), 2) ElseIf OperandIsRgn(tl.Token(1).Text) Then 'out R,I DESCRIPTION: Compilation pass 3: address backpatching ctl = Chr(&HE0 + Asc(UCase(tl.Token(1).Text)) - 66) + 1. Check for "label already declared' Check for "undeclared reference Dec2Chr(CIm16(tl.Token(2).Text, Len(.Code) + 1, tl.CodeLine), 2) Else Call AddErr("Syntax error in operand OR opcode and operand incompatible. Offending operand: '" + tl.Token(1).Text + "'.", tl.CodeLine, "EC2005") p - should contain tokenized program (Tkn*) ctl = "" End If

Listings

Private Sub CompilePass3()	ctl = cl), 2)
Dim i As Integer, A As Integer, t As Integer	Else
With Dura D	ma =
With Proj.P	If Or
'Check for "label already declared"	·
For i = 0 To UBound(.Ref)	۱ <u>–</u> ۰
<pre>For A = i + 1 To UBound(.Ref) If UCase(.Ref(i).Name) = UCase(.Ref(A).Name) Then</pre>	ct] Else
Call AddErr("Label already declared: '" + .Ref(A).Name +	115e
"'. Previous declaration on line " + CStr(.Ref(i).CodeLine) + ".",	·
.Ref(A).CodeLine, "EC3001")	'
End If Next	ct] Else
Next	'
	۰ <u>-</u> -
'Check for "undeclared reference"	' 11
<pre>For i = 0 To UBound(.Backpatch) For A = 0 To UBound(.Ref)</pre>	11 If
<pre>If UCase(.Backpatch(i).Name) = UCase(.Ref(A).Name) Then GoTo</pre>	1
fnd	r
Next Call AddErr("Undeclared reference: '" + .Backpatch(i).Name +	End
"'.", .Backpatch(i).CodeLine, "EC3002")	
fnd:	cl), <mark>2</mark>)
Next	r
'Backpatch	End
For i = 0 To UBound(.Backpatch)	If
For A = 0 To UBound(.Ref)	1
<pre>If .Backpatch(i).Name = .Ref(A).Name Then If .Backpatch(i).IsDW Then</pre>	0, 1, 2 0
Mid(.Code, .Backpatch(i).Addr + 1, 2) =	0, 1, 20
Dec2Chr(.Ref(A).Addr, 2)	
Else	H
t = .Ref(A).AddrBackpatch(i).RelTo If t < 0 Then t = <mark>256</mark> + t	1
<pre>Mid(.Code, .Backpatch(i).Addr + 1, 1) = Chr(t)</pre>	
End If	0, 1, 2 0
End If Next	
Next	I
	1
'Match variable references against reference names For i = 0 To UBound(.Vars)]
.Vars(i).Name = "[unnamed]"	T End
.Vars(i). Name = "[unnamed]" Next	End 'Ch
Next For i = 0 To UBound(.Ref)	End 'Ch 'ju
Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars)	End 'Ch 'ju 11
Next For i = 0 To UBound(.Ref)	End 'Ch 'ju
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next</pre>	End 'Ch 'ju ll ctl End I End If
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name</pre>	End 'Ch 'ju ll ctl End I End If <u>'Check</u>
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end</pre>	End 'Ch 'ju ll ctl End I End If
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars)</pre>	End 'Ch 'ju End If End If 'Check If ctl compilati
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next /Add HALT at the end .Code = .Code + Chr(&H75)</pre>	End 'Ch ctl End If Check If ctl compilati Compilet
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end</pre>	End 'Ch 'ju End If End If 'Check If ctl compilati
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Next 'Add HALT at the end .Code = .Code + Chr(&H75) End With</pre>	End 'C' 'Ju End I End If 'Check If ctl compilati 'Return Compile End Funct
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub</pre>	End 'Ch ctl End If Check If ctl compilati Compilet
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub '</pre>	End 'Ch 'ju In ctl End If check If ctl Compilati 'Return Compilat End Funct
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Next Next Mext Ind With End With End Sub Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String</pre>	End 'C' 'Ju End I End If 'Check If ctl compilati 'Return Compile End Funct
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub '</pre>	End 'C' 'Ju ctl End If 'Check If ctl compilati 'Return Compile End Funct
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Next Next Mext Ind With End With End With End Sub Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String</pre>	End 'C' 'ju Il End If 'Check If ctl Compilati 'Return Compilati 'Return 'Private 'Private 'PARAMET 'C - a '
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub '</pre>	End 'C' 'Ju ctl End If 'Check If ctl compilati 'Return Compile End Funct
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Mext Mext Mext Meth End Sub</pre>	End 'C' 'ju ll ctl End If 'Check If ctl compilati 'Return Compilati 'Return Compilati 'Private 'Private 'End Funct
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub '</pre>	End 'C' 'Ju ctl End If 'Check If ctl compilati 'Return Compilati 'Return Compilati 'Return 'Private 'Private 'RETURNS 'C - a 'I' 'RETURNS 'L' 'R' 'R' 'R' 'R' 'R' 'R' 'R'
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub ' ' ' Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String ' DESCRIPTION: compiles addressing from assembly language ' into machine code. ' ' PARAMETERS: ' c - contains addressing operand in assembly language ' adr - machine code offset for where the addressing will ' ' ' ' ' Parameters: ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' '</pre>	End 'C' 'Ju ctl End If 'Check If ctl compilati 'Return Compilati 'Return Compilati 'Private 'Private 'End Funct
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub '</pre>	End 'C' 'J' End If 'Check If ctl compilati 'Return Compilati 'Return 'Private 'Private 'Private 'Private 'End Funct 'Private 'Private 'I' 'Case(s) Operational Operations
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub </pre>	End 'C' 'Ju ctl End If 'Check If ctl compilati 'Return Compilati 'Return Compile End Funct 'Private 'Private 'Private I If UCase(s) Opera Else Opera End If
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next 'Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub '</pre>	End 'C' 'J' End If 'Check If ctl compilati 'Return Compilati 'Return 'Private 'Private 'Private 'Private 'End Funct 'Private 'Private 'I' 'Case(s) Operational Operations
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Add HALT at the end .Code = .Code + Chr(&H75) End With End Sub </pre>	End 'C' 'Ju ctl End If 'Check If ctl compilati 'Return Compilati 'Return Compile End Funct 'Private 'Private 'Private I If UCase(s) Opera Else Opera End If
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Mext Mext Mext Meth End With End With End Sub</pre>	End Chi End If End If Check If ctl compilati Return Compilati
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Mext Mext Muth End Sub DESCRIPTION: compileMemoryAddressing(c As String, add As Long, cl As Integer) As String DESCRIPTION: compiles addressing from assembly language into machine code. PARAMETERS: c - contains addressing operand in assembly language addr - machine code offset for where the addressing will start (required for backpatching) cl - code line containing the addressing for error log. p - program structure with error log etc RETURNS: machine codes ready to be added to p.Code Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String Dim ctl As String, ma As String, ll As Long, l2 As Long</pre>	End Chi End I End I End I Check If ctl compilati Compilati Compilati Compilati Compilati Compilati Private Private Private UCase(s) Opera Else Opera End If End Funct
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Mext Mext Mext Meth End With End With End Sub</pre>	End 'Check If ctl End If 'Check If ctl compilet 'Return Compilet 'Return Compilet 'Private 'Private 'Private 'Return 'Private 'Private 'Private 'Return 'Private 'Private 'Private 'End Funct 'Private 'End Funct 'Private 'End If End Funct 'Private 'Priv
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Next Mext Muth End With End With End Sub Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String DESCRIPTION: compiles addressing from assembly language into machine code. PARAMETERS: c - contains addressing operand in assembly language adr - machine code offset for where the addressing will start (required for backpatching) cl - code line containing the addressing for error log. p - program structure with error log etc RETURNS: machine codes ready to be added to p.Code Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String Dim ctl As String, ma As String, ll As Long, l2 As Long ma = c ctl = ""</pre>	End 'C' 'J' End If 'Check If ctl compilati 'Return Compilati 'Return 'Private 'Private 'Private 'Private 'End Funct 'Private 'End Funct 'Private
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Mext Mext Mext Mext Muth End Sub</pre>	End Chi End If End If Check If ctl compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Private Private If Cca Return Compilati Return Compilati Return Compilati Return Private If Cca Second Private If Cca Second Private Private If Cca Second Private If Cca Second Private Private If Cca Second Private Pri
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Next Mext Muth End With End With End Sub Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String DESCRIPTION: compiles addressing from assembly language into machine code. PARAMETERS: c - contains addressing operand in assembly language adr - machine code offset for where the addressing will start (required for backpatching) cl - code line containing the addressing for error log. p - program structure with error log etc RETURNS: machine codes ready to be added to p.Code Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String Dim ctl As String, ma As String, ll As Long, l2 As Long ma = c ctl = ""</pre>	End Chi End If End If Check If ctl compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Return Compilati Compilati Compilati Compilati Compilati Compilati Return Compilati C
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Mext Mext Mext Mext Mext Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String DESCRIPTION: compiles addressing from assembly language into machine code. PARAMETERS: c - contains addressing operand in assembly language adr - machine code offset for where the addressing will start (required for backpatching) cl - code line containing the addressing for error log. Private Function CompileMemoryAddressing(c As String, adr As Long, cl - conde line codes ready to be added to p.Code Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String Dim ctl As String, ma As String, ll As Long, l2 As Long ma = c ctl = "" Detect type If Left(ma, 2) = "[[" And Right(ma, 2) = "]]" Then </pre>	End 'C' 'J' End If 'Check If ctl compilati 'Return Compilati 'Return Compilati 'Return 'Private 'Private 'Private 'Private 'End Funct 'Private 'Private 'Returns 'Private
<pre>Next For i = 0 To UBound(.Ref) For A = 0 To UBound(.Vars) If .Ref(i).Addr = .Vars(A).Addr Then .Vars(A).Name = .Ref(i).Name Next Next Mext Next Mext Mext Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String DESCRIPTION: compiles addressing from assembly language into machine code. PARAMETERS: c - contains addressing operand in assembly language adr - machine code offset for where the addressing will start (required for backpatching) cl - code line containing the addressing for error log. p - program structure with error log etc RETURNS: machine codes ready to be added to p.Code Private Function CompileMemoryAddressing(c As String, adr As Long, cl As Integer) As String Dim ctl As String, ma As String, ll As Long, l2 As Long ma = c ctl = "" Detect type If Left(ma, 2) = "[[" And Right(ma, 2) = "]]" Then</pre>	End 'Ch 'Ji End If Check If ctl compilati 'Return Compilati 'Return Compilati 'Return 'Private 'Private 'Private 'Private 'Returns 'Private 'Returns 'Private

= Chr(1) + Dec2Chr(CIm16(Mid(ma, 3, Len(ma) - 4), adr + 1, Mid(ma, 2, Len(ma) - 2) perandIsIm16(ma) Then Memory Direct Immediate 1 = Chr(0) + Dec2Chr(CIm16(ma, adr + 1, cl), 2) If OperandIsRgn(ma) Then - Memory Indirect Register 1 = Chr(64 + (Asc(UCase(ma))) -66)) UCase(Left(ma, 2)) = "B+" Then 11 = 11 + 64ma = Mid(ma, 3)nd If InStr(ma, "+") > 0 Then ctl = Dec2Chr(CIm16(Mid(ma, InStr(ma, "+") + 1), adr + 1, ma = Left(ma, InStr(ma, "+") - 1) 11 = 11 + 32nd If InStr(ma, "*") > 0 Then If Not StringIsLong(Mid(ma, InStr(ma, "*") + 1)) Then Call AddErr("Memory addressing scaling factor should be or 4.", cl, "EC2011") ctl = "" Return End If Call AddErr("Memory addressing scaling factor should be or 4.", cl, "EC2011") ctl = "" Return End If **If** 12 = **4** Then 12 = **3** '00=0,01=1,10=2,11=4 11 = 11 + 12 * 4ma = Left(ma, InStr(ma, "*") - 1) nd If ecks have been completed before, ust the indexation register B, C, D or E = 11 + (Asc(UCase(ma)) - 66) cl = Chr(11) + ctlIf = "" Then Call Errr("pCompile.CompilePass3: mem addr ion successful but returns nothing. Contact the author.") rn compiled addressing .eMemoryAddressing = ctl tion Function OperandIsRg(s As String) As Boolean IS: True if operand is a register (A,B,C,D,E) Function OperandIsRg(s As String) As Boolean ase(s) = "A" Or UCase(s) = "B" Or UCase(s) = "C" Or = "D" Or UCase(s) = "E" Then randIsRg = <mark>True</mark> randIsRg = False tion Function OperandIsRgn(s As String) As Boolean S: True if operand is a GP register (B,C,D,E)

Private Function OperandIsRgn(s As String) As Boolean If UCase(s) = "B" Or UCase(s) = "C" Or UCase(s) = "D" Or JCase(s) = "E" Then

Listings

OperandIsRgn = True	' NUMBER'
Else OperandIsRgn = False	''
End If	Dim s As String, testval As Long, minus As Boo.
End Function	<pre>If Len(c) = 0 Then GoTo IsNot</pre>
	s = UCase(c)
'' ' Private Function OperandIsMem(c As String) As Boolean '	<pre>minus = False If Left(s, 1) = "-" Then</pre>
	If $Len(s) = 1$ Then GoTo IsNot
' PARAMETERS: '	s = Mid(s, 2)
c - assembly language operand to be tested	minus = True
'	<pre>End If If Right(s, 1) = "H" Or Right(s, 1) = "B" Then</pre>
''	Then GoTo IsNot
Private Function OperandIsMem(c As String) As Boolean	'Check charset and try to convert (overflow wi
Dim s As String, st As String s = UCase(c)	If Right(s, 1) = "H" Then s = Left(s, Len(s) - 1)
'Checks	If Not TestCharset(Left(s, 1), "0123456789")
If $Len(s) < 3$ Then	<pre>If Not TestCharset(s, "0123456789ABCDEF") The</pre>
OperandIsMem = False	testval = Hex2Dec(s)
Exit Function End If	<pre>ElseIf Right(s, 1) = "B" Then s = Left(s, Len(s) - 1)</pre>
<pre>If Left(s, 1) <> "[" Or Right(s, 1) <> "]" Then</pre>	If Not TestCharset(s, "01") Then GoTo IsNot
OperandIsMem = False	<pre>testval = Bin2Dec(s)</pre>
Exit Function End If	<pre>Else If Not TestCharset(s, "0123456789") Then GoT</pre>
s = Mid(s, 2, Len(s) - 2)	testval = CLng(s)
	End If
'Indirect immediate	'Check range
<pre>If Left(s, 1) = "[" And Right(s, 1) = "]" Then OperandIsMem = OperandIsIml6(Mid(s, 2, Len(s) - 2))</pre>	<pre>If minus Then testval = -testval If testval < -128 Or testval > 255 Then GoTo Is</pre>
Exit Function	'Everything is fine
	OperandIsIm8 = True
'Indirect register	
ElseIf OperandIsRgn(s) Then OperandIsMem = True	<pre>ElseIf UCase(c) = "A" Or UCase(c) = "B" Or UCase UCase(c) = "D" Or UCase(c) = "E" Then</pre>
Exit Function	OperandIsIm8 = False
	Else
'Direct (immediate) ElseIf OperandIsIml6(s) Then	'' ' VARIABLE'
OperandIsMem = True	' VARIABLE'
Exit Function	
	<pre>OperandIsIm8 = TestCharset(c, CharsetLabel)</pre>
'Either memory indexed or not memory Else	End If
'Base register	Exit Function
If Left $(s, 2) = "B+"$ Then $s = Mid(s, 3)$	IsNot:
<pre>'Offset If InStr(s, "+") > 0 Then</pre>	OperandIsIm8 = False End Function
If OperandIsIm16(Mid(s, InStr(s, "+") + 1)) Then	
s = Left(s, InStr(s, "+") - 1)	(
Else	' Private Function OperandIsIm16(c As String) As Bo
OperandIsMem = False Exit Function	PARAMETERS:
End If	c - assembly language operand to be tested
End If	
<pre>'Scale If InStr(s, "*") > 0 Then</pre>	' RETURNS: True if operand is a 16-bit immediate ' constant, and range checks are passed.
st = Mid(s, InStr(s, "*") + 1)	'
<pre>If TestCharset(st, "0123456789") Then 'no range check here -</pre>	Private Function OperandIsIm16(c As String) As Boo
<pre>s = Left(s, InStr(s, "*") - 1) 'so can generate error Else 'msg with explanation</pre>	On Error GoTo IsNot
Else 'msg with explanation OperandIsMem = False 'later	'Check type depending on first symbol
Exit Function	If TestCharset(Left(c, 1), "-0123456789") Then
End If	··
End If 'Register	' NUMBER'
'What is left by now should be the central register which is	
compulsory	Dim s As String, testval As Long, minus As Boo
OperandIsMem = OperandIsRgn(s)	<pre>If Len(c) = 0 Then GoTo IsNot</pre>
End If End Function	s = UCase (c) minus = False
	If Left(s, 1) = "-" Then
··	If Len(s) = 1 Then GoTo IsNot
' Private Function OperandIsIm8(c As String) As Boolean '	s = Mid(s, 2) minus = True
' PARAMETERS:	End If
c - assembly language operand to be tested	If Right(s, 1) = "H" Or Right(s, 1) = "B" Then
	Then GoTo IsNot
' RETURNS: True if operand is an 8-bit immediate ' ' constant, and range checks are passed. '	<pre>'Check charset and try to convert (overflow wi If Right(s, 1) = "H" Then</pre>
''	s = Left(s, Len(s) - 1)
Private Function OperandIsIm8(c As String) As Boolean	<pre>If Not TestCharset(Left(s, 1), "0123456789")</pre>
On Error GoTo IsNot	<pre>If Not TestCharset(s, "0123456789ABCDEF") The testual = Her2Dec(a)</pre>
'Check type depending on first symbol	<pre>testval = Hex2Dec(s) ElseIf Right(s, 1) = "B" Then</pre>
<pre>If TestCharset(Left(c, 1), "-0123456789") Then</pre>	s = Left(s, Len(s) - 1)
۲ <u></u> ۲	<pre>If Not TestCharset(s, "01") Then GoTo IsNot</pre>

As String, testval As Long, minus As Boolean en(c) = 0 Then GoTo IsNot JCase(c) = False eft(s, 1) = "-" Then Len(s) = 1 Then GoTo IsNot Mid(s, 2) nus = True f ight(s, 1) = "H" Or Right(s, 1) = "B" Then If Len(s) = 1 IsNot ck charset and try to convert (overflow will be trapped) ight(s, 1) = "H" Then Left(s, Len(s) - 1) Not TestCharset(Left(s, 1), "0123456789") Then GoTo IsNot Not TestCharset(s, "0123456789ABCDEF") Then GoTo IsNot stval = Hex2Dec(s) If Right(s, 1) = "B" Then
= Left(s, Len(s) - 1) Not TestCharset(s, "01") Then GoTo IsNot stval = Bin2Dec(s) Not TestCharset(s, "0123456789") Then GoTo IsNot stval = CLng(s) f ck range inus Then testval = -testval estval < -128 Or testval > 255 Then GoTo IsNot ything is fine andIsIm8 = True UCase(c) = "A" Or UCase(c) = "B" Or UCase(c) = "C" Or = "D" Or UCase(c) = "E" Then andIsIm8 = False andIsIm8 = TestCharset(c, CharsetLabel) unction dIsIm8 = False cion Function OperandIsIm16(c As String) As Boolean sembly language operand to be tested ant, and range checks are passed. Punction OperandIsIml6(c As String) As Boolean or GoTo IsNot type depending on first symbol Charset(Left(c, 1), "-0123456789") Then As String, testval As Long, minus As Boolean en(c) = 0 Then GoTo IsNot JCase(c) = False eft(s, 1) = "-" Then
Len(s) = 1 Then GoTo IsNot Mid(s, 2) nus = True Гf ight(s, 1) = "H" Or Right(s, 1) = "B" Then If Len(s) = 1IsNot ck charset and try to convert (overflow will be trapped) ight(s, 1) = "H" Then Left(s, Len(s) - 1) Not TestCharset(Left(s, 1), "0123456789") Then GoTo IsNot Not TestCharset(s, "0123456789ABCDEF") Then GoTo IsNot stval = Hex2Dec(s)If Right(s, 1) = "B" Then

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11	' ' Private Function CleanSpaces(ByVal src As String)
' Private Sub AddErr(ByRef ErrL As PErrLog, Message As '	· · · · · · · · · · · · · · · · · · ·
' String, LineNum As Integer, ErrCode As String) '	' DESCRIPTION: Converts all tab/" "/"," sequences wit
	' a single space character.
' DESCRIPTION: Appends an error to the given error log '	·
	PARAMETERS:
' PARAMETERS: ' ' ErrL - error log to add to '	' src - string to process
' Message - error message to be displayed '	' RETURNS: processed string with converted characters
' LineNum - offending line number in source code '	·
' ErrCode - error code to manage help '	Private Function CleanSpaces(ByVal src As String)
··	Dim i As Integer, b As Boolean
Private Sub AddErr(Message As String, LineNum As Integer, ErrCode	Dim s As String, c As String
As String)	
Dim i As Integer	b = True
'Add element	s = ""
ReDim Preserve Proj.P.ErrL.lError(-1 To	For i = 1 To Len(src)
UBound(Proj.P.ErrL.lError) + 1)	c = Mid(src, i, 1)
ReDim Preserve Proj.P.ErrL.sError(-1 To UBound(Proj.P.ErrL.sError) + 1)	If Asc(c) > 127 Then GoTo skp If b Then
ReDim Preserve Proj.P.ErrL.nError(-1 To	If $c = Chr(9)$ Or $c = "$ " Or $c = ","$ Then
UBound (Proj. P. ErrL. nError) + 1)	s = s + "
'Set new element	b = False
Proj.P.ErrL.lError(UBound(Proj.P.ErrL.lError)) = LineNum	Else
Proj.P.ErrL.SError(UBound(Proj.P.ErrL.SError)) = Message	s = s + c
Proj.P.ErrL.nError(UBound(Proj.P.ErrL.nError)) = ErrCode	End If
End Sub	Else
	If c <> $Chr(9)$ And c <> " " And c <> "," Then
۱۱	s = s + c
' Private Sub AddWng(ByRef ErrL As PErrLog, Message As '	b = True
' String, LineNum As Integer, WngCode As String) '	End If
	End If
DESCRIPTION: Appends a warning to given error log	skp:
· · · · · · · · · · · · · · · · · · ·	Next
' ErrL - error log to add to '	CleanSpaces = s
' Message - warning message to be displayed '	End Function
' LineNum - offending line number in source code '	
' WngCode - warning code to manage help '	'Warnings:
''	'WC1001, WC1002
Private Sub AddWng(Message As String, LineNum As Integer, WngCode	
As String)	'Errors:
Dim i As Integer	'EC1001, EC1002, EC1003, EC1004
'Add element	
ReDim Preserve Proj.P.ErrL.lWarning(-1 To	'EC2001 - operand takes X operands, not Y.
UBound(Proj.P.ErrL.lWarning) + 1)	'EC2002
ReDim Preserve Proj.P.ErrL.sWarning(-1 To UBound(Proj.P.ErrL.sWarning) + 1)	'EC2005 - Syntax error in operand OR opcode and opera incompatible. Offending operand: 'Y'.
ReDim Preserve Proj.P.ErrL.nWarning(-1 To	'EC2006
UBound (Proj. P. ErrL. nWarning) + 1)	'EC2007 - 16 bit immediate constant is out of range.
'Set new element	'EC2008, EC2009, EC2010, EC2011, EC2013, EC2014,
Proj.P.ErrL.lWarning(UBound(Proj.P.ErrL.lWarning)) = LineNum	'EC2015, EC2016, EC2017, EC2018, EC2019
Proj.P.ErrL.sWarning(UBound(Proj.P.ErrL.sWarning)) = Message	
<pre>Proj.P.ErrL.nWarning(UBound(Proj.P.ErrL.nWarning)) = WngCode</pre>	'EC3001, EC3002
End Sub	

1002, EC1003, EC1004 nd takes X or erands, not Y. Offending operand: 'Y' immediate constant is out of ra EC2010, EC2011, EC2013, EC2014, EC2018 EC3002

22.5. pExec

Option Explicit

·		- 1
' Public declarat	ions in this module:	
1		
' VARIABLES:		
' CPU - current	CPII state	
, CIO CUIICIIC	cro beace	
CONSTANTS:		
	nstants for CPU.FLAGS	
, IIA - IIAY CO.	IIStants for CPU.FLAGS	
		-
' PROCEDURES:		
' exeInit - ini	tialises this module	1
' Tick - advance	e simulation by 1 tick	1
' DI2Str - CStr	(microinstruction)	
		_ •
'Decoded microins	truction	
Private Type TpDI		
	up has no built in sur	
-	'VB has no built-in sup	-
Sig2 As Long	' 64-bit integers, so	have to s

egers, so have to split them Sig2 As Long nToIDB As Long nAluOpNum As Integer

nJmpCond As Integer nAdrMul As Integer nAluSh As Integer nIntIS As Integer End Type



'General purpose A As Long R(0 To 3) As Long

'Special purpose IP As Long SP As Long FLAGS As Long 'Internal MAR As Long MDR As Long CIB As String

Fetch As Boolean FREM As Integer	Private JmpOpType4Hex(0 To 11) As Integer
fremMem As Boolean 'invisible register to allow for FREM+ IS As Long	Public Sub exeInit()
'' ' Execution'	' Initializes this module '
··	Public Sub exeInit()
eSelectedReg As Integer '0X: b,c,d,e,sp,ip	Dim i As Integer
eIDB As Long eIAB As Long	'Reset arrays For i = 0 To 255
eDP As Integer 'which microinstruction doing	InstructionLen $(i) = 1$ 'let it fetch and try to de
eEDB As Long eEAB As Long	<pre>InstructionMem(i) = False Next</pre>
Breakpoint() As Long 'address at which we are when first exec	
nd Type	'Cannot do it in a loop - there are holes everywher
Signal numbers	InstructionLen(&H0) = 1 $InstructionLen(&H1) = 1$
rivate Const reg_sb = 0	InstructionLen(&H2) = 1
rivate Const reg_sc = 1	InstructionLen(&H3) = 1
rivate Const reg_sd = 2 rivate Const reg_se = 3	InstructionLen(&H4) = 1 $InstructionLen(&H5) = 1$
rivate Const reg_ssp = 4	<pre>InstructionLen(&H6) = 1</pre>
rivate Const reg_sip = 5 rivate Const reg_r = 6	InstructionLen(&H7) = 1
rivate Const reg $w = 7$	InstructionLen(&HF) = 1 InstructionLen($\&H10$) = 1
rivate Const reg_ipi = 8	InstructionLen(&H11) = 1
rivate Const reg_spi = 9	InstructionLen(&H13) = 1
rivate Const reg_spd = 10 rivate Const adr_br = 11	InstructionLen($\&$ H14) = 1 InstructionLen($\&$ H15) = 1
rivate Const adr_sr = 12	InstructionLen(&H16) = 1
rivate Const adr_im = 13	InstructionLen(&H17) = 1
rivate Const adr_c = 14 rivate Const lea_ad = 15	InstructionLen($\&$ H54) = 1 InstructionLen($\&$ H55) = 1
rivate Const acc_r = 16	InstructionLen(&H56) = 1
rivate Const acc_w = 17	InstructionLen(&H57) = 1
rivate Const flg_r = 18 rivate Const flg_w = 19	InstructionLen($\&H64$) = 1 InstructionLen($\&H65$) = 1
rivate Const ctl_mr = 20	InstructionLen(&H66) = 1
rivate Const ctl_mw = 21	InstructionLen($\&$ H67) = 1
rivate Const ctl_pr = 22 rivate Const ctl_pw = 23	InstructionLen(&H72) = 1 InstructionLen(&H73) = 1
rivate Const mdr_ri = 24	<pre>InstructionLen(&H75) = 1</pre>
rivate Const mdr_re = 25 rivate Const mdr_wi = 26	InstructionLen($\&H76$) = 1
rivate Const mdr_wi = 20	InstructionLen($\&$ H77) = 1 InstructionLen($\&$ H8F) = 1
rivate Const mar_ri = 28	InstructionLen(&H9F) = 1
rivate Const mar_re = 29 rivate Const mar_wi = 30 'skip 31 because there may be errors	InstructionLen($\&$ HAO) = 1 InstructionLen($\&$ HAI) = 1
rivate Const mar_we = 32 ' associated with sign	InstructionLen($\&HA2$) = 1
rivate Const alu_swp = 33	InstructionLen(&HA3) = 1
rivate Const ctl_halt = 34 rivate Const lea_da = 35	InstructionLen(&HA4) = 1 InstructionLen(&HA5) = 1
rivate Const flg_stz = 36	InstructionLen(&HA6) = 1
rivate Const flg_clz = 37	InstructionLen(&HA7) = 1
rivate Const flg_stc = 38 rivate Const flg_clc = 39	InstructionLen($\&$ HA8) = 1 InstructionLen($\&$ HA9) = 1
rivate Const flg_sto = 40	InstructionLen(&HAA) = 1
rivate Const flg_clo = 41	InstructionLen(&HAB) = 1
rivate Const flg_sts = 42 rivate Const flg_cls = 43	InstructionLen($\&$ HAC) = 1 InstructionLen($\&$ HAD) = 1
rivate Const flg_sti = 44	InstructionLen(&HAE) = 1
rivate Const flg_cli = 45	InstructionLen(&HAF) = 1
civate Const op_alu_sh = 58 'these indicate that the	InstructionLen($\&$ HBC) = 1 InstructionLen($\&$ HBD) = 1
rivate Const op_jmp_cond = 59 'next operand in a call to	InstructionLen(&HBE) = 1
rivate Const op_idb_im = 60 'SFlg is a respective parameter	
rivate Const op_adr_mm = 61 rivate Const op_alu_c = 62	InstructionLen(&HDC) = 1 InstructionLen(&HDD) = 1
	<pre>InstructionLen(&HDE) = 1</pre>
lag constants	InstructionLen(&HDF) = 1
ublic Const flZ = 1 ublic Const flS = 2	InstructionLen($\&$ HE6) = 1 InstructionLen($\&$ HE7) = 1
ublic Const fl0 = 4	<pre>InstructionLen(&HF0) = 1</pre>
ublic Const flC = 8	InstructionLen(&HF1) = 1
ublic Const flI = 16 ublic Const flN = 256	InstructionLen($\&$ HF2) = 1 InstructionLen($\&$ HF3) = 1
ublic Const flP = 512	<pre>InstructionLen(&HF4) = 1</pre>
Potch simplification arrayo	InstructionLen(&HF5) = 1
Fetch simplification_arrays ublic InstructionLen(0 To 255) As Integer 'public because othe:	InstructionLen(&HF6) = 1 r InstructionLen(&HF7) = 1
ublic InstructionMem(0 To 255) As Boolean 'units may need the	-
engths	InstructionLen($\&$ H25) = 2 InstructionLen($\&$ H25) = 2
<pre>"Massive decode" simplification rivate AluOpTypelHex(0 To 7) As Integer</pre>	InstructionLen($\&$ H35) = 2 InstructionLen($\&$ H74) = 2
rivate AluOpType2Hex(0 To 13) As Integer	InstructionLen(&H80) = 2
rivate AluOpType3Hex1(0 To 4) As Integer	<pre>InstructionLen(&H83) = 2</pre>

<pre>InstructionLen(&H86) = 2</pre>	
InstructionLen($\&H89$) = 2	
InstructionLen($\&H8C$) = 2	
InstructionLen(&H90) = 2	
InstructionLen($\&H93$) = 2	
InstructionLen(&H96) = 2	
InstructionLen(&H99) = 2	
InstructionLen(&H9C) = 2	
InstructionLen($\&HB0$) = 2	
InstructionLen($\&$ HB3) = 2	
InstructionLen(&HB6) = 2	
InstructionLen(&HB9) = 2	
<pre>InstructionLen(&HCO) = 2</pre>	
InstructionLen(&HC1) = 2	
InstructionLen(&HC2) = 2	
InstructionLen(&HC3) = 2	
InstructionLen(&HC4) = 2	
InstructionLen(&HC5) = 2	
InstructionLen(&HC6) = 2	
InstructionLen(&HC7) = 2	
<pre>InstructionLen(&HD0) = 2</pre>	
<pre>InstructionLen(&HD1) = 2</pre>	
InstructionLen(&HD2) = 2	
<pre>InstructionLen(&HD3) = 2</pre>	
InstructionLen(&HD4) = 2	
InstructionLen(&HD6) = 2	
InstructionLen(&HD8) = 2	
<pre>InstructionLen(&HD9) = 2</pre>	
InstructionLen(&HDA) = 2	
<pre>InstructionLen(&HDB) = 2</pre>	
<pre>InstructionLen(&HE4) = 2</pre>	
InstructionLen(&HE5) = 2	
InstructionLen(&H12) = 3	
InstructionLen(&H20) = 3	
InstructionLen(&H21) = 3	
InstructionLen(&H22) = 3	
InstructionLen(&H23) = 3	
InstructionLen($\&$ H24) = 3	
<pre>InstructionLen(&H30) = 3</pre>	
InstructionLen($\&$ H31) = 3	
InstructionLen(&H32) = 3	
InstructionLen(&H33) = 3	
<pre>InstructionLen(&H34) = 3 InstructionLen(&H81) = 3</pre>	
<pre>InstructionLen(&H81) = 3 InstructionLen(&H84) = 3</pre>	
<pre>InstructionLen(&H87) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H91) = 3</pre>	
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<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H91) = 3 InstructionLen(&H91) = 3 InstructionLen(&H97) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H84) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H91) = 3 InstructionLen(&H97) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9D) = 3 InstructionLen(&H9D) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H91) = 3 InstructionLen(&H94) = 3 InstructionLen(&H94) = 3 InstructionLen(&H97) = 3 InstructionLen(&H91) = 3 InstructionLen(&H81) = 3 InstructionLen(&H81) = 3 InstructionLen(&H84) = 3 InstructionLen(&H84) = 3</pre>	
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<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H91) = 3 InstructionLen(&H91) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9D) = 3 InstructionLen(&HB1) = 3 InstructionLen(&HB4) = 3 InstructionLen(&HB7) = 3 InstructionLen(&HBA) = 3 InstructionLen(&HBA) = 3 InstructionLen(&HBA) = 3 InstructionLen(&HBA) = 3 InstructionLen(&HBA) = 3 InstructionLen(&HBA) = 3 InstructionLen(&HEC) = 3 InstructionLen(&HE1) = 3 InstructionLen(&HE2) = 3 InstructionLen(&HE2) = 3 InstructionLen(&HE3) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H97) = 3 InstructionLen(&H97) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 3</pre>	
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 3</pre>	
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$eq:linear_line$	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H91) = 3 InstructionLen(&H94) = 3 InstructionLen(&H94) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8B) = 3 InstructionLen(&H8B) = 3 InstructionLen(&H8B) = -1 InstructionLen(&H8) = -1 InstructionLen(&HA) = -1</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = -1 InstructionLen(&HB) = -1 InstructionLen(&HB) = -1 InstructionLen(&HB) = -1</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = -1 InstructionLen(&HBA) = -1 InstructionLen(&HBA) = -1 InstructionLen(&HBA) = -1 InstructionLen(&HBA) = -1 InstructionLen(&HBAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 1 InstructionLen(&HAB) = -1 InstructionLen(&HB) = -1 InstructionLen(&HBA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 3 InstructionLen(&HED) = 3 InstructionLen(&HED) = 4 InstructionLen(&HBB) = -1 InstructionLen(&HB) = -1 InstructionLen(&HB) = -1 InstructionLen(&HB) = -1 InstructionLen(&HB) = -1 InstructionLen(&HAD) = -1 InstructionLen(&HAD) = -1 InstructionLen(&HAD) = -1</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 1 InstructionLen(&HBA) = -1 InstructionLen(&HA) = -1 InstructionLen(&HA) = -1 InstructionLen(&HA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 1 InstructionLen(&HBA) = -1 InstructionLen(&HAB) = -1 InstructionLen(&HAD) = -1 InstructionL</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = -1 InstructionLen(&HB) = -1 InstructionLen(&HB) = -1 InstructionLen(&HB) = -1 InstructionLen(&HBA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9D) = 3 InstructionLen(&HBA) = 3 InstructionLen(&HEA) = 3 InstructionLen(&HEA) = 3 InstructionLen(&HEA) = 3 InstructionLen(&HEA) = 3 InstructionLen(&HEA) = 3 InstructionLen(&HEA) = 1 InstructionLen(&HAA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA</pre>	'minus saves code lines - see below
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<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H9J) = 3 InstructionLen(&H9J) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 1 InstructionLen(&HBA) = -1 InstructionLen(&HBA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA</pre>	'minus saves code lines - see below
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<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H91) = 3 InstructionLen(&H91) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 1 InstructionLen(&HBA) = 1 InstructionLen(&HAB) = -1 InstructionLen(&HAB) = -1 InstructionLen(&HAB) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H91) = 3 InstructionLen(&H91) = 3 InstructionLen(&H94) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HB4) = 3 InstructionLen(&HB4) = 3 InstructionLen(&HB4) = 3 InstructionLen(&HB7) = 4 InstructionLen(&H88) = -1 InstructionLen(&H89) = -1 InstructionLen(&H80) = -1 InstructionL</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HBA) = 1 InstructionLen(&HBA) = 1 InstructionLen(&HBA) = -1 InstructionLen(&HBA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAA) = -1 InstructionLen(&HAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA</pre>	'minus saves code lines - see below
<pre>InstructionLen(&H87) = 3 InstructionLen(&H8A) = 3 InstructionLen(&H8D) = 3 InstructionLen(&H91) = 3 InstructionLen(&H91) = 3 InstructionLen(&H94) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&H9A) = 3 InstructionLen(&HB4) = 3 InstructionLen(&HB4) = 3 InstructionLen(&HB4) = 3 InstructionLen(&HB7) = 4 InstructionLen(&H88) = -1 InstructionLen(&H89) = -1 InstructionLen(&H80) = -1 InstructionL</pre>	'minus saves code lines - see below

InstructionLen(&H8B) = -1	
InstructionLen(&H8E) = -1	
InstructionLen(&H92) = -1	
InstructionLen(&H95) = -1	
InstructionLen(&H98) = -1	
InstructionLen(&H9B) = -1	
<pre>InstructionLen(&H9E) = -1</pre>	
InstructionLen(&HB2) = -1	
InstructionLen(&HB5) = -1	
InstructionLen(&HB8) = -1	
<pre>InstructionLen(&HBB) = -1</pre>	
InstructionLen(&H26) = -2	
InstructionLen(&H27) = -2	
InstructionLen($\&H36$) = -2	
InstructionLen(&H37) = -2	
For $i = 0$ To 255	
<pre>If InstructionLen(i) = -2 Or InstructionLen(i) = -1 Then</pre>	
InstructionMem(i) = True	
<pre>InstructionLen(i) = -InstructionLen(i)</pre>	
End If	
Next	
'Initialise array AluOpTypelHex	
AluOpTypelHex(0) = &HCO	
AluOpTypelHex(1) = &HC1	
AluOpTypelHex(2) = &HC2	
AluOpTypelHex(3) = $\&$ HC3	
AluOpTypelHex(4) = &HC4	
AluOpTypelHex(5) = &HC5	
AluOpTypelHex(6) = &HC6	
AluOpType1Hex(7) = &HC7	
'Initialise array AluOpType2Hex	
AluOpType2Hex(0) = $\&H80$	
AluOpType2Hex(1) = &H83	
AluOpType2Hex(2) = &H86	
AluOpType2Hex(3) = &H89	
AluOpType2Hex(4) = &H8C	
AluOpType2Hex(5) = &H90	
AluOpType2Hex(6) = &H93	
AluOpType2Hex $(7) = \&H96$	
AluOpType2Hex(8) = &H99	
AluOpType2Hex(9) = &H9C	
AluOpType2Hex(10) = &HB0	
AluOpType2Hex(11) = &HB3	
AluOpType2Hex(12) = &HB6	
AluOpType2Hex(13) = &HB9	
AluOpType2Hex(13) = &HB9 'Initialise array AluOpType3Hex	
AluOpType2Hex(13) = &HB9 Initialise array AluOpType3Hex AluOpType3Hex1(0) = &HA0	
AluOpType2Hex(13) = &HB9 'Initialise array AluOpType3Hex	
AluOpType2Hex(13) = &HB9 Initialise array AluOpType3Hex AluOpType3Hex1(0) = &HA0 AluOpType3Hex1(1) = &HA4	
AluOpType2Hex(13) = &HB9 Initialise array AluOpType3Hex AluOpType3Hex1(0) = &HA0 AluOpType3Hex1(1) = &HA4 AluOpType3Hex1(2) = &HA8	
AluOpType2Hex(13) = &HB9 Initialise array AluOpType3Hex AluOpType3Hex1(0) = &HA0 AluOpType3Hex1(1) = &HA4 AluOpType3Hex1(2) = &HA8 AluOpType3Hex1(3) = &HBC	
AluOpType2Hex(13) = &HB9 Initialise array AluOpType3Hex AluOpType3Hex1(0) = &HA0 AluOpType3Hex1(1) = &HA4 AluOpType3Hex1(2) = &HA8 AluOpType3Hex1(3) = &HBC AluOpType3Hex1(4) = &HDC	
AluOpType2Hex(13) = &HB9 Initialise array AluOpType3Hex AluOpType3Hex1(0) = &HA0 AluOpType3Hex1(1) = &HA4 AluOpType3Hex1(2) = &HA4 AluOpType3Hex1(3) = &HBC AluOpType3Hex1(4) = &HDC AluOpType3Hex2(0) = &HAC	
Alu0pType2Hex(13) = &HB9 Thitalise array Alu0pType3Hex Alu0pType3Hex1(0) = &HA0 Alu0pType3Hex1(1) = &HA4 Alu0pType3Hex1(2) = &HA8 Alu0pType3Hex1(3) = &HBC Alu0pType3Hex1(4) = &HBC Alu0pType3Hex2(0) = &HAC Alu0pType3Hex2(1) = &HAD	
AluOpType2Hex(13) = &HB9 Initialise array AluOpType3Hex AluOpType3Hex1(0) = &HA0 AluOpType3Hex1(1) = &HA4 AluOpType3Hex1(2) = &HA8 AluOpType3Hex1(3) = &HBC AluOpType3Hex2(4) = &HDC AluOpType3Hex2(0) = &HAC AluOpType3Hex2(1) = &HAD AluOpType3Hex2(2) = &HAE	
Alu0pType2Hex(13) = &HB9 Thitalise array Alu0pType3Hex Alu0pType3Hex1(0) = &HA0 Alu0pType3Hex1(1) = &HA4 Alu0pType3Hex1(2) = &HA8 Alu0pType3Hex1(3) = &HBC Alu0pType3Hex1(4) = &HBC Alu0pType3Hex2(0) = &HAC Alu0pType3Hex2(1) = &HAD	
Alu0pType2Hex(13) = &HB9 Initialise array Alu0pType3Hex Alu0pType3Hex1(0) = &HA0 Alu0pType3Hex1(1) = &HA4 Alu0pType3Hex1(2) = &HA4 Alu0pType3Hex1(2) = &HAC Alu0pType3Hex2(0) = &HAC Alu0pType3Hex2(1) = &HAD Alu0pType3Hex2(2) = &HAE Alu0pType3Hex2(3) = &HAF Alu0pType3Hex2(4) = &HAF	
Alu0pType2Hex(13) = &HB9 Initialise array Alu0pType3Hex Alu0pType3Hex1(0) = &HA0 Alu0pType3Hex1(1) = &HA4 Alu0pType3Hex1(2) = &HA4 Alu0pType3Hex1(2) = &HAC Alu0pType3Hex2(0) = &HAC Alu0pType3Hex2(1) = &HAD Alu0pType3Hex2(2) = &HAE Alu0pType3Hex2(3) = &HAF Alu0pType3Hex2(4) = &HAF	
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Call Errr("pExec.GFlg: index should be between 0 and 62 excluding 31.")

Private Sub SFlg(di As TpDI, ParamArray arr()) ' DESCRIPTION: adds the specified signals to the ' decoded instruction buffer element.	Private Sub eTick() If Proj.CPU.Fetch Then
DESCRIPTION: adds the specified signals to the '	-
decoded instruction buffer element.	eFetch
	Else
PARAMETERS:	eExecute End If
di - decoded microinstruction to be modigied '	devTick
arr() - list of signal constants	Proj.TickCount = Proj.TickCount + 1
	End Sub
NOTES: using a signal constant beginning with '	
op_ will make SFlg assume that the next	11
arr() element is the signal-specific param '	' Private Sub eFetch()
to be stored in di.	' Advances simulation by one tick when CPU is '
.vate Sub SFlg(di As TpDI, ParamArray arr())	' in fetch mode. Does NOT check if it is. After
Dim i As Integer	' fetching the last byte automatically initiates '
li.nAdrMul = 1	' eDecode, so that decoding takes no further ticks '
li.nAluOpNum = -1	··································
li.nJmpCond = -1	Private Sub eFetch()
i.nToIDB = -1	Dim b As Byte
li.nAluSh = 0	With Proj.CPU
i.Sig1 = 0 $i.Sig2 = 0$	I Datab a but o into a temporary location
i.Sig2 = 0 or i = 0 To UBound(arr)	'Fetch a byte into a temporary location b = Proj.RAM(.IP)
'Add signal	.IP = .IP + 1
If $(arr(i) \ge 0)$ And $(arr(i) \le 30)$ Then	
di.Sig1 = di.Sig1 Or (2 ^ arr(i))	'First loop of fetch cycle
ElseIf $(arr(i) \ge 32)$ And $(arr(i) \le 62)$ Then	If .CIB = "" Then
di.Sig2 = di.Sig2 Or (2 ^ (arr(i) - 32))	'How many bytes to fetch
End If 'Check if signal parameter present	.FREM = InstructionLen(b) .fremMem = InstructionMem(b)
If arr(i) = op_adr_mm Then	.iremment = instructionment(b)
di.nAdrMul = arr(i + 1)	'Middle loop & recalc of FREM needed
i = i + 1	ElseIf (.FREM = 0) And .fremMem Then
<pre>ElseIf arr(i) = op_alu_c Then</pre>	If (b And 192) = 0 Then .FREM = 3
di.nAluOpNum = arr(i + 1)	If (b And 192) = 64 Then .FREM = 1
i = i + 1	If (b And 192) = 0 Then .FREM = 3
ElseIf arr(i) = op_idb_im Then	If (b And 160) = 160 Then .FREM = 3
di.nToIDB = arr(i + 1) i = i + 1	<pre>If (b And 160) = 128 Then .FREM = 1 .fremMem = False</pre>
ElseIf arr(i) = op_jmp_cond Then	End If
di.nJmpCond = $arr(i + 1)$	
i = i + 1	'Add fetched byte to CIB
<pre>ElseIf arr(i) = op_alu_sh Then</pre>	.CIB = .CIB + Chr(b)
di.nAluSh = arr(i + 1)	'Update remaining bytes register
i = i + 1	.FREM = .FREM - 1
End If Mext	Itest less of fotab avala
l Sub	'Last <u>loop of fetch cycle</u> Dim i As Integer
	If .FREM = 0 And .fremMem = False Then
	.Fetch = False
Public Sub Tick()	eDecode
	'Check breakpoints
dvances simulation by one clock tick by	For i = 0 To UBound(.Breakpoint)
etching, decoding or executing something. '	If .IP = .Breakpoint(i) Then
lic Sub Tick()	Proj.Paused = <mark>True</mark> C all fiMain.UpdateAll(True)
Tick	End If
No updates - update manually for more control	Next
Sub	End If
	End With
ublic Sub Tick()	End Sub
dvances simulation by CPU instruction,	11
etching the next one.	' Private Sub eDecode '
'	
lic Sub Step()	' DESCRIPTION: Decodes fetched instruction '
<pre>im wasIS As Integer 'finish step at interrupt</pre>	' from CPU.CIB, placing decoded microprogram '
hile Not Proj.CPU.Fetch	' into CPU.DIB. '
wasIS = Proj.CPU.IS	Private Cub oDecedo()
eTick If wasIS <> Proj.CPU.IS Then Exit Sub	Private Sub eDecode() With Proj.CPU
IF Wasis <> Proj.CPU.IS Then Exit Sub	MIGH FLUJ.CFU
Mile Proj.CPU.Fetch	'Get first byte into b
-	If .CIB = "" Then
Wabib - riul.Cru.ib	Call Errr("pExec.execDecode: CIB is empty. Contact
wasIS = Proj.CPU.IS eTick	
	author.")
eTick If wasIS <> Proj.CPU.IS Then Exit Sub Nend	End If
eTick If wasIS <> Proj.CPU.IS Then Exit Sub Mend No updates - update manually for more control	End If Dim b As Byte
eTick If wasIS <> Proj.CPU.IS Then Exit Sub Mend	End If
eTick If wasIS <> Proj.CPU.IS Then Exit Sub end No updates - update manually for more control	End If Dim b As Byte

If $(b \ge \&H30)$ And $(b \le \&H37)$ Then b = b - &H10If (b >= &H20) And (b <= &H23) Then 'ld Reg,I ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), reg_sX(b And 3), reg_w, op_idb_im, Chr2Dec(Mid(.CIB, 2, 2))) Exit Sub ElseIf b = &H24 Ther <u>'ld A,I</u> ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), acc_w, op_idb_im, Chr2Dec(Mid(.CIB, 2, 2))) Exit Sub ElseIf b = &H25 Then 'ld R,R b = Asc(Mid(.CIB, 2, 1))If (b And 192) = 192 Then 'ld A,A ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), acc_r, acc_w) ElseIf (b And 192) = 128 Then 'ld A,Rn ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), reg_sX(b And 3), reg_r, acc_w) ElseIf (b And 192) = 64 Then 'ld Rn,A ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), reg_sX((b And 24) / 8), reg_w, acc_r) Else <u>'ld Rn,Rn</u> ReDim .DIB(-1 To 1) Call SFlg(.DIB(0), reg_sX(b And 3), reg_r, mdr_wi) Call SFlg(.DIB(1), reg_sX((b And 24) / 8), reg_w, mdr_ri) End If Exit Sub ElseIf b = &H26 Then 'ld R,M ReDim .DIB(-1 To -1) Call DecodeMem(3) b = Asc(Mid(.CIB, 2, 1))**ReDim** Preserve .DIB(-1 To UBound(.DIB) + 2) Call SFlg(.DIB(UBound(.DIB) - 1), mar_re, ctl_mr, mdr_we) If (b And 4) = 4 Then Call SFlg(.DIB(UBound(.DIB)), mdr_ri, acc_w) Else Call SFlg(.DIB(UBound(.DIB)), mdr_ri, reg_sX(b And 3), reg_w) End If Exit Sub ElseIf b = &H27 Then 'ld M,R ReDim .DIB(-1 To -1) Call DecodeMem(3) b = Asc(Mid(.CIB, 2, 1))**ReDim** Preserve .DIB(-1 To UBound(.DIB) + 2) If (b And 4) = 4 Then Call SFlg(.DIB(UBound(.DIB) - 1), acc_r, mdr_wi) Else Call SFlg(.DIB(UBound(.DIB) - 1), reg_sX(b And 3), reg_r, mdr wi) End If Call SFlg(.DIB(UBound(.DIB)), mar_re, mdr_re, ctl_mw) Exit Sub End If TYPE 1 Dim i As Integer For i = 0 To 7 If (b = AluOpTypelHex(i)) Then b = Asc(Mid(.CIB, 2, 1))If (b And 128) = 0 Then 'op A,Rn ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), reg_sX((b And 96) $\ 32$), reg_r, op_alu_c, 19 + i, acc_w) Else ReDim .DIB(-1 To 0) If (b And 16) = 0 Then 'op Rn,N Call SFlg(.DIB(0), reg_sX((b And 96) \ 32), reg_r, op_alu_sh, b And 15, op_alu_c, 27 + i, reg_w) Else 'op A,N

Call SFlg(.DIB(0), acc_r, op_alu_sh, b And 15, op_alu_c, 27 + i, acc_w) End If End If Exit Sub End If Next - TYPE 2 For i = 0 To 13 If (b = AluOpType2Hex(i)) Then b = Asc(Mid(.CIB, 2, 1))If (b And 4) = 0 Then If (b And 8) = 0 Then 'op A,Rn ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), reg_sX(b And 3), reg_r, acc_w, op_alu_c, i) Else 'op Rn,A ReDim .DIB(-1 To 1) Call SFlg(.DIB(0), reg_sX(b And 3), reg_r, op_alu_c, i, alu_swp) Call SFlg(.DIB(1), reg_sX(b And 3), reg_w) End If Else 'op A,A ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), acc_r, acc_w, op_alu_c, i) End If Exit Sub ElseIf (b = AluOpType2Hex(i) + 1) Then 'op A,I ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), acc_w, op_alu_c, i, op_idb_im, Chr2Dec(Mid(,CIB, 2, 2))) Exit Sub ElseIf (b = AluOpType2Hex(i) + 2) Then 'op A,M ReDim .DIB(-1 To -1) Call DecodeMem(2) **ReDim** Preserve .DIB(-1 To UBound(.DIB) + 2) Call SFlg(.DIB(UBound(.DIB) - 1), mar_re, ctl_mr, mdr_we) Call SFlg(.DIB(UBound(.DIB)), mdr_ri, acc_w, op_alu_c, i) Exit Sub End If Next TYPE 3 For i = 0 To 4 If (b = AluOpType3Hex2(i)) Then 'op A ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), acc_r, op_alu_c, 14 + i, acc_w) Exit Sub ElseIf (b >= AluOpType3Hex1(i)) And (b <= AluOpType3Hex1(i) +</pre> 3) Then 'op Rn ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), reg_sX(b And 3), reg_r, op_alu_c, 14 + i, reg_w) Exit Sub End If Next TYPE 4 For i = 0 To 11 ijxx M
If b = JmpOpType4Hex(i) Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), op_jmp_cond, i) Call DecodeMem(2) **ReDim** Preserve .DIB(-1 To UBound(.DIB) + 1) Call SFlg(.DIB(UBound(.DIB)), mar_ri, lea_ad, reg_sip, reg_w) Exit Sub End If Next If b = & H70 Then ReDim .DIB(-1 To -1)

Call DecodeMem(2) **ReDim** Preserve .DIB(-1 To UBound(.DIB) + 1) Call SFlg(.DIB(UBound(.DIB)), mar_ri, lea_ad, reg_sip, reg_w) Exit Sub End If CALL --If b = &H71 Then ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), reg_sip, reg_r, mdr_wi, reg_spi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_mw) Call DecodeMem(2) $\label{eq:reduced_relation} \textbf{ReDim} \ \texttt{Preserve} \ . \texttt{DIB}(\ -1 \ \textbf{To} \ \textbf{UBound}(\ . \texttt{DIB}) \ + \ \textbf{1})$ Call SFlg(.DIB(UBound(.DIB)), mar_ri, lea_ad, reg_sip, reg_w) Exit Sub End If '--- RET ---' If b = &H72 Then ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg.spd, reg.ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), mar_re, ctl_mr, mdr_we) Call SFlg(.DIB(2), reg_sip, mdr_ri, reg_w) Exit Sub End If - HALT --If b = &H75 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), ctl_halt) Exit Sub End If '--- PUSH --- ' If (b >= &H0) And (b <= &H3) Then 'push Rn ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), reg_sX(b), reg_r, mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_mw, reg_spi) Exit Sub ElseIf b = &H10 Then 'push A ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_ssp, adr_sr, adr_c, mar_wi) **Call** SFlg(.DIB(1), acc_r, mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_mw, reg_spi) Exit Sub ElseIf b = &H12 Then 'push I
ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), op_idb_im, Chr2Dec(Mid(.CIB, 2, 2)), mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_mw, reg_spi) Exit Sub ElseIf b = &H13 Then 'pushpc ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), reg_sip, reg_r, mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_mw, reg_spi) Exit Sub ElseIf b = &H14 Then 'pushsp ReDim DIB(-1 To 2) Call SFlg(.DIB(0), reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), reg_ssp, ddr_sr, ddr_e, mdr_wi) Call SFlg(.DIB(1), reg_ssp, reg_r, mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_mw, reg_spi) Exit Sub ElseIf b = &H15 Then 'pushfl ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), flg_r, mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_mw, reg_spi) Exit Sub End If $\frac{1}{1} - \frac{1}{2} \frac{$ 'pop Rn ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_spd, reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), mar_re, ctl_mr, mdr_we) Call SFlg(.DIB(2), reg_sX(b And 3), mdr_ri, reg_w) Exit Sub

ElseIf b = &H11 Then

'pop A ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_spd, reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), mar_re, ctl_mr, mdr_we) **Call** SFlg(.DIB(2), mdr_ri, acc_w) Exit Sub ElseIf b = &H16 Then 'popsp ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_spd, reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), mar_re, ctl_mr, mdr_we) Call SFlg(.DIB(2), reg_ssp, mdr_ri, reg_w) Exit Sub ElseIf b = &H17 Then 'popfl
ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_spd, reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), mar_re, ctl_mr, mdr_we) Call SFlg(.DIB(2), mdr_ri, flg_w) Exit Sub End If '--- SP2B ---' If b = &HF Then ReDim .DIB(-1 To 1) Call SFlg(.DIB(0), reg_ssp, reg_r, mdr_wi) Call SFlg(.DIB(1), mdr_ri, reg_sb, reg_w) Exit Sub End If LEA ---' If $(b \ge \&H8)$ And $(b \le \&HB)$ Then ReDim .DIB(-1 To -1) Call DecodeMem(2) ReDim Preserve .DIB(-1 To UBound(.DIB) + 1) Call SFlg(.DIB(UBound(.DIB)), mar_ri, lea_ad, reg_sX(b And 3), reg_w) Exit Sub ElseIf b = &H2F Then ReDim .DIB(-1 To -1) Call DecodeMem(2) **ReDim** Preserve .DIB(-1 To UBound(.DIB) + 1) Call SFlg(.DIB(UBound(.DIB)), mar_ri, lea_ad, acc_w) Exit Sub End If '--- XCHG ---' If (b >= &HF0) And (b <= &HF3) Then 'xchg A,Rn ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), acc_r, mdr_wi) Call SFlg(.DIB(1), reg_sX(b And 3), reg_r, acc_w) Call SFlg(.DIB(2), mdr_ri, reg_sX(b And 3), reg_w) Exit Sub ElseIf b = HE6 Then 'xchg c,d
ReDim .DIB(-1 To 3) Call SFlg(.DIB(0), reg_sc, reg_r, mdr_wi) Call SFlg(.DIB(1), reg_sd, reg_r, lea_da, mar_wi) Call SFlg(.DIB(2), mdr_ri, reg_sd, reg_w) Call SFlg(.DIB(3), mar_ri, lea_ad, reg_sc, reg_w) Exit Sub ElseIf b = &HE7 Then 'xchg c,e ReDim .DIB(-1 To 3) Call SFlg(.DIB(0), reg_sc, reg_r, mdr_wi) Call SFlg(.DIB(1), reg_se, reg_r, lea_da, mar_wi) Call SFlg(.DIB(2), mdr_ri, reg_se, reg_w) Call SFlg(.DIB(3), mar_ri, lea_ad, reg_sc, reg_w) Exit Sub ElseIf b = &HF4 Then 'xchg d,e ReDim .DIB(-1 To 3) Call SFlg(.DIB(0), reg_sd, reg_r, mdr_wi) Call SFlg(.DIB(1), reg_se, reg_r, lea_da, mar_wi) Call SFlg(.DIB(2), mdr_ri, reg_se, reg_w) Call SFlg(.DIB(3), mar_ri, lea_ad, reg_sd, reg_w) Exit Sub ElseIf b = &HF5 Then 'xchq b,c ReDim .DIB(-1 To 3) Call SFlg(.DIB(0), reg_sb, reg_r, mdr_wi) Call SFlg(.DIB(1), reg_sc, reg_r, lea_da, mar_wi) Call SFlg(.DIB(2), mdr_ri, reg_sc, reg_w) Call SFlg(.DIB(3), mar_ri, lea_ad, reg_sb, reg_w) Exit Sub ElseIf b = &HF6 Then 'xchq b.d

ReDim .DIB(-1 To 3) Call SFlg(.DIB(0), reg_sb, reg_r, mdr_wi) Call SFlg(.DIB(1), reg_sd, reg_r, lea_da, mar_wi) Call SFlg(.DIB(2), mdr_ri, reg_sd, reg_w) Call SFlg(.DIB(3), mar_ri, lea_ad, reg_sb, reg_w) Exit Sub ElseIf b = &HF7 Then 'xchg b,e ReDim .DIB(-1 To 3) Call SFlg(.DIB(0), reg_sb, reg_r, mdr_wi) Call SFlg(.DIB(1), reg_se, reg_r, lea_da, mar_wi) Call SFlg(.DIB(2), mdr_ri, reg_se, reg_w) Call SFlg(.DIB(3), mar_ri, lea_ad, reg_sb, reg_w) Exit Sub End If - NOP ---If b = &H8F Then ReDim .DIB(-1 To -1) Exit Sub End If If b = &H54 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_stz) Exit Sub ElseIf b = &H55 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_clz) Exit Sub ElseIf b = &H56 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_stc) Exit Sub ElseIf b = &H57 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_clc) Exit Sub ElseIf b = &H64 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_sto) Exit Sub ElseIf b = &H65 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_clo) Exit Sub ElseIf b = &H66 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_sts) Exit Sub ElseIf b = &H67 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_cls) Exit Sub ElseIf b = &H76 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_sti) Exit Sub ElseIf b = &H77 Then ReDim .DIB(-1 To 0) Call SFlg(.DIB(0), flg_cli) Exit Sub End If '--- IN ---' If b = &HE4 Then b = Asc(Mid(.CIB, 2, 1))ReDim .DIB(-1 To 2) $\mbox{Call SFlg(.DIB(0), IIf((b And 16) > 0, -1, reg_sX(b And 3)), } \\$ $IIf((b And 16) > 0, acc_r, reg_r), adr_im, adr_c, mar_wi)$ Call SFlg(.DIB(1), mar_re, ctl_pr, mdr_we) $\label{eq:call_sflg(.DIB(2), mdr_ri, IIf((b And 32) > 0, -1, reg_sX((b And 12) \ \ 32)), IIf((b And 32) > 0, acc_w, reg_w))$ Exit Sub **ElseIf** (b >= &HD8) And (b <= &HDB) Then ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), op_idb_im, Asc(Mid(.CIB, 2, 1)), adr_im, adr c, mar wi) Call SFlg(.DIB(1), mar_re, ctl_pr, mdr_we) Call SFlg(.DIB(2), mdr_ri, reg_sX(b And 3), reg_w) Exit Sub ElseIf b = &HE5 Then ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), op_idb_im, Asc(Mid(.CIB, 2, 1)), adr_im, adr_c, mar_wi) Call SFlg(.DIB(1), mar_re, ctl_pr, mdr_we) Call SFlg(.DIB(2), mdr_ri, acc_w)

End If '--- OUT ---' If b = &HD4 Then b = Asc(Mid(.CIB, 2, 1))ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), IIf((b And 32) > 0, -1, reg_sX((b And 12) $\$ 4)), IIf((b And 32) > 0, acc_r, reg_r), adr_im, adr_c, mar_wi) Call SFlg(.DIB(1), IIf((b And 16) > 0, -1, reg_sX(b And 3)), IIf((b And 16) > 0, acc_r, reg_r), mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_pw) Exit Sub ElseIf (b >= &HE0) And (b <= &HE3) Then</pre> ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), reg_sX(b And 3), reg_r, adr_im, adr_c, mar wi) Call SFlg(.DIB(1), op_idb_im, Chr2Dec(Mid(.CIB, 2, 2)), mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_pw) Exit Sub ElseIf (b >= &HD0) And (b <= &HD3) Then</pre> ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), op_idb_im, Asc(Mid(.CIB, 2, 1)), adr_im, adr_c, mar_wi) Call SFlg(.DIB(1), reg_sX(b And 3), reg_r, mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_pw) Exit Sub ElseIf b = &HD5 Then ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), acc_r, adr_im, adr_c, mar_wi) Call SFlg(.DIB(1), op_idb_im, Chr2Dec(Mid(.CIB, 2, 2)), mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_pw) Exit Sub ElseIf b = &HD6 Then ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), op_idb_im, Asc(Mid(.CIB, 2, 1)), adr_im, adr c, mar wi) Call SFlg(.DIB(1), acc_r, mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_pw) Exit Sub ElseIf b = &HD7 Then ReDim .DIB(-1 To 2) Call SFlg(.DIB(0), op_idb_im, Asc(Mid(.CIB, 2, 1)), adr_im, adr_c, mar_wi) Call SFlg(.DIB(1), op_idb_im, Chr2Dec(Mid(.CIB, 3, 2)), mdr_wi) Call SFlg(.DIB(2), mdr_re, mar_re, ctl_pw) Exit Sub End If If b = &H73 Then ReDim .DIB(-1 To 5) 'pop fl Call SFlg(.DIB(0), reg_spd, reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(1), mar_re, ctl_mr, mdr_we) Call SFlg(.DIB(2), mdr_ri, flg_w) Call SFlg(.DIB(3), reg_spd, reg_ssp, adr_sr, adr_c, mar_wi) Call SFlg(.DIB(4), mar_re, ctl_mr, mdr_we) Call SFlg(.DIB(5), reg_sip, mdr_ri, reg_w) Exit Sub

'Could not dec

Exit Sub

Call Errr("Failed to decode instruction starting with " + Dec2Hex(CInt(b), 2) + "h.")

End With End Sub

End If

Private Sub eExecute() Advances simulation by one tick when CPU is in execute mode. Does NOT check if it is. Executes current microinstruction, switches CPU to Fetch mode after executing last microinstruction.

Private Sub eExecute(With Proj.CPU

Dim adr_rslt As Long

'Check for no action
If UBound(.DIB) = -1 Then GoTo ItsOver

'Halt If GFlg(.DIB(.eDP), ctl_halt) Then Proj.Paused = True Proj.Running = True Proj.Halted = True fiMain.UpdateAll Call MsgBox("CPU halted", vbOKOnly + vbInformation) GoTo ItsOver End If 'Condition If GFlg(.DIB(.eDP), op_jmp_cond) Then zf=0 and sf=o If .DIB(.eDP).nJmpCond = 0 Then If ((.FLAGS And flz) = 0) And ((.FLAGS And flS) = 0) Then GoTo weContinue 'jle - zf=1 or sf<>of If .DIB(.eDP).nJmpCond = 1 Then If ((.FLAGS And fls) = fls) Or (((.FLAGS And flS) = flS) <> ((.FLAGS And flO) = flO)) Then GoTo weContinue If .DIB(.eDP).nJmpCond = 2 Then If ((.FLAGS And flS) = flS) <> ((.FLAGS And fl0) = fl0) Then GoTo weContinue 'iqe - sf=of If .DIB(.eDP).nJmpCond = 3 Then If ((.FLAGS And fls) = fls) = ((.FLAGS And fl0) = fl0) Then GoTo weContinue 'jz_- z<u>f=</u>1 If .DIB(.eDP).nJmpCond = 4 Then If ((.FLAGS And flz) = flz) Then GoTo weContinue 'jnz zf=0 If .DIB(.eDP).nJmpCond = 5 Then If ((.FLAGS And flZ) = 0) Then GoTo weContinue 'ic cf=1 If .DIB(.eDP).nJmpCond = 6 Then If ((.FLAGS And flC) = flC) Then GoTo weContinue 'jnc If .DIB(.eDP).nJmpCond = 7 Then If ((.FLAGS And flC) = 0) Then GoTo weContinue 'jo_- o<u>f</u>=1 65535 If .DIB(.eDP).nJmpCond = 8 Then If ((.FLAGS And flo) = flo) Then GoTo weContinue of=0 If .DIB(.eDP).nJmpCond = 9 Then If ((.FLAGS And flo) = 0) Then GoTo weContinue js - sf=1 If .DIB(.eDP).nJmpCond = 10 Then If ((.FLAGS And fls) = fls) Then GoTo weContinue 'ins - sf=0 true If .DIB(.eDP).nJmpCond = 11 Then If ((.FLAGS And fls) = 0) Then GoTo weContinue GoTo ItsOver weContinue: End If 'Decremennt SP If GFlg(.DIB(.eDP), reg_spd) Then .SP = .SP - 2 'Flags If GFlg(.DIB(.eDP), flg_stz) Then .FLAGS = .FLAGS Or flZ If GFlg(.DIB(.eDP), flg_stc) Then .FLAGS = .FLAGS Or flC If GFlg(.DIB(.eDP), flg_sto) Then .FLAGS = .FLAGS Or fl0 If GFlg(.DIB(.eDP), flg_sts) Then .FLAGS = .FLAGS Or flS If GFlg(.DIB(.eDP), flg_sti) Then .FLAGS = .FLAGS Or fll If GFlg(.DIB(.eDP), flg_clz) Then .FLAGS = .FLAGS And Not flZ If GFlg(.DIB(.eDP), flg_clc) Then .FLAGS = .FLAGS And Not flC If GFlg(.DIB(.eDP), flg_clo) Then .FLAGS = .FLAGS And Not fl0 If GFlg(.DIB(.eDP), flg_cls) Then .FLAGS = .FLAGS And Not flS If GFlg(.DIB(.eDP), flg_cli) Then .FLAGS = .FLAGS And Not fll 'Select register If GFlg(.DIB(.eDP), reg_sb) Then .eSelectedReg = 0 If GFlg(.DIB(.eDP), reg_sc) Then .eSelectedReg = 1 If GFlg(.DIB(.eDP), reg_sd) Then .eSelectedReg = 2 If GFlg(.DIB(.eDP), reg_se) Then .eSelectedReg = 3 If GFlg(.DIB(.eDP), reg_ssp) Then .eSelectedReg = 4 If GFlg(.DIB(.eDP), reg_sip) Then .eSelectedReg = 5 lead operations for ID If GFlg(.DIB(.eDP), reg_r) Then If .eSelectedReg = 4 Then .eIDB = .SP ElseIf .eSelectedReg = 5 Then .eIDB = .IP ElseIf .eSelectedReg >= 0 And .eSelectedReg <= 3 Then</pre> .eIDB = .R(.eSelectedReg) End If End If If GFlg(.DIB(.eDP), acc_r) Then .eIDB = .A If GFlg(.DIB(.eDP), flg_r) Then .eIDB = .FLAGS If GFlg(.DIB(.eDP), mdr_ri) Then .eIDB = .MDR

'Immediate If GFlg(.DIB(.eDP), op_idb_im) Then .eIDB = .DIB(.eDP).nToIDB If GFlg(.DIB(.eDP), adr_c) Then adr_rslt = 0 If GFlg(.DIB(.eDP), adr_sr) Then If .eSelectedReg = 4 Then adr_rslt = adr_rslt + .SP * .DIB(.eDP).nAdrMul ElseIf .eSelectedReg = 5 Then adr rslt = adr rslt + .IP * .DIB(.eDP).nAdrMul Else adr_rslt = adr_rslt + .R(.eSelectedReg) * .DIB(.eDP).nAdrMul End If End If If GFlg(.DIB(.eDP), adr_br) Then adr_rslt = adr_rslt + .R(0) If GFlg(.DIB(.eDP), adr_im) Then adr_rslt = adr_rslt + .eIDB .eIAB = adr rslt End If 'MAR If GFlg(.DIB(.eDP), mar_ri) Then .eIAB = .MAR 'External If GFlg(.DIB(.eDP), mar_re) Then .eEAB = .MAR If GFlg(.DIB(.eDP), mdr_re) Then .eEDB = .MDR If GFlg(.DIB(.eDP), lea_ad) Then .eIDB = .eIAB If GFlg(.DIB(.eDP), lea_da) Then .eIAB = .eIDB If GFlg(.DIB(.eDP), ctl_mr) Then .eEDB = Proj.RAM(.eEAB) * CLng(256) + Proj.RAM(.eEAB + 1) if GFlg(.DIB(.eDP), ctl_mw) Then Proj.RAM(.eEAB) = Int(.eEDB / 256) And &HFF: Proj.RAM(.eEAB + 1) = .eEDB And &HFF If GFlg(.DIB(.eDP), ctl_pr) Then .eEDB = devPortRead(.eEAB And If GFlg(.DIB(.eDP), ctl_pw) Then Call devPortWrite(.eEAB And 65535, .eEDB) 'ALU Dim 11 As Long, 12 As Long, 13 As Long Dim i As Integer, lf As Long Dim bl As Boolean, b2 As Boolean, b3 As Boolean Dim scf As Integer 'set carry flag: -1 leave, 0 0 set false. scf = -1If GFlg(.DIB(.eDP), op_alu_c) Then i = .DIB(.eDP).nAluOpNum If GFlg(.DIB(.eDP), alu_swp) Then ll = .eIDB 12 = .A Else 11 = .A 12 = .eIDB End If sult to be saved sult to be used for 13 = 11 lf = 11 b1 = (11 And 32768) > 0 is negative b2 = (12 And 32768) > 012 is n If (i >= 27) And (i <= 34) Then i = i - 8 l1 = .DIB(.eDP).nAluSh End If If i = 0 Then 'add 13 = 11 + 12 1f = 13 ElseIf i = 1 Then 'sub 13 = 11 - 12 lf = 13ElseIf i = 2 Then 'adc 13 = 11 + 12 + IIf((.FLAGS And flC) > 0, 1, 0) 1f = 13ElseIf i = 3 Then sbb 13 = 11 - 12 - IIf((.FLAGS And flC) > 0, 1, 0) lf = 13 ElseIf i = 4 Then 'cmp lf = 11 - 12ElseIf i = 5 Then 'mul 13 = 11 * 12 lf = 13 ElseIf i = 6 Then 'div 13 = 11 \ 12 1f = 13ElseIf i = 7 Then 'imul 13 = (Abs(11) And &H7FFF) * (Abs(12) And &H7FFF) **If** (b1 Xor b2) **Then** 13 = -13 lf = 13

ElseIf i = 8 Then 'idiv $13 = (13 \setminus 2) + IIf(scf = 1, 32768, 0)$ 13 = Int((Abs(11) And &H7FFF) / (Abs(12) And &H7FFF)) scf = IIf((lf And 1) = 0, 0, 1)If (b1 Xor b2) Then 13 = -13 Next 1f = 131f = 13ElseIf i = 9 Then 'mod End If 13 = 11 Mod 12 'Result to IDB .eIDB = (13 And 65535) lf = 13 ElseIf i = 10 Then 'and hate VB. 1 13 = 11 And 12 1f = 13 13 and &HFFFF) returned weird re (65535=&HFFFF ret ElseIf i = 11 Then 'or fals 13 = 11 Or 12 lf = 13 ElseIf i = 12 Then 'xor 13 = 11 Xor 12 1f = 13 Apparently anything bigger than 32767 bec ElseIf i = 13 Then 'test lf = 11 And 12 I ElseIf i = 14 Then 'inc 13 = 12 + 1 1f = 13 wonder how many more errors like that there are in my ElseIf i = 15 Then 'dec 'Flags 13 = 12 - 1 b3 = (13 And 32768) > 0lf = 13 .FLAGS = .FLAGS And (Not (flZ + flS + flO + flC + flP + flN)) .FLAGS = .FLAGS Or IIf(lf = 0, flZ, 0) . FLAGS = .FLAGS Or IIf((lf And 32768) > 0, flS, 0) ElseIf i = 16 Then 'neg 13 = (65535 - 12) + 1f = 13If bl = b2 Then .FLAGS = .FLAGS Or IIf(bl Xor b3, flo, 0) .FLAGS = .FLAGS Or IIf(lf And &HFFFF0000) > 0, flC, 0) ElseIf i = 17 Then 'not 13 = 65535 - 12 b1 = (.FLAGS And flS) lf = 13 b2 = (.FLAGS And fl0) ElseIf i = 18 Then 'bswp .FLAGS = .FLAGS Or IIf(b1 <> b2, flN, 0) 13 = (12 And 65280) \ 256 + (12 And 255) * 256 1f = 13 .FLAGS = .FLAGS Or IIf((b1=b2) And ((.FLAGS And flZ) = 0) , flp, ()) ElseIf i = 19 Then 'lshl A,Rn If scf = 1 Then .FLAGS = .FLAGS Or flC i = 11 And 15 If scf = 0 Then .FLAGS = .FLAGS And Not flC scf = IIf((12 And (2 (16 - i))) = 0, 0, 1)End If 13 = 12 * (<mark>2</mark> ^ i) lf = 13 ernal buses write ElseIf i = 20 Then 'lshr A,Rn If GFlg(.DIB(.eDP), mar_we) Then .MAR = .eEAB
If GFlg(.DIB(.eDP), mdr_we) Then .MDR = .eEDB i = 11 And 15 If i > 0 Then 'Write operations for IDB scf = IIf((12 And (2 ^ (i - 1))) = 0, 0, 1) If GFlg(.DIB(.eDP), reg_w) Then 13 = 12 \ (<mark>2</mark> ^ i) If .eSelectedReg = 4 Then .SP = .eIDB ElseIf .eSelectedReg = 5 Then Else 13 = 12 End If .IP = .eIDB ElseIf .eSelectedReg >= 0 And .eSelectedReg <= 3 Then</pre> lf = 13 ElseIf i = 21 Then 'ashl A,Rn .R(.eSelectedReg) = .eIDB i = 11 And 15 End If $scf = IIf((12 And (2 \land (16 - i))) = 0, 0, 1)$ End If If GFlg(.DIB(.eDP), acc_w) Then .A = .eIDB
If GFlg(.DIB(.eDP), flg_w) Then .FLAGS = .eIDB $13 = 12 * (2 ^ i)$ lf = 13 ElseIf i = 22 Then 'ashr A,Rn If GFlg(.DIB(.eDP), mdr_wi) Then .MDR = .eIDB 13 = 12 For i = 1 To 11 And 15 scf = IIf((12 And 1) = 0, 0, 1)If GFlq(.DIB(.eDP), mar wi) Then .MAR = .eIAB $13 = 12 \setminus 2$ If (12 And 32768) = 1 Then 13 = 13 Or 32768 Next If GFlg(.DIB(.eDP), reg_ipi) Then .IP = .IP + 1 lf = 13If GFlg(.DIB(.eDP), reg_spi) Then .SP = .SP + 2 ElseIf i = 23 Then 'rol A,Rn Return internal registers to zeroes in order to pr possibility of data transferred between microinstr 13 = 12For i = 1 To (11 And 15) 13 = ((13 * 2) And 65535) + IIf((13 And 32768) = 0, 0, 1) impossible in real Next scf = IIf((13 And 1) = 0, 0, 1)will not do this for now - wait until I will fix a bug with decoding ALU instructions -relied on values staying in these registers lf = 13ElseIf i = 24 Then 'ror A,Rn 13 = 12 For i = 1 To (11 And 15) 'DONE! - take next microinstruction $13 = (13 \setminus 2) + IIf((13 \text{ And } 1) = 0, 0, 32768)$ Next eDP = eDP + 1scf = IIf((13 And 32768) = 0, 0, 1)lf = 13 'Last loop of execution cycle
If .eDP = UBound(.DIB) + 1 Th ElseIf i = 25 Then 'rcl A,Rn + 1 Ther 13 = 12 scf = IIf((.FLAGS And flC) > 0, 1, 0)'Prepare registers For i = 1 To (11 And 15) .CIB = " .Fetch = True lf = 13 13 = ((13 * 2) And 65535) + IIf(scf = 1, 1, 0) .eDP = 0 scf = IIf((lf And 32768) = 0, 0, 1) ReDim .DIB(-1 To -1) Next 'Process interru 1f = 13eInterrupt ElseIf i = 26 Then 'rcr A,Rn End If 13 = 12scf = IIf((.FLAGS And flC) > 0, 1, 0)End With For i = 1 To (11 And 15) End Sub 1f = 13



If GFlg(db, alu_swp) Then s = s + "alu_swp, "
If GFlg(db, op_alu_sh) Then s = s + "alu_sh(" + CStr(db.nAluSh) +
"), "
If GFlg(db, op_alu_c) Then s = s + "alu_c(" + CStr(db.nAluOpNum) +
"), "

External buses write If GFlg(db, mar_we) Then s = s + "mar_we, " If GFlg(db, mdr_we) Then s = s + "mdr_we, " 'Write operations for IDB If GFlg(db, reg_w) Then s = s + "reg_w, " If GFlg(db, acc_w) Then s = s + "acc_w, " If GFlg(db, flg_w) Then s = s + "flg_w, " If GFlg(db, mdr_wi) Then s = s + "mdr_wi, "

22.6. plO

Option Explicit

interfaced through this unit (apart from window-related functions such as .Hide). Device modules should 'export' the following functions: Init - whatever intialisation devices want. This is not simulation-related - rather, interface-related, as this will be called only once at program startup. Reset - simulation-related initialisations. Should be similar to what a real device would do when it receives Reset signal. Reset will be sent when user presses the computer Reset button, or during program a device would do with every rise of OSC in a real CPU is what should be done in Tick. PortRead - will simulate arrival of PortRead signal. The function should return whatever the device would if the device decides to ignore the signal. PortWrite - simulates arrival of PortWrite signal. Public Sub ioInit() Initializes this module Public Sub ioInit() 1 End Sub Public Sub devInit() Public Sub devInit() fdVideo Init fdKeyboard.Init fdSpeaker.Init End Sub Public Sub devReset() Resets all devices

Public Sub devReset() fdVideo.Reset fdKeyboard.Reset fdSpeaker.Reset End Sub

' Public Sub devTick() ' ' Resets all devices '

Public Sub devTick()

fdVideo.Tick fdKeyboard.Tick



End Function

fdSpeaker.Tick End Sub



fsRegs.Update fhCPU.Update fhCU.Update

End With

End Function

22.7. fhCPU

General Wn MDP W MDP III III MDP III III IIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIII	🙀 CLab - I	Aicrosoft Visual Basic [design] - [fhCPU (Form)]	
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General Wn MDP W MDP III III MDP III III IIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIII		・ 🍸 🚔 🖥 🍓 🍋 🕫 🔹 🕨 💷 🔰 📾 📽 😭 🕂 🎖 🛠 🔂 📋 5160, 2520 🛛 🦉 1095 x 855	
B: -15384 -16384 2:0 0:0 0:0 0:0 0:0 0:0 P:0 0:0 P:0 0:0 P:0 0:0 B: -16384 B: -16384 B: -16384 C: Textl Internition B: Fice B: Fice B: -16384 C: Textl Internition B: Fice B: Fice B: -16384 D: Textl Internition B: Fice B: Fice B: -16384 P: Textl Internition B: Fice B: Fice B: -16384 P: Textl MAR -16384 Properties Properties </th <th>General Win k</th> <th>MDR</th> <th>CLab (CLab A</th>	General Win k	MDR	CLab (CLab A
Image: SP: Text1 MAR 16384 Alphabetic Call. Image: SP: Text1 SP: Text1 Image: SP: Text1 Audice: False Audice: False Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text1 Image: SP: Text2 Image: SP: Text2 Image: SP: Text2 Image: SP: Text2 Image: SP: Text2 <td></td> <td>B: -16384 -16384 C: 0 9: 0 0: 0 0: 0 P: 0 N: 0 T: 0 D: Text1 E: Text1 E: Text1</td> <td>C fhCu C fhRA C fhCa C fiRa C fiRa</td>		B: -16384 -16384 C: 0 9: 0 0: 0 0: 0 P: 0 N: 0 T: 0 D: Text1 E: Text1 E: Text1	C fhCu C fhRA C fhCa C fiRa C fiRa
Cycle (F/E) Execute Interrupt statu (B): To fork (F/EM): 2 0000000000000 To exec (EREM): 2 00000000000000 (Name) Returns the name used in code to identify an Form Layout X		MAR -16384	(Name) PctBask A Align 0 - None Appearan 1 - 3D AutoRedr.False BackColor 8460 BorderSty 1 - Fixec
		To fetch (FREM) • FEDCBA976543210 To exec (EREM) 2 00000000000000	ClpContrc True DataField (Name) Returns the name used in code to identify an Form Layout X

Option Explicit

''
' Public declarations in this module: '
t de la construcción de la constru
' PROCEDURES: '
' Reset '
··
11
' Public Sub Reset() '
r I
' Initialises the module and '
' related variables.
·
Public Sub Reset()
'Empty the DIB array
ReDim Proj.CPU.DIB(-1 To -1)
1002m 110J.000.010(1 10 1)
'Init registers
Proj.CPU.A = 0
Proj.CPU.R(0) = 0
Proj.CPU.R(1) = 0
Proj.CPU.R(2) = 0
Proj.CPU.R(3) = 0
Proj. CPU. IP = 0
Proj.CPU.SP = 24576 'meaning 6000h
Proj.CPU.FLAGS = fli
Proj.CPU.MAR = 0
Proj.CPU.MDR = 0
Proj.CPU.CIB = ""
Proj.CPU.Fetch = True
Proj.CPU.FREM = 0
Proj.CPU.fremMem = False
Proj.CPU.IS = 0
Proj.CPU.eDP = 0
Proj.CPU.eEAB = 0
Proj.CPU.eEDB = 0
Proj.CPU.eIAB = 0
Proj.CPU.eIDB = 0
Proj.CPU.eSelectedReg = 0
Proj.P.Code = ""

Proj.P.CompileNeeded = True



Listings

```
LFN.Caption = IIf((Proj.CPU.FLAGS And 256) > 0, "1", "0")
  LFP.Caption = IIf((Proj.CPU.FLAGS And 512) > 0, "1", "0")
  'Control Unit
 LFE.Caption = IIf(Proj.CPU.Fetch, "Fetch", "Execute")
 If Proj.CPU.Fetch Then
   If Proj.CPU.CIB = "" Then
     LFREM.Caption = "?"
   Else
     LFREM.Caption = IIf(Proj.CPU.fremMem, CStr(Proj.CPU.FREM + 1) +
"+", CStr(Proj.CPU.FREM))
   End If
   LEREM.Caption = "N/A"
 Else
   LEREM.Caption = CStr(UBound(Proj.CPU.DIB) - Proj.CPU.eDP + 1)
   LFREM.Caption = "N/A"
 End If
 LCIR.Caption = Str2Chr(Proj.CPU.CIB)
 LIS.Caption = Dec2Bin(Proj.CPU.IS, 16)
End Sub
```

blic Sub SetComplexit

Public Sub SetComplexity(If Proj.Complexity = 0 Then PctPanel.Picture = PctBasic.Picture LMDR.Visible = False LMAR.Visible = False LCIR.Visible = False LFZ.Visible = False LFS.Visible = False LFO.Visible = False LFC.Visible = False LFP.Visible = False LFN.Visible = False LFI.Visible = False llMDR.Visible = False llMAR.Visible = False llCIR.Visible = False llFLAGS.Visible = False llFZ.Visible = False llFS.Visible = False llFO.Visible = False llFC.Visible = False llFP.Visible = False llFN.Visible = False llFI.Visible = False LFE.Visible = False llFE.Visible = False llFREM.Visible = False LFREM.Visible = False LEREM.Visible = False llEREM.Visible = False LIS.Visible = False 11IS.Visible = False lllIS.Visible = False ElseIf Proj.Complexity = 1 Then PctPanel.Picture = PctAlevel.Picture LMDR.Visible = True LMAR.Visible = True LCIR.Visible = True

LFZ.Visible = True LFS.Visible = True LFO.Visible = True LFC.Visible = True LFP.Visible = True LFN.Visible = True LFI.Visible = True llMDR.Visible = True llMAR.Visible = True llCIR.Visible = True llFLAGS.Visible = True llFZ.Visible = True llFS.Visible = True llFO.Visible = True llFC.Visible = True llFP.Visible = True llFN.Visible = True llFI.Visible = True LFE.Visible = True llFE.Visible = True llFREM.Visible = True LFREM.Visible = True LEREM.Visible = True llEREM.Visible = True LIS.Visible = True llIS.Visible = True lllIS.Visible = True Else PctPanel.Picture = PctFull.Picture LMDR.Visible = True LMAR.Visible = True LCIR.Visible = True LFZ.Visible = True LFS.Visible = True LFO.Visible = True LFC.Visible = True LFP.Visible = True LFN.Visible = True LFI.Visible = True llMDR.Visible = True llMAR.Visible = True llCIR.Visible = True llFLAGS.Visible = True llFZ.Visible = True llFS.Visible = True llFO.Visible = True llFC.Visible = True llFP.Visible = True llFN.Visible = True llFI.Visible = True LFE.Visible = True llFE.Visible = True llFREM.Visible = True LFREM.Visible = True LEREM Visible = True llEREM.Visible = True LIS.Visible = True llIS.Visible = True lllIS.Visible = True End If End Sub

22.8. fhCU

🛠 CPU - Control Unit	2	×
Cycle (F/E): Execute To fetch (FREM): - To exec (EREM): 2	Interrupt status (IS): FEDCBA9876543210 000000000000000000	•
Current Instruction Register: 00 00 00 00 00 Decoded Instruction Register	r.	•
		-
		•

Option Explicit

<pre>Public declarations in this module: PROCEDURES: Reset Update '' 'Stores the string values for all registers displayed '' ' on this form to track changes and highlight respectively Private LastStr(0 To 8) As String '' ' Public Sub Init()</pre>	I LFE I LFR I LER I LIS I LCI S
' Initializes this module ' '' Public Sub Init() 'Update register values Update 'Save changes SaveLast	D L I Lis
"Update again to highlight nothing Update End Sub	DI2
<pre>' Private Sub Form_Unload(Cancel As Integer) ' ' DESCRIPTION: Event handler for Form_Unload ' ' unloads the form if the application is ' ' really shutting down, and just hides the ' ' form in case the user requested to close ' ' it.</pre>	E End ' ' P ' S
Private Sub Form_Unload(Cancel As Integer) If Not Appp.Terminating Then Cancel = 1 fhCU.Hide End If End Sub	' h ' U ' Pri L L L
Public Sub Update() Updates the contents of the window to reflect changes to the state of the simulation.	L End ' ' P
<pre>Public Sub Update()</pre>	' U ' Pri L f. End
LEREM.Caption = "N/A" Else LEREM.Caption = CStr(UBound(Proj.CPU.DIB) - Proj.CPU.eDP + 1) LFREM.Caption = "N/A" End If	' U ' c ' Pri

LCIB.Caption = Str2Chr(Proj.CPU.CIB) LIS.Caption = Dec2Bin(Proj.CPU.IS, 16)

```
'Register colors
If LastStr(0) <> LFE.Caption Then LFE.ForeColor = &HFF Else
E.ForeColor = 0
If LastStr(1) <> LFREM.Caption Then LFREM.ForeColor = &HFF Else
REM.ForeColor = 0
If LastStr(2) <> LEREM.Caption Then LEREM.ForeColor = &HFF Else
REM.ForeColor = 0
If LastStr(3) <> LIS.Caption Then LIS.ForeColor = &HFF Else
S.ForeColor = 0
If LastStr(7) <> LCIB.Caption Then LCIB.ForeColor = &HFF Else
IB.ForeColor = 0
SaveLast
'Dec
     oded instru
                 tion buffer
Dim i As Integer
ListDIB.Clear
If Proj.CPU.Fetch Then
 Call ListDIB.AddItem("<N/A>")
Else
 If UBound(Proj.CPU.DIB) = -1 Then Call
stDIB.AddItem("<Empty>")
For i = 0 To UBound(Proj.CPU.DIB)
  Call ListDIB.AddItem(IIf(Proj.CPU.eDP = i, "--> ", " ") +
2Str(Proj.CPU.DIB(i)))
 Next
End If
```

fhCU.Caption = CStr(Proj.TickCount)

d Sub

Saves the state of all register in order to nighlight them as they change. Is called by Jpdate() after getting new values for them ivate Sub SaveLast() LastStr(0) = LFE.Caption LastStr(1) = LFREM.Caption LastStr(2) = LEREM.Caption LastStr(3) = LIS.CaptionLastStr(7) = LCIB.Captiond Sub Private Sub ListDIB_Click() clicks on it to remove the blue line ivate Sub ListDIB_Click() ListDIB.ListIndex = -1 fhCU.SetFocus d Sub rivate Sub ListDIB_GotFocus()

Unselects the DIB every time the user clicks on it to remove the blue line ivate Sub ListDIB_GotFocus()

ListDIB.ListIndex = -1fhCU.SetFocus

End Sub

Update End Sub

End Sub

Public Sub Reset()

Public Sub Reset() 'Clear RAM memory Dim i As Long

22.9. fhRAM



Option Explicit

()
' Public declarations in this module: '
' PROCEDURES: '
' Reset '
'Update '
''
'Prevent specific events
Private BlockSetEditCell As Boolean
'Fonts
Private fFixed As TFnt 'Fixed cells
Private fDef As TFnt 'Usual cells
Private fSel As TFnt 'Cursor
Private fIP As TFnt 'IP
Private fSP As TFnt 'SP
Private fIPil As TFnt 'IP cur instruction
Private fIPi2 As TFnt 'IP cur instruction (uncertain) Private fCode As TFnt 'Memory loaded with compiled code
Private fCode As TFnt 'Memory loaded with compiled code Private fVal As TFnt 'Other memory with nonzero values
VIIVate Ival as Irne Other memory with homzero varues
()
' Public Sub Init()
n de la constante de la constan
' DESCRIPTION: initialises this module. '
Public Sub Init() 'SG fonts - initial params
fSel.FaceName = "Courier New"
fSel.Size = -11
fSel.Weight = 400
fFixed = fSel
fDef = fSel
fSP = fSel
fIP = fSel
fIPi1 = fSel
fIPi2 = fSel fCode = fSel
fVal = fSel
'SG fonts - colors
fFixed.ForeColor = GetSysColor(COLOR_BTNTEXT)
fFixed.BackColor = GetSysColor(COLOR_BTNFACE)
fDef.ForeColor = &HCOCOCO
fDef.BackColor = &HFFFFFF
fSel.ForeColor = GetSysColor(COLOR_HIGHLIGHTTEXT)
<pre>fSel.BackColor = GetSysColor(COLOR_HIGHLIGHT) fSP.ForeColor = 0</pre>
fSP.BackColor = 65280
fIP.ForeColor = &H80
fIP.BackColor = &HFFB366
fIP.Weight = 800
fIPil.ForeColor = &H80
fIPil.BackColor = &HFFEOC1
fIPil.Weight = 800
fIPi2.ForeColor = &H80

fIPi2.BackColor = &HFFF2E6 fIPi2.Weight = 800 fCode.ForeColor = 0fCode.BackColor = &H1FFFF
fVal.ForeColor = 0 fVal.BackColor = &HEOEOFF 'SG f CreateFnt fFixed CreateFnt fDef CreateFnt fSel CreateFnt fIP CreateFnt fSP CreateFnt fIPil CreateFnt fIPi2 CreateFnt fCode CreateFnt fVal 'Set SG option SG.Option(goEditing) = True SG.Option(goColMoving) = False SG.Option(goColSizing) = False SG.Option(goRangeSelect) = False SG.Option(goRowMoving) = False SG.Option(goRowSizing) = False SG. Option (goThumbTracking) = True 'SG visual SG.ColCount = 17 SG.RowCount = 4097 SG.FixedCols = 1 SG.FixedRows = 1 Dim sz As Size Call SelectObject(SG.hdc, fDef.fntFont) Call GetTextExtentPoint32(SG.hdc, "FF", Len("FF"), sz) GG.DefaultColWidth = sz.cx + 4 SG.DefaultColWidth = sz.cx + 4 SG.DefaultRowHeight = sz.cy + 2 Call SelectObject(SG.hdc, fDef.fntFont) Call GetTextExtentPoint32(SG.hdc, "FFFF", Len("FFFF"), sz) SG.ColWidths(0) = sz.cx + 4BlockSetEditCell = False Reset

ReDim Proj.RAM(0 To 65535) 'Redim sets everything to zeroes

Listings



<pre>' Call Errr("The string you entered is not a number. It should contain 0-9 for decimal numbers, 0-9 & A-F for hexadecimal numbers ending with 'h', and 0-1 for binary numbers ending with 'b'.") ' GoTo retry ' End If ' n = Hex2Dec(Left(Addr, Len(Addr) - 1)) ' ElseIf Right(Addr, 1) = "b" Then ' If Not TestCharset(Left(Addr, Len(Addr) - 1), "01") Then ' Call Errr("The string you entered is not a number. It should contain 0-9 for decimal numbers, 0-9 & A-F for hexadecimal numbers ending with 'h', and 0-1 for binary numbers ending with 'b'.") ' GoTo retry ' End If ' n = Bin2Dec(Left(Addr, Len(Addr) - 1)) ' Else</pre>	<pre>' GoTo retry ' End If ' rgFrom = n \ 16 ' If rgFrom > 4095 - 7 Then rgFrom = 4095 - 7 ' Update ' Exit Sub ' 'Errrr: ' Call Errr("The address should be between 0 and 65535 (0000h and 0FFFFh)") ' Resume retry End Sub Private Sub BtnShowIP_Click() ' rgFrom = Proj.CPU.IP \ 16</pre>
 If Not TestCharset(Left(Addr, Len(Addr) - 1), "0123456789") Then Call Errr("The string you entered is not a number. It should 	' If rgFrom > 4095 - 7 Then rgFrom = 4095 - 7 ' Update
contain 0-9 for decimal numbers, 0-9 & A-F for hexadecimal numbers ending with 'h', and 0-1 for binary numbers ending with 'b'.")	End Sub
<pre>ending with n , and of for Dinary numbers ending with D .) GoTo retry End If n = CLng(Addr) End If If n < 0 Or n > 65535 Then Call Errr("The address should be between 0 and 65535 (0000h and 0FFFFh)")</pre>	<pre>Private Sub BtnShowSP_Click() ' rgFrom = Proj.CPU.SP \ 16 ' If rgFrom > 4095 - 7 Then rgFrom = 4095 - 7 ' Update End Sub</pre>

22.10. fdKeyboard

	🔀 Keyboard	controller ?X
	DBO	Keyboard controller Port 60h, IRQ1
	СВО	Pending-key: 2Eh Right-arrow Last key: 2Dh Left arrow
	ABO	IRQ: Accepted
Option Explicit	-	KeyName(49) = "Circle"
'Key code to key name map Private KeyName(0 To 51) As String		KeyName(50) = "Square" KeyName(51) = "Triangle" "Reset once
'Key code to be sent to programs Private CurKey As Integer 'Last key pressed - just for user informatic	on	Reset End Sub
Private LastKey As Integer 'Whether interrupt has been sent Private InterruptSent As Boolean		' Public Sub Reset() '
'' ' Public Sub Init() '		' Callback for pIO. ' '' Public Sub Reset()
' Initializes this module '		CurKey = -1 LastKey = -1 InterruptSent = False
Public Sub Init() 'Initialize key names Dim i As Integer		Update End Sub
<pre>For i = 0 To 25 KeyName(i) = Chr\$(i + 65) Next</pre>		'' ' Public Sub Tick() '
KeyName(26) = "." KeyName(27) = "Enter" KeyName(28) = "Spacebar"		' Callback for pIO. ' '' Public Sub Tick()
KeyName(29) = "=" For i = 0 To 9 KeyName(i + 30) = Chr\$(i + 48)		<pre>Send interrupt If CurKey <> -1 And Not InterruptSent Then InterruptSent = IRO(1)</pre>
Next KeyName(40) = "Numpad ."		Update End If
KeyName(41) = "/" KeyName(42) = "*" KeyName(43) = "-"		End Sub
KeyName(44) = "+" KeyName(45) = "Left arrow" KeyName(46) = "Right arrow"		' Public Function PortRead(PortNum ' ' As Integer) As Long ' '
KeyName(47) = "Up arrow" KeyName(48) = "Down arrow"		' Callback for pIO. '

Listings



22.11. fdSpeaker



TmrFreg.Enabled = True	
Else TmrFreg.Enabled = False	' Private Sub Pct Paint() '
End If	··
Update	Private Sub Pct_Paint()
End Sub	If spkState Then
	Call Pct.PaintPicture(PctH.Picture, 0, 0)
··	Else
' Private Sub Update() '	Call Pct.PaintPicture(PctL.Picture, 0, 0)
	End If
' Updates screen to reflect '	End Sub
' current situation. '	
··	()
Private Sub Update()	' Private Sub TmrFreg Timer() '
'Refresh picture	··
<pre>If LState.Caption <> IIf(spkState, "High", "Low") Then Pct_Paint</pre>	Private Sub TmrFreq_Timer()
LState.Caption = IIf (spkState, "High", "Low")	spkState = Not spkState
LFreq.Caption = IIf(spkFreq = 0 Or spkFreq = 1, "N/A", CStr(Int(10	Update
/ 65535 * spkFreg * 100) / 100) + " Hz")	End Sub
End Sub	

22.12. fdVideo



Option Explicit

((
' Public declarations in this module: '
PROCEDURES:
' Reset '
' SetVideoMode '
' UpdateScr '
··
۱۱
' PORT STATE VARIABLES: '
··
' port51state values: '
' -1: no operations '
' other: first word written to 51h '
··
' port53state values: '
1
' 0000h: no operation '
' 0100h: set pixel '
' 0500h: set pen colour '
' 0600h: set brush colour '
' 1000h: draw line '
' 1010h: set pen position '
' 1020h: continue line '
' 2000h: draw empty circle '
' 2001h: draw filled circle '
' 3000h: draw empty rectangle '
' 3001h: draw filled rectangle '
Prinche problighte be tere
Private port51state As Long Private port53state As Long
Private port53state As Long Private port53var1 As Long 'depends on function
Private port53var1 As Long 'depends on function
rivate portosvarz As holig depends on function
11
' Private Sub Form_Unload(Cancel As Integer) '





Next PortRead = IIf(Proj.Video.autoUpdate, 1, 0) Update screen Call StretchBlt(fiComp.PctScr.hdc, 0, 0, fiComp.PctScr.Width, Case Else fiComp.PctScr.Height, PortRead = 999999 Proj.Video.vDC.hdc, 0, 0, Proj.Video.vDC.Width, Exit Function Proj.Video.vDC.Height, SRCCOPY) End Select Call StretchBlt(fiDisplay.PctScr.hdc, 0, 0, fiDisplay.PctScr.Width, End Function fiDisplay.PctScr.Height, Proj.Video.vDC.hdc, 0, 0, Proj.Video.vDC.Width, Proj.Video.vDC.Height, SRCCOPY) Public Sub PortWrite(PortNum End Sub As Integer, Dt As Long) Public Sub Init() Public Sub PortWrite(PortNum As Integer, Dt As Long) Initializes this module Dim 11 As Long Select Case PortNum Public Sub Init() Case &H50 !---'Create SCREEN MODE Set Proj.Video.vDC = New VirtualDC If ((Dt >= 1) And (Dt <= 7)) Or ((Dt >= 129) And (Dt <= 135)) Then 'Reset Reset Call SetVideoMode(Dt) Update UpdateScr End Sub End If Case &H51 Public Sub Reset() '--- PALETTE If port51state = -1 Then Callback for pIO. port51state = Dt Else 11 = port51state And 255 Proj.Video.PalMem(11) = CLng((port51state And 65280) \ Public Sub Reset() Dim i As Integer, A As Integer Dim R As Long, g As Long, b As Long 256) + Dt * CLng(256) port51state = -1 'Init palette memory For i = 1 To 7 End If For A = 0 To 31 Proj.Video.PalMem((i - 1) * 32 + A) = CLng(IIf((i And 1) > 0))Case &H52 A, 0)) * 8 + CLng(IIf((i And 2) > 0, A, 0)) * 2048 + CLng(IIf((i And MEMORY '--- MEMORY ---' Proj.Video.MemOff = Dt 4) > 0, A, 0)) * 524288Next Case &H53 Next 'RAM video memory offset -- DRAWING ---If port53state = Proj.Video.MemOff = 57344 '&HE000 - &HEFFF 0 Then 'Get function Call SetVideoMode(1) port53state = Dt If (port53state = &H100) Or (port53state = &H500) Or _ 'Reset port state (port53state = &H600) Or (port53state = &H1000) Or _ (port53state = &H1010) Or (port53state = &H1020) Or _ (port53state = &H2000) Or (port53state = &H2001) Or _ port51state = -1port53state = 0 'Other variables (port53state = &H3000) Or (port53state = &H3001) _ Proj.Video.autoUpdate = True Then port53state = 0 End Sub Else Select Case port53state Case &H100 Public Sub Tick() '--- Set pixel Case &H50 '--- Set Case &H600 '--- Set h Public Sub Tick(Case &H1000 End Sub Case &H101 Case &H102 Case &H2000 Callback for pIO. Draw Case &H2 Public Function PortRead(PortNum As Integer) As Long Select Case PortNum Case &H300 Case &H50 '--- Draw SCREEN MODE Case &H3001 PortRead = Proj.Video.Mode '--- Dr Case Else filled rectangle port53state = 0 Case &H51 - PALETTE End Select PortRead = 0 End If Case &H52 Case &H54 '--- MEMORY ---' PortRead = Proj.Video.MemOff UPDATINC If Dt = 0 Then Proj.Video.autoUpdate = False Case &H53 ElseIf Dt. = 1 Then DRAWING Proj.Video.autoUpdate = True PortRead = 0 Else If Not fiMain.MIRefNothing.Checked Then UpdateScr Case &H54 End If

End Select	' Private Sub BtnRefresh_Click() '
End Sub	Private Sub BtnRefresh_Click() UpdateScr
۱ <u>ـــــ</u> ۱	End Sub

22.13. fiMain



Option Explicit

	'Unhook this window to be sure to terminate nicely		
'' ' Public declarations in this module: '	Unhook 'Unload all loaded forms to terminate application.		
vubic declaracions in this module.	'We should NOT use the End statement, otherwise		
' PROCEDURES: '	' VB will crash. Took me an hour to figure it out.		
' Init '	Dim i As Integer		
' WindowProc	<pre>For i = Forms.Count - 1 To 0 Step -1 Unload Forms(i)</pre>		
	Next		
'Indicates whether a form was visible prior to minimizing	End Sub		
Private frmWasVisible(0 To 20) As Boolean			
'Time (GetTickCount) when the last Step was executed (Run mode) Private lastStepTime As Long	'' ' Private Sub Hook() '		
··	' Hook this window to process minimize event. '		
' Public Sub Init()			
' . ' Initialises this form. This is called '	<pre>Private Sub Hook() Appp.PrevWndProc = SetWindowLong(fiMain.hwnd, GWL WNDPROC,</pre>		
' before the form is actually shown. '	AddressOf pGlobals.WindowProc)		
··	End Sub		
Public Sub Init() 'Need to set Top to 0 because when visual styles are used			
'it is positions of client areas, and not captions, that	' Private Sub Unhook()		
'stay fixed. So a window's real Top is higher on screen			
'with visual styles.	' Unhook this window prior to unloading it.		
Top = 0	'' Private Sub Unhook()		
'Change visuals depending on OS	'Set previous (VB) window procedure		
If Appp.RunningOnWinXP Then	Call SetWindowLong(fiMain.hwnd, GWL_WNDPROC, Appp.PrevWndProc)		
fiMain.BackColor = &HEDEFEF 'comct132 v6.0 gives us toolbars with	End Sub		
this particular color Height = 1830	11		
Else	' Public Function WindowProc(ByVal hw As Long, '		
fiMain.BackColor = GetSysColor(COLOR_BTNFACE)	' ByVal uMsg As Long, ByVal wParam As Long, '		
Height = 1710 End If	ByVal lParam As Long) As Long		
End II	' DESCRIPTION: Window procedure for fiMain '		
'Install own window procedure	··		
Hook	Public Function WindowProc(ByVal hw As Long, ByVal uMsg _		
'Init toolbars	As Long, ByVal wParam As Long, ByVal lParam As Long) As Long		
SetToolbarButtons	If uMsg = WM_SYSCOMMAND Then		
End Sub	<pre>If wParam = SC_MINIMIZE Then</pre>		
	'Minimize everything, saving state		
'' ' Private Sub Form_Unload(Cancel As Integer) '	<pre>frmWasVisible(0) = fhCPU.Visible frmWasVisible(1) = fhCU.Visible</pre>		
	frmWasVisible(2) = fhRAM.Visible		
' DESCRIPTION: Unlike event handlers for	<pre>frmWasVisible(3) = fdVideo.Visible</pre>		
Unload on all other forms, this handler ' should shut down the application in a '	<pre>frmWasVisible(4) = fiComp.Visible frmWasVisible(5) = fsCode.Visible</pre>		
proper way, which involves unloading all '	frmWasVisible(6) = fsRegs.Visible		
forms. After that the application will '	frmWasVisible(7) = fsStack.Visible		
' terminate. '	<pre>frmWasVisible(8) = fsVars.Visible</pre>		
' NOTES: prior to unloading forms, this proc '	fhCPU.Hide fhCU.Hide		
will query every form whether it is OK	fhRAM.Hide		
' for it to shut down, giving a chance to '	fdVideo.Hide		
' save work for instance. If a form wants '	fiComp.Hide		
' to cancel shutdown, it will return False ' ' as the result of ShutdownQuery. '	fsCode.Hide fsRegs.Hide		
''	fsStack.Hide		
Private Sub Form_Unload(Cancel As Integer)	fsVars.Hide		
'Indicate we're really shutting down	ElseIf wParam = SC_RESTORE Then		
Appp.Terminating = True	<pre>fhCPU.Visible = frmWasVisible(0)</pre>		
fhou uisible - formus uisible	. (1)		1
---	--	---------	---
<pre>fhCU.Visible = frmWasVisible fhRAM.Visible = frmWasVisibl fdVideo.Visible = frmWasVisib fiComp.Visible = frmWasVisib fsCode.Visible = frmWasVisib faRegs.Visible = frmWasVisib fsStack.Visible = frmWasVisib</pre>	Le(2) ible(3) ole(4) ole(5) ole(6)		Private Sub MIIDisplay_Clic Private Sub MIIDisplay_Click(fiDisplay.Show End Sub
fsVars.Visible = frmWasVisib End If 'Call VB window proc + default			'' Private Sub MIIKbd_Click()
WindowProc = CallWindowProc(Ag lParam) ElseIf uMsg = WM_NCLBUTTONDOWN 7 If wParam = HTCAPTION Then fil	opp.PrevWndProc, hw, uMsg, T hen	wParam,	Private Sub MIIKbd_Click() fiKeyboard.Show End Sub
<pre>'Call VB window proc + default WindowProc = CallWindowProc(Ag lParam) Else</pre>	ppp.PrevWndProc, hw, uMsg,	wParam,	' ' Private Sub MIISpeaker_Clic ' Private Sub MIISpeaker_Click(
<pre>'Call VB window proc + default WindowProc = CallWindowProc(Ag lParam)</pre>		wParam,	End Sub
End If End Function			' ' Private Sub MIHRAM_Click() '
' ' Private Sub BtnWritePrg_Click() ' Private Sub BtnWritePrg_Click()	- 1 - 1 - 7		Private Sub MIHRAM_Click() fhRAM.Show End Sub
fsCode.Show End Sub			' ' Private Sub MIHBuses_Click('
<pre>Private Sub Command3_Click() Step fiMain.UpdateAll End Sub</pre>			Private Sub MIHBuses_Click() End Sub
Private Sub Command4_Click() Tick			' Private Sub MIHCPU_Click()
fiMain.UpdateAll End Sub	_		Private Sub MIHCPU_Click() fhCPU.Show End Sub
' Private Sub MICompBasic_Click()			'' ' Private Sub MIHCU_Click() '
<pre>Private Sub MICompBasic_Click() Proj.Complexity = 0 fhCPU.SetComplexity MICompBasic.Checked = True MICompMed.Checked = False IICompMed.Checked = False</pre>			Private Sub MIHCU_Click() fhCU.Show End Sub
MICompFull.Checked = False End Sub			'' Private Sub MIHALU_Click()
' ' Private Sub MICompMed_Click() ' '' Private Sub MICompMed_Click()			Private Sub MIHALU_Click()
Proj.Complexity = 1 fhCPU.SetComplexity MICompBasic.Checked = False MICompMed.Checked = True			' ' Private Sub MIHKbd_Click() ' Private Sub MIHKbd_Click()
MICompFull.Checked = False End Sub			fdKeyboard.Show End Sub
'' ' Private Sub MICompFull_Click() '			'' Private Sub MIHVid_Click()
<pre>Private Sub MICompFull_Click() Proj.Complexity = 2 fhCPU.SetComplexity MICompBasic.Checked = False</pre>	-		Private Sub MIHVid_Click() fdVideo.Show End Sub
MICompMed.Checked = False MICompFull.Checked = True End Sub			' Private Sub MIHSpk_Click()
'' ' Private Sub MIExit_Click() '			Private Sub MIHSpk_Click() fdSpeaker.Show End Sub
Private Sub MIExit_Click() Unload fiMain End Sub			' Private Sub MIDCode_Click()
'' ' Private Sub MIIComp_Click() ' ''			Private Sub MIDCode_Click() fsCode.Show End Sub
Private Sub MIIComp_Click() fiComp.Show End Sub			' Private Sub MIDRegs_Click()
			Private Sub MIDRegs_Click()

fsRegs.Show End Sub	' Private Sub MISpd05_Click() ' ' Private Sub MISpd05_Click()
'' ' Private Sub MIDVars_Click() '	MISpdMax.Checked = False MISpd01.Checked = False MISpd05.Checked = True
Private Sub MIDVars_Click() fsVars.Show	MISpd20.Checked = False MISpd60.Checked = False
End Sub	End Sub
'' ' Private Sub MIDStack_Click() '	'' ' Private Sub MISpd20_Click() ' '
Private Sub MIDStack_Click() fsStack.Show	Private Sub MISpd20_Click() MISpdMax.Checked = False
End Sub	MISpd01.Checked = False MISpd05.Checked = False
' Private Sub MINBin_Click() '	MISpd20.Checked = True MISpd60.Checked = False End Sub
Private Sub MINBin_Click() Proj.NmbRep = 1	··
Call fiMain.UpdateAll(True) MINBin.Checked = True	' Private Sub MISpd60_Click() ' ''
MINDecS.Checked = False MINDecU.Checked = False	Private Sub MISpd60_Click() MISpdMax.Checked = False
MINHex.Checked = False End Sub	MISpd01.Checked = False MISpd05.Checked = False
۰۰	MISpd20.Checked = False MISpd60.Checked = True
' Private Sub MINDecS_Click() '	End Sub
Private Sub MINDecS_Click() Proj.NmbRep = 3	'' ' Private Sub MIRefNothing_Click() '
Call fiMain.UpdateAll(True) MINBin.Checked = False	<pre>'' Private Sub MIRefNothing_Click()</pre>
MINDecS.Checked = True MINDecU.Checked = False	MIRefNothing.Checked = True MIRefUser.Checked = False
MINHex.Checked = <mark>False</mark> End Sub	MIRefAll.Checked = False End Sub
'' ' Private Sub MINDecU_Click() '	'' ' Private Sub MIRefUser_Click() '
Private Sub MINDecU_Click()	Private Sub MIRefUser_Click()
Proj.NmbRep = 2 Call fiMain.UpdateAll(True)	MIRefNothing.Checked = False MIRefUser.Checked = True
MINBin.Checked = False MINDecS.Checked = False	MIRefall.Checked = False End Sub
MINDecU.Checked = True	
MINHex.Checked = False End Sub	' Private Sub MIRefAll_Click() '
'' ' Private Sub MINHex_Click() '	<pre>Private Sub MIRefAll_Click() MIRefNothing.Checked = False</pre>
Private Sub MINHex Click()	MIRefUser.Checked = False MIRefAll.Checked = True
Proj.MubRep = 0 Call fiMain.UpdateAll(True)	End Sub
MINBin.Checked = False MINDecS.Checked = False	'' ' Private Sub TmrRun_Timer(Index As Integer) '
MINDecU Checked = False MINDecU Checked = True	' Calls Step if project is running. The
End Sub	' procedure is not executed unless a given ' ' time interval has passed (speed control). '
	' An array of timers is used to increase '
' Private Sub MISpdMax_Click() '	call frequency.
Private Sub MISpdMax_Click() MISpdMax.Checked = True	' NOTE: UpdateAll without enforcement will ' ' be called after each Step '
MISpd01.Checked = <mark>False</mark> MISpd05.Checked = <mark>False</mark>	Private Sub TmrRun_Timer(Index As Integer)
MISpd20.Checked = <mark>False</mark> MISpd60.Checked = <mark>False</mark>	'Check timing Dim i As Integer
End Sub	<pre>If MISpdMax.Checked Then i = 0 If MISpd01.Checked Then i = 100</pre>
'' ' Private Sub MISpd01_Click() '	If MISpd05.Checked Then i = 500 If MISpd20.Checked Then i = 2000
Private Sub MISpd01_Click()	<pre>If MISpd60.Checked Then i = 6000 If GetTickCount - lastStepTime < i Then Exit Sub</pre>
MISpdMax.Checked = False MISpd01.Checked = True	Execute one step If Proj.Running And Not Proj.Paused Then
MISpd05.Checked = False	lastStepTime = GetTickCount
MISpd20.Checked = False MISpd60.Checked = False	Step UpdateAll
End Sub	End If End Sub
··	

11	'TICK'
Public Sub UpdateAll(Optional Force	fsCode.MITick_ForFIMAIN
As Boolean = False)	ElseIf Button.Index = 2 Then 'STEP'
' Updates all windows as necessary. Set '	fsCode.MIStep_Click
' Force=True to update all windows in ' ' any situation. '	End If End Sub
· <u></u> ·	
Public Sub UpdateAll(Optional Force As Boolean = False) If Not Force And Proj.Running And Not Proj.Paused Then	۲۲
If MIRefNothing.Checked Then Exit Sub	' Private Sub ToolbarWnd_ButtonClick(ByVal Button As
End If If Force Or Not MIRefUser.Checked Or Not Proj.Running Then	ComctlLib.Button) '
'Update all windows	··
fhCPU.Update	Private Sub ToolbarWnd_ButtonClick(ByVal Button As
fhCU.Update fhRAM.Update	ComctlLib.Button) If Button.Index = 1 Then
fsCode.Update	'CPU'
fsRegs.Update fsStack.Update	fhCPU.Show ElseIf Button.Index = 2 Then
fsVars.Update	'RAM'
fdVideo.Update If Not Proj.Paused And fiMain.MISpdMax.Checked Then Call	fhRAM.Show ElseIf Button.Index = 3 Then
<pre>fsCode.RTB.MarksSetVisible(0, False)</pre>	'Registers'
Else 'Update user interface windows only	fsRegs.Show ElseIf Button.Index = 4 Then
fdVideo.Update	'Variables'
End If End Sub	fsVars.Show ElseIf Button.Index = 5 Then
	'Stack'
'' ' Public Sub ResetAll()	fsStack.Show End If
1	End Sub
' Resets the computer, initialising ' ' all hardware	1
''	
Public Sub ResetAll() Proj.TickCount = 0	' Private Sub ToolbarWnd2_ButtonClick(ByVal Button As ComctlLib.Button) '
Proj.Running = False	Comecting.Button)
Proj.Paused = False	'
Proj.P.CompileNeeded = True fhCPU. Reset	Private Sub ToolbarWnd2_ButtonClick(ByVal Button As ComctlLib.Button)
fhRAM.Reset	<pre>If Button.Index = 1 Then</pre>
devReset fsCode.RTB.ReadOnly = <mark>False</mark>	fiComp.Show
Call fsCode.RTB.MarksSetVisible(0, False)	ElseIf Button.Index = 2 Then
End Sub	'Display' fiDisplay.Show
·	Elself Button.Index = 3 Then
' ' Private Sub ToolbarProgram_ButtonClick(ByVal Button As	'Keyboard' fiKeyboard.Show
ComctlLib.Button) '	ElseIf Button.Index = 4 Then
··	' <u>Speaker</u> ' fdSpeaker.Show
Private Sub ToolbarProgram_ButtonClick(ByVal Button As	End If
ComctlLib.Button) 'WRITE A PROGRAM'	End Sub
If Button.Index = 1 Then fsCode.Show	
End Sub	' Public Sub SetToolbarButtons() ' ''
۲	Public Sub SetToolbarButtons()
' ' Private Sub ToolbarRun_ButtonClick(ByVal Button As	'Set text and image for Start If Proj.Running Then
ComctlLib.Button) '	If Proj.Paused Then
// '	ToolbarRun.Buttons(1).Caption = "Continue" ToolbarRun.Buttons(1).Image = 1
Private Sub ToolbarRun_ButtonClick(ByVal Button As ComctlLib.Button)	Else
If Button.Index = 1 Then START	ToolbarRun.Buttons(1).Caption = "Pause" ToolbarRun.Buttons(1).Image = 2
fsCode.MIStart_Click	End If
ElseIf Button.Index = 2 Then 'RESET'	Else ToolbarRun.Buttons(1).Caption = "Start"
fsCode.MIReset_Click	ToolbarRun.Buttons(1).Image = 1
End If End Sub	End If 'Set enabled for Stop
	If Proj.Running Or Proj.Halted Then
'' '	ToolbarRun.Buttons(2).Enabled = True Else
' Private Sub ToolbarStep_ButtonClick(ByVal Button As	ToolbarRun.Buttons(2).Enabled = False
ComctlLib.Button) '	End If 'Set Enabled for step
'	ToolbarStep.Buttons(1).Enabled = Not Proj.Running Or Proj.Paused
Private Sub ToolbarStep_ButtonClick(ByVal Button As ComctlLib.Button) If Button.Index = 1 Then	ToolbarStep.Buttons(2).Enabled = Not Proj.Running Or Proj.Paused
II Button.index = 1 inen	End Sub

22.14. fiComp



Option Explicit



in the keyboard "controller".

22.15. fiKeyboard



22.16. fiDisplay





22.17. fsCode



form in case the user requested to close

Option Explicit

Public declarations in this module: PROCEDURES: Init Update Private sub Form_Unload(Cancel As Integer) If Not Appp.Terminating Then Cancel = 1 fscode.Hide Bnd If Private errLine As Integer 'highlight an error Private Multication Private Sub Form_Resize() Private Sub Init() 'Public Sub Init() Public Sub Init() Private Sub Form_Resize() LERT.Left = 8 LERT.Width = PnlErr.Width - 16 Private Sub Form_Resize() LERT.Left = 8 LERT.Width = PnlCode.Height - 16 RTB.Height = PnlCode.Height - 16 RTB.MarksZoffset = 0 RTB.MarksZoffset = 0 RTB.StoptionFlag(coAutoIndent, True) Call RTB.SetOptionFlag(coTrimTrailingSpaces, True) RTB.StoptionFlag(coTrimTrailingSpaces, True) RTB.StoptionFlag(coTrimTrailingSpaces, True) RTB.StoptionFlag(coTrimTrailingSpaces, True) RTB.StoptionFlag(coTrimTrailingSpaces, True) RTB.StoptionFlag(coTrimTrailingSpaces, True) RTB.StoptionFlag(coTrimTrailingSpaces, True) RTB.S
PROCEDURES: Init Init Gancel = 1 Update fsCode.Hide Private errLine As Integer 'highlight an error Private wngLine As Integer 'highlight an error Private Bkpt() As Integer 'breakpoints line numbers Public Sub Init() 'Resizes all controls on the form in order 'bESCRIPTION: initialises this module. 'resizes all controls on the form in order 'bESCRIPTION: initialises this module. 'resizes all controls on the form in order 'bescription: 'nesizes all controls on the form in order 'bescription: 'resizes all controls on the form in order 'bescription: 'resizes all controls on the form in order 'bescription: 'resizes all controls on the form in order 'bescription: 'resizes all controls on the form in order 'bescription: 'resizes all controls on the form in order 'bescription: 'resizes all controls on the form size.' 'resizes all controls on the form in order 'resizes all controls on the form size.' 'resizes all controls on the form in order 'resizes all controls on the form size.' 'resizes all controls on the form in order 'resizes all controls on the form size.' 'resizes all
PROCEDURES: Init Update Cancel = 1 fsCode.Hide End If End Sub Private errLine As Integer [highlight an error Private wngLine As Integer [highlight a warning] Private sub Form_Resize() Private Sub Form_Resize() Private Sub Form_Resize() Private Sub Form_Resize() Private Sub Form_Resize() DESCRIPTION: initialises this module. Private Sub Form_Activate() Private Sub Form_Activate() Private Sub Form_Activate()
<pre>Init Update Init Update Init Update Init Update Init Ipdate Init Init InitIpdate I</pre>
Update End If Image: Private errLine As Integer 'highlight an error Private wngLine As Integer 'highlight a warning Private wngLine As Integer 'breakpoints line numbers 'Public Sub Init() 'resizes all controls on the form in order 'Public Sub Init() 'resizes sub Form_Resize() 'Private Sub Form_Resize() 'LErr.Left = 8 'Public Sub Init() 'LErr.Width = PnlErr.Width - 16 Public Sub Init() Private Sub Form_Resize() 'Public Sub Init() Private Sub Form_Resize() 'Private Sub Form_Resize() Err.Width = PnlErr.Width - 16 Problic Sub Init() Private Sub Form_Resize() 'RTB highlighting PnlCode.Height = PnlErr.Top errLine = -1 RTB.Top = 8 wngLine = -1 RTB.WarksLoffset = 0 RTB.MarksLoffset = 0 RTB.MarksLoffset = 0 RTB.MarksLofftMargin = 2 'Private Sub Form_Activate() Call RTB.SetOptionFlag(eoSmartTabs, False) 'Private Sub Form_Activate() Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
Image: Private errLine As Integer Private wngLine As Integer (highlight a warning) Private Bkpt() As Integer (breakpoints line numbersEnd SubPrivate sub Form_Resize() (breakpoints line numbers/ Private Sub Form_Resize() (breakpoints line numbersPublic Sub Init() (DESCRIPTION: initialises this module./ Private Sub Form_Resize() (breakpoints line numbersPublic Sub Init() (DESCRIPTION: initialises this module.Private Sub Form_Resize() (breakpoints line numbers)Public Sub Init() (PTE highlighting errLine = -1 Rebim Bkpt(-1 To -1) (TTE MarksLoffset = 0) RTB_MarksLoffset = 0 RTB_MarksLoffset = 0 RTB_MarksLoffset = 0 RTB_SetOptionFlag(eoXmartTabs, False) Call RTB_SetOptionFlag(eoTrimTrailingSpaces, True)Private Sub Form_Activate()Private Sub Form_Activate()/
<pre>Private errLine As Integer 'highlight an error Private wngLine As Integer 'highlight a warning Private Sub Form_Resize() ' Public Sub Init() ' DESCRIPTION: initialises this module. ' Public Sub Init() ' DESCRIPTION: initialises this module. ' Private Sub Form_Resize() ' LErr.Left = 8 LErr.Width = PnlErr.Width - 16 Public Sub Init() ' RTB highlighting errLine = -1 wngLine = -1 ReDim Bkpt(-1 To -1) ' RTB MarksZOffset = 0 RTB.MarksZOffset = 0 RTB.MarksZOffset = 0 RTB.MarksZOffset = 0 RTB.SetOptionFlag(eoXmartTabs, False) Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True)</pre>
<pre>Private wngLine As Integer 'highlight a warning Private Bkpt() As Integer 'breakpoints line numbers ' Resizes all controls on the form in order ' Resizes all controls on the form size. ' ' Resizes all controls on the form size. ' ' Resizes all controls on the form in order ' Resizes all controls on the form size. ' ' Resizes all controls on the form in order ' Resizes all controls on the form in order ' Resizes all controls on the form size. ' ' Resizes all controls on the form in order ' Resizes all controls on the form in order ' Resizes all controls on the form in order ' Resizes all controls on the form size. ' ' Resizes all controls on the form in order ' Resizes all controls on the form size. ' ' ' Resizes all controls on the form size. ' ' Resizes all controls on the form size. ' ' Resizes all controls on the form size. ' ' Resizes all controls on the form size. ' ' Resizes all controls on the form size. ' ' Resizes all controls on the form. ' ' Resizes all controls on the f</pre>
<pre>Private Bkpt() As Integer breakpoints line numbers Public Sub Init() DESCRIPTION: initialises this module. Private Sub Form_Resize() LETr.Left = 8 LETr.Width = PnlETr.Width - 16 Public Sub Init() Private Sub Form_Resize() LETr.Width = PnlETr.Top RTB.Top = 8 RTB.Height = PnlCode.Height - 16 RTB.MarksLeftMargin = 2 Call RTB.SetOptionFlag(eoXmartTabs, False) Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate() Private Sub Form_Activate() Private Sub Form_Activate() Private Sub Form_Activate()</pre>
<pre>' Resizes all controls on the form in order ' to allow for adjustable form size. ' ' to allow for adjustable form size. ' ' ' ' to allow for adjustable form size. ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' '</pre>
<pre>' to allow for adjustable form size. ' ' DESCRIPTION: initialises this module. ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' '</pre>
<pre>Public Sub Init() Private Sub Form_Resize() LErr.Left = 8 LErr.Width = PnlErr.Width - 16 Public Sub Init() Public Sub I</pre>
Private Sub Form_Resize() LErr.Left = 8 LErr.Width = PnlErr.Width - 16 Public Sub Init() Private Sub Form_Resize() LErr.Left = 8 LErr.Width = PnlErr.Width - 16 Public Sub Init() Private Sub Form_Resize() LErr.Width = PnlErr.Width - 16 Public Sub Init() Private Sub Form_Resize() RTB highlighting errLine = -1 ReDim Bkpt(-1 To -1) RTB Expoperties RTB.MarksLofftMargin = 2 Call RTB.SetOptionFlag(eoAutoIndent, True) Call RTB.SetOptionFlag(eoSmartTabs, False) Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
<pre>LET.Left = 8 LET.Left = 8 LET.Width = PnlET.Width - 16 Public Sub Init()</pre>
Public Sub Init() IERT. Width = PnlErr.Width - 16 Public Sub Init() PnlCode.Height = PnlErr.Top RTB highlighting PnlCode.Height = PnlErr.Top errLine = -1 RTB.Top = 8 wngLine = -1 RTB.Height = PnlCode.Height - 16 REDim Bkpt(-1 To -1) RTB.Width = PnlCode.Width - 8 - RTB.Left 'RTB properties End Sub RTB.MarksLoffset = 0 'Private Sub Form_Activate()' Call RTB.SetOptionFlag(eoXmartTabs, False) 'Private Sub Form_Activate() Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
Public Sub Init() PnlCode.Height = PnlErr.Top "RTB highlighting PnlCode.Height = PnlErr.Top errLine = -1 RTB.Top = 8 wngLine = -1 RTB.Height = PnlCode.Height - 16 ReDim Bkpt(-1 To -1) RTB.TB.Width = PnlCode.Width - 8 - RTB.Left 'RTB properties End Sub RTB.MarksLeftMargin = 2 '
'RTB highlighting PnlCode.Height = PnlErr.Top errLine = -1 RTB.Top = 8 wngLine = -1 RTB.Top = 8 ReDim Bkp(-1 To -1) RTB.Height = PnlCode.Height - 16 'RTB properties RTB.WarksLeftMargin = 2 Call RTB.SetOptionFlag(eoXmartTabs, False) '' Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
errLine = -1 RTB.Top = 8 wngLine = -1 RTB.Top = 8 ReDim Bkpt(-1 To -1) RTB.Height = PnlCode.Height - 16 RTB.MarksLeftMargin = 2 RTB.MarksLeftMargin = 2 Call RTB.SetOptionFlag(eoAutoIndent, True) / Private Sub Form_Activate() ' Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
wngLine = -1 RTB.Height = PnlCode.Height - 16 ReDim Bkpt(-1 To -1) RTB.Width = PnlCode.Width - 8 - RTB.Left IRTB properties End Sub RTB.MarksXOffset = 0 Image: Call RTB.SetOptionFlag(eoAutoIndent, True) Call RTB.SetOptionFlag(eoSmartTabs, False) '' Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
ReDim Bkpt(-1 To -1) RTB.Width = PnlCode.Width - 8 - RTB.Left 'RTB properties End Sub RTB.MarksZoffset = 0
'RTB properties End Sub RTB.MarksXOffset = 0 ''''''''''''''''''''''''''''''''''''
RTB.MarksXOffset = 0 RTB.MarksLeftMargin = 2 Call RTB.SetOptionFlag(eoAutoIndent, True) 'Private Sub Form_Activate() Call RTB.SetOptionFlag(eoSmartTabs, False) Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
Call RTB.SetOptionFlag(eoAutoIndent, True) ' Private Sub Form_Activate() ' Call RTB.SetOptionFlag(eoSmartTabs, False) '' Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
Call RTB.SetOptionFlag(eoAutoIndent, True) ' Private Sub Form_Activate() ' Call RTB.SetOptionFlag(eoSmartTabs, False) '' Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
Call RTB.SetOptionFlag(eoSmartTabs, False) '' Call RTB.SetOptionFlag(eoTrimTrailingSpaces, True) Private Sub Form_Activate()
RTB.TabWidth = 6 RTB.SetFocus
Call RTB.MarksCreateImgList(AppDir + "gutter.bmp", 16, 16) End Sub
'Add one mark for instruction pointer and error pointer
Call RTB.MarksAdd(0, 0, False)
End Sub ' Public Sub Update() '
'' ' Highlights current execution point. '
' Private Sub Form Unload(Cancel As Integer) '
Private Sub Form_onload(cancer As Integer)
' DESCRIPTION: Event handler for Form Unload ' InvalidateRTB
' unloads the form if the application is ' End Sub
introdus clie form in clie application is interesting down, and just hide ste

Listings





If Not GetFilename(False, s, "", "Assembly language files |*.asm |All files |*.*", "asm", "Import assembly language program") Then Exit Sub 'Load as RTF if ext is rtf, text otherwise SetPanelText "Loading file... Please wait" Open s For Binary As 1 s = String(LOF(1), "")Get 1, , s RTB.Text = s Close SetPanelText "" 'Need to compile before running Proj.P.CompileNeeded = True End Sub Private Sub MIExport_Click() 'Get file name Dim s As String: s = "" If Not GetFilename(True, s, "", "Assembly language files|*.asm|All files|*.*", "asm", "Export assembly language program") Then Exit Sub 'Save file Open s For Binary As 1 Put 1, , RTB.Text Close End Sub Private Sub MIRun_Click(Private Sub MIRun_Click() If Proj.Running Then If Proj.Paused Then MIStart.Caption = "Continue" Else MIStart.Caption = "Pause" End If Else MIStart.Caption = "Start" End If MIStep.Enabled = Not Proj.Running Or Proj.Paused End Sub Public Sub MIStart Click() If Proj.Halted Then Proj.Running = False Proj.Paused = False Proj.Halted = False End If If Proj.Running Then If Not ConfirmCompile Then Exit Sub Proj.Paused = Not Proj.Paused If Not Proj.Paused And fiMain.MISpdMax.Checked Then Call RTB.MarksSetVisible(0, False) Else fiMain.ResetAll If Not ConfirmCompile Then Exit Sub TransferBreakpoints Proj.Running = True Proj.Paused = False End If RTB.ReadOnly = True fsCode.SetFocus fiMain.SetToolbarButtons End Sub Public Sub MIStep_Click(If Proj.Halted Then Proj.Running = False Proj.Paused = False Proj.Halted = False End If If Proj.Running Then If Not ConfirmCompile Then Exit Sub Step Else fiMain.ResetAll

If Not ConfirmCompile Then Exit Sub

TransferBreakpoints Proj.Running = True Proj.Paused = True Step RTB.ReadOnly = True End If fiMain.UpdateAll fsCode.SetFocus fiMain.SetToolbarButtons End Sub	<pre>'Check error messages If UBound(Proj.P.ErrL.sError) = -1 Then</pre>
Public Sub MIReset_Click() fiMain.ResetAll fiScode.SetFocus fiMain.SetToolbarButtons End Sub	<pre>Private Sub MINextErr_Click() Private Sub MINextErr_Click() If LErr.ListCount = 0 Then Exit Sub If LErr.ListIndex >= LErr.ListCount - 1 Then LErr.ListIndex = 0 Else LErr.ListIndex = LErr.ListIndex + 1</pre>
<pre>' Private Sub MIBreakpoint_Click() ' '' Private Sub MIBreakpoint_Click() Call RTB_OnGutterClick(0, 0, RTB.CaretY) End Sub</pre>	End If LErr_DblClick End Sub '' ' Public Sub MITick_ForFIMAIN() ' '' Public Sub MITick_ForFIMAIN()
<pre>'Private Sub MISyntax_Click() ' ' Private Sub MISyntax_Click() 'Try to compile PrgCompile 'Check error messages</pre>	'fiMain form needs to invoke Tick, but it has no access to 'some important functions which are private fsCode functions. If Proj.Halted Then Proj.Running = False Proj.Paused = False Proj.Halted = False End If
<pre>If UBound(Proj.P.ErrL.sError) = -1 Then</pre>	<pre>If Proj.Running Then If Not ConfirmCompile Then Exit Sub Tick Else fiMain.ResetAll If Not ConfirmCompile Then Exit Sub TransferBreakpoints Proj.Running = True</pre>
End Sub ' ' Private Sub MIDoCompile_Click() ' Private Sub MIDoCompile_Click() ' Try to compile PrgCompile	Proj.Paused = True Tick RTB.ReadOnly = True End If fiMain.UpdateAll fsCode.SetFocus fiMain.SetToolbarButtons End Sub

22.18. fsRegs

🜌 Registers 🔀		
F/E: Exec (3)		
A: 0000h		
B: 0000h C: 0000h D: 0000h E: 0000h		
PC: 0000h SP: 0000h		
FLAGS:		
Z: 0 S: 0 0: 0 C: 0 P: 0 N: 0 I: 0		

Option Explicit

Public declarations in this module:



Stores the string values for all registers displayed

' on this form to track changes and highlight respectively	LFP.ForeColor = IIf(LastStr(13) = LFP.Caption, 0, &HFF) LFN Exceller = IIf(LastStr(14) = LFN Caption 0, SHFF)
Private LastStr(0 To 15) As String	LFN.ForeColor = IIf(LastStr(14) = LFN.Caption, 0, &HFF) LFI.ForeColor = IIf(LastStr(15) = LFI.Caption, 0, &HFF)
' Public Sub Init() '	'Remember last status SaveLast End Sub
' NOTES: Must be called *after* fhCPU.Init '	' Private Sub SaveLast()
Public Sub Init() 'Set initial values and colors	' Saves the state of all register in order to '
Update SaveLast	' highlight them as they change. Is called by ' ' Update() after getting new values for them '
Update End Sub	Private Sub SaveLast()
··	LastStr(0) = LFE.Caption LastStr(1) = LA.Text
' Private Sub Form_Unload(Cancel As Integer) '	LastStr(2) = LB.Text LastStr(3) = LC.Text
' DESCRIPTION: Event handler for Form_Unload ' ' unloads the form if the application is '	LastStr(4) = LD.Text LastStr(5) = LE.Text
<pre>' really shutting down, and just hides the ' ' form in case the user requested to close '</pre>	LastStr(6) = LIP.Text LastStr(7) = LSP.Text
it	LastStr(8) = LFLAGS.Text LastStr(9) = LFZ.Caption
Private Sub Form_Unload(Cancel As Integer) If Not Appp.Terminating Then	LastStr(10) = LFS.Caption LastStr(11) = LFO.Caption
Cancel = 1 fsReqs.Hide	LastStr(12) = LFC.Caption LastStr(13) = LFP.Caption
End If End Sub	LastStr(14) = LFN.Caption LastStr(15) = LFI.Caption
11	End Sub
Public Sub Update()	Private Sub LA_KeyPress(KeyAscii As Integer) If KeyAscii = 13 Then
' Updates the contents of the window to ' ' reflect changes to the state of the '	If IsFmt16(LA.Text) Then Proj.CPU.A = Fmt2Dec16(LA.Text)
' simulation. '	Call Update Else
Public Sub Update()	Call MsgBox("Please enter a number between -32768 and 65535, in decimal, hexadecimal or binary. Hexadecimal numbers must be
NRun.Visible = Not Proj.Running NCtl.Visible = Proj.Running	followed by letter 'h', binary numbers - by 'b'.", vbExclamation + vbOKOnly)
Dim s As String	End If End If
<pre>s = IIf(Proj.CPU.Fetch, "Fetch (", "Exec (")</pre>	End Sub
<pre>If Proj.CPU.Fetch Then If Proj.CPU.CIB = "" Then s = s + "?"</pre>	Private Sub LB_KeyPress(KeyAscii As Integer) If KeyAscii = 13 Then
<pre>S - S + : Else S = S + IIf(Proj.CPU.fremMem, CStr(Proj.CPU.FREM + 1) + "+",</pre>	If Refrection 1 men If IsFmtl6(LB.Text) Then Proj.CPU.R(0) = Fmt2Decl6(LB.Text)
CStr(Proj.CPU.FREM)) End If	Call Update Else
Else s = s + CStr(UBound(Proj.CPU.DIB) - Proj.CPU.eDP + 1)	Call MsgBox("Please enter a number between -32768 and 65535, in decimal, hexadecimal or binary. Hexadecimal numbers must be
End If LFE.Caption = s + ")"	followed by letter 'h', binary numbers - by 'b'.", vbExclamation + vbOKOnly)
LA.Text = Dec2Fmt16(Proj.CPU.A, Proj.NmbRep) LB.Text = Dec2Fmt16(Proj.CPU.R(0), Proj.NmbRep)	End If End If
LC.Text = Dec2Fmt16(Proj.CPU.R(1), Proj.NmbRep) LD.Text = Dec2Fmt16(Proj.CPU.R(2), Proj.NmbRep)	End Sub
LE.Text = Dec2Fmt16(Proj.CPU.R(3), Proj.NmbRep) LIP.Text = Dec2Fmt16(Proj.CPU.IP, Proj.NmbRep)	Private Sub LC_KeyPress(KeyAscii As Integer) If KeyAscii = 13 Then
LSP.Text = Dec2Fmt16(Proj.CPU.SP, Proj.NmbRep) LFLAGS.Text = Dec2Fmt16(Proj.CPU.FLAGS, Proj.NmbRep)	<pre>If IsFmt16(LC.Text) Then Proj.CPU.R(1) = Fmt2Dec16(LC.Text)</pre>
LFZ.Caption = IIf((Proj.CPU.FLAGS And 1) > 0, "1", "0") LFS.Caption = IIf((Proj.CPU.FLAGS And 2) > 0, "1", "0")	Call Update Else
LFO.Caption = IIf((Proj.CPU.FLAGS And 4) > 0, "1", "0") LFC.Caption = IIf((Proj.CPU.FLAGS And 8) > 0, "1", "0")	Call MsgBox("Please enter a number between -32768 and 65535, in decimal, hexadecimal or binary. Hexadecimal numbers must be
LFI.Caption = IIf((Proj.CPU.FLAGS And 16) > 0, "1", "0") LFN.Caption = IIf((Proj.CPU.FLAGS And 256) > 0, "1", "0")	<pre>followed by letter 'h', binary numbers - by 'b'.", vbExclamation + vbOKOnly)</pre>
LFP.Caption = $IIf((Proj.CPU.FLAGS And 512) > 0, "1", "0")$	End If End If
<pre>'Colors LFE.ForeColor = IIf(LastStr(0) = LFE.Caption, 0, &HFF)</pre>	End Sub
LA.ForeColor = IIf(LastStr(1) = LA.Text, 0, &HFF) LB.ForeColor = IIf(LastStr(2) = LB.Text, 0, &HFF)	Private Sub LD_KeyPress(KeyAscii As Integer) If KeyAscii = 13 Then
LC.ForeColor = IIf(LastStr(3) = LC.Text, 0, &HFF) LD.ForeColor = IIf(LastStr(4) = LD.Text, 0, &HFF)	<pre>If IsFmt16(LD.Text) Then Proj.CPU.R(2) = Fmt2Dec16(LD.Text)</pre>
LE.ForeColor = IIf(LastStr(5) = LE.Text, 0, &HFF) LIP.ForeColor = IIf(LastStr(6) = LIP.Text, 0, &HFF)	Call Update Else
LSP.ForeColor = IIf(LastStr(7) = LSP.Text, 0, &HFF) LFLAGS.ForeColor = IIf(LastStr(8) = LFLAGS.Text, 0, &HFF)	Call MsgBox("Please enter a number between -32768 and 65535, in decimal, hexadecimal or binary. Hexadecimal numbers must be
<pre>LFZ.ForeColor = IIf(LastStr(9) = LFZ.Caption, 0, &HFF) LFS.ForeColor = IIf(LastStr(10) = LFS.Caption, 0, &HFF)</pre>	followed by letter 'h', binary numbers - by 'b'.", vbExclamation + vbOKOnly)
LFO.FORECOLOR = IIf(LastStr(1) = LFO.Caption, 0, &HFF) LFO.ForeCoLor = IIf(LastStr(12) = LFO.Caption, 0, &HFF)	End If End If

End Sub

```
Private Sub LE_KeyPress(KeyAscii As Integer)
If KeyAscii = 13 Then
     If IsFmt16(LE.Text) Then
        Proj.CPU.R(3) = Fmt2Dec16(LE.Text)
        Call Update
     Else
       Call MsgBox("Please enter a number between -32768 and 65535, in
decimal, hexadecimal or binary. Hexadecimal numbers must be followed
by letter 'h', binary numbers - by 'b'.", vbExclamation + vbOKOnly)
     End If
  End If
End Sub
Private Sub LSP_KeyPress(KeyAscii As Integer)
If KeyAscii = 13 Then
     If IsFmt16(LSP.Text) Then
        Proj.CPU.SP = Fmt2Dec16(LSP.Text)
        Call Update
     Else
       Call MsgBox("Please enter a number between -32768 and 65535, in
decimal, hexadecimal or binary. Hexadecimal numbers must be followed
by letter 'h', binary numbers - by 'b'.", vbExclamation + vbOKOnly)
     End If
  End If
End Sub
```





Option Explicit

'Font to draw string grid Private drwFont As Long 'Brush for selected variable Private brushSel As Long

Prevent specific events Private BlockSelectCell As Boolean Private BlockSetEditCell As Boolean

Public Sub Init() DESCRIPTION: initialises this module.

Public Sub Init() 'SG font drwFont = CreateFont(-11, 0, 0, 0, 400, False, False, False, 1, 0, 0, 0, "Courier New" Ο, Dim LB As LOGBRUSH LB.lbColor = GetSysColor(COLOR_HIGHLIGHT): LB.lbHatch = 0: LB.lbStyle = 0 brushSel = CreateBrushIndirect(LB) 'Set SG options SG. Option (goEditing) = True SG.Option(goColMoving) = False SG.**Option**(goColSizing) = True SG.Option(goRangeSelect) = False SG.Option(goRowMoving) = False SG.Option(goRowSizing) = False SG.Option(goThumbTracking) = True 'SG visual SG.ColCount = 2 SG.RowCount = 20

```
SG.FixedCols = 0
```

```
SG.FixedRows = 1
```

Private Sub LIP_KeyPress(KeyAscii As Integer)
If KeyAscii = 13 Then
If IsFmt16(LIP.Text) Then
Proj.CPU.IP = Fmt2Dec16(LIP.Text)
Call fimin.UpdateAll
Else
Call MsgBox("Please enter a number between -32768 and 65535,
in decimal, hexadecimal or binary. Hexadecimal numbers must be
followed by letter 'h', binary numbers - by 'b'.", vbExclamation +
vb0KOnly)
End If
End If
End Sub
Private Call MSDBOX("E Decr(M. Decime a External)

Private Sub LFLAGS_KeyPress(KeyAscii As Integer)
If KeyAscii = 13 Then
If IsFmtl6(LFLAGS.Text) Then
Proj.CPU.FLAGS = Fmt2Decl6(LFLAGS.Text)
Else
Call MsgBox("Please enter a number between -32768 and 65535,
in the integration of the second sec

in decimal, hexadecimal or binary. Hexadecimal numbers must be
followed by letter 'h', binary numbers - by 'b'.", vbExclamation +
vbOKOnly)
End If

```
End If
End Sub
```

SG.Cells(0, 0) = "Variable" SG.Cells(1, 0) = "Value" SG.ColWidths(0) = 100 SG.ColWidths(1) = 70 BlockSelectCell = False BlockSetEditCell = False

Update End Sub









Private Sub SG_OnMouseDown(ByVal MouseButton As StringGridVBProj.TxMouseButton) If MouseButton = mbRight Then PopupMenu MIOptions End Sub



23. Assembly language testing

In this section I will first test different types of operands using the 1d instruction. I will then test the most frequently used opcodes.

Note that this section provides no visual proof that tests were passed – mainly because there are too many tests, and supporting each one with a picture would be very excessive. I will support my tests in the other two sections, which will at the same time partly support these tests too.

23.1. Operand testing

No.	Instruction	Expected result	Passed
1	ld a, 10	10 loaded into accumulator	OK
2	ld a, 10b	2 loaded into accumulator	OK
3	ld a, 0FFFFh	65535 loaded into accumulator	OK
4	ld a, -8000h	-32768 loaded into accumulator	OK
5	ld a, 66000	Error message – overflow in 16-bit constant	OK
6	ld a, 0FF0	Error message – incorrect syntax	OK
7	ld b, 0F00h	3840 loaded into B register	OK
8	ld a, b	Contents of B copied into accumulator	OK
9	ld a, myvar	Contents of myvar loaded into accumulator	OK
10	ld myvar, a	Accumulator stored in myvar	OK
11	ld myvar, 10	Error message – cannot load a constant into a variable	OK
12	ld c, myvar ;myvar is not declared	Error message – myvar is not declared	OK

23.2. Opcode testing

No.	Instruction	Expected result	Passed
13	add a, 10	10 added to accumulator; result stored in accumulator	OK
14	sub b, a	A subtracted from B; result stored in B	OK
15	cmp a, b ; $a = b$	Zero flag set to true	OK
16	mul a, b	A multiplied by B; result stored in A	OK
17	mul b, c	Error message – incompatible operands	OK
18	div a, myvar	A divided by myvar; result stored in A	OK
19	neg c	Sign of C changed	OK
20	not c	All bits in C changed	OK
21	and a, myvar	Bitwise And performed on A and myvar; result in A	OK
22	or b, a	Bitwise Or on B and A; result stored in B	OK
23	xor b, a	Bitwise Xor performed on B and A; result stored in B	OK
24	xor a,a	Accumulator set to zero	OK
25	lshr b, 5	Bits in B shifted by 5 bits to the right	OK
26	lshl a, 3	Bits in A shifted by 3 bits to the left	OK
27	ashr a, 5	Bits in A shifted by 5 bits to the right; sign preserved	OK
28	ashl b, a	Bits in B shifted by A bits to the right; sign preserved	OK
29	ashl b, 20	Error message – number of shifts exceeds 15	OK
30	jg~lbl~; after comparing 2 and 5	Jump not performed	OK
31	jg~lbl~; after comparing 5 and 2	Jump performed	OK
32	$jl \ lbl$; after comparing 2 and 5	Jump performed	OK
33	jmp mylabel	Jump to mylabel	OK

34	call myproc	Myproc called; previous address pushed on stack	OK
35	ret ; after a call	Jump to the instruction following the call opcode	OK
36	halt	Program execution stopped	OK

24. Window testing

24.1. Code window

No.	Test	Expected result	Passed	Reference
37	Write a program	Syntax highlighting; editing facilities	OK	24.1.1
38	Step through a program	Current line highlighted	OK	24.1.2
39	Set a breakpoint	Breakpoint line highlighted; execution should pause	OK	24.1.3
40	Run a program with	Error listed at the bottom; The first error displayed	OK	24.1.4
	many errors	in an error message.		





24.2. CPU window

No.	Test	Expected result	Passed	Reference
41	Run a program	Register contents should be displayed, including	OK	24.2.1
		Current Instruction Register and MAR/MDR		
42	Switch complexity mode	Window layout should change – to a simple one in	OK	24.2.2
		this case		







24.3.	RAM	window

No.	Test	Expected result	Passed	Reference
43	Run a program	Program code should be displayed and highlighted	OK	24.3.1
		with yellow. Current instruction should be		
		highlighted with blue		
44	Edit a cell	Memory byte must be updated; the new byte must	OK	24.3.2
		be highlighted with red		



24.3.2

24.4. Registers window

No.	Test	Expected result	Passed	Reference
45	Run a program	Register values should be displayed. Registers that		24.4.1
		change must be highlighted with red.		
46	Edit a register	Register value must be updated	OK	N/A
47	Change number format	Values must be shown in the new number format	OK	24.4.2
	_	(decimal unsigned in this case)		



24.5. Variables window

No.	Test	Expected result	Passed	Reference
48	Run a program	All variables declared in the program must be	OK	24.5.1
		listed in the grid.		
49	Edit a variable	Variable value must be updated	OK	N/A



24.6. Stack window

No.	Test	Expected result	Passed	Reference
50	Push a value onto stack	The value should be displayed in the grid; stack		24.6.1
		pointer (red) must move down		
51	Return from a call	Stack pointer must move up	OK	N/A



cli

25. Overall testing

25.1. Print keys program

This program installs a keyboard interrupt service procedure which prints every key that the user presses on the screen. The source code is as follows:

```
;Install interrupt handler
      ld a,offset(isp_kbd)
      ld [0FF02h],a
      ;Initialise screen
      in b,52h ;get video memory offset
out 50h,2 ;color text mode
out 54h,0 ;manual refresh
      ;Variables
      ld d, 0 ;number of chars read
ld e, 0 ;exit flag - nonzero to terminate
      sti
      ;Main loop
lp:
    xor a,a
              ;test exit flag
;continue if zero
     cmp e,a
      jz lp
      ;Terminate program
     halt
;----;
;--- ISP for keyboard interrupts ---;
;-----
       -----;
isp_kbd:
     push a
     push c
     ;Get pressed key code
     in c,60h
      ;Check if it is any of special keys
      ld a,c
      cmp a,51
                                    ;these are after
      jz key_exit
      cmp a,50
                                    ;this procedure to
      jz key_cls
                                   ;make this clearer
continue:
      ;Output key to screen
      ld c,[c*1+offset(kbd_xlat)] ;get key symbol
      lshr c,8
                                    ;remove second symbol - have read a word
      ld a,c
      ashl a,8
                                    ;char
      add a,0F0h
                                    ;color
      ld c,a
      ld [b+d*2],c
                                    ;print char
     out 54h,2
                                    ;refresh screen
      inc d
                                    ; increment printed num
      ;This is the end
exit_isp_kbd:
     pop c
      pop a
      iret
key_exit:
      ld e,1
      jmp exit_isp_kbd
key_cls:
     ;out FUNCTION NOT SUPPORTED YET BY VIDEOCARD
      jmp continue
```

The program is started, and the input is a sequence of keys pressed to type, "testing keyboard and screen.". The program worked fine, and the ouput was:



25.2. Factorial program

This is an assembly language version of the following high-level language code:

```
Result = Factorial(Number)
END
Sub Factorial(n)
   if n=1 then return 1 and exit sub
   Factorial = Factorial(n-1)*n
Exit Sub
```

This is a rather ineffective recursive implementation of the algorithm, but it helps to show why stack is needed for recursive procedures. It is also a good test of how well stack works. This is the code:

```
;Run Factorial with parameter Number
ld a,Number
call Factorial
st a,Result
halt
```

```
Factorial:
      ;Check if we need to exit
      cmp a,1
      jz Factorial_done
      ;We need to preserve A - otherwise recursion
      ; will modify it incorrectly
      push a
      ;Call Factorial(a-1)
      dec a
      call Factorial
      ;Get old parameter into b
      pop b
      ;Multiply the two numbers and return the result
      mul a,b
Factorial_done:
      ret
Number:
            dw 5
Result:
            dw ?
```

If the Number variable is set to 5, the output is 120, as expected (picture 25.2.1). If Number is 8, Result becomes 40320 (picture 25.2.2).



🛚 Variables		X	
Variable	Value		
Number	8		
Result	40320		
	ł		,

Maintenance

26. Introduction

This section describes the way the program functions in order to allow any programmer to correct minor bugs, or make modifications to the program.

It is important to understand that a lot was written about the way the system functions in the Design part. In this section, I will assume that the reader has the knowledge of the terms and algorithms described in that part.

27. Organisation and conventions

27.1. Modules

The program is composed of several modules, each of which has a specific function. There are two kinds of modules – procedural (*.bas) and form (*.frm). Procedural module names start with a "p", so in general procedural modules are named p*.bas. Form modules are subdivided into four categories – hardware modules (fh*.frm), device modules (fd*.frm), interface modules (fi*.frm) and system modules (fs*.frm). The functions of these modules are as follows:

- *Procedural* contain procedures, functions, types and variables related to a particular process rather than to a form.
- *Hardware* contain code related to a particular piece of computer hardware. These modules, unlike *device* modules, provide user with a way to view component's operation and edit its contents if any. Operation of this hardware is simulated in procedural modules.
- *Device* each such module describes interface *as well as* operation of a device. Device modules exist for all external devices (i.e. all devices which communicate with processes through input/output ports).
- *Interface* modules which simulate peripherals such as keyboard or screen. These devices are those that a user sees when sitting at a PC.
- *System* these provide development windows such as code development window, register & variables contents windows etc. These windows could be thought of as an IDE (Integrated Development Environment).

Below is a list of all modules with a short description:

Procedural

- **pGlobals** defines application entry point, declares two global variables *Proj* and *Appp* (discussed later), defines window procedure wrapper.
- **pWinAPI** imports Windows API functions.
- **pUtils** implements a set of general functions not available in VB, such as conversion between different number representations etc.
- **pCompile** contains local variables and procedures required to compile a program into machine code.

- **pExec** contains local variables and code required to execute machine code. Exports functions to perform one clock tick, execute one machine instruction etc.
- **pIO** provides general functions for interfacing device modules, such as PortWrite, for simple operations with multiple external devices.

Hardware

- **fhCPU** code associated with CPU window and a little bit with CPU itself.
- **fhCU** code associated with control unit interface window.
- **fhRAM** code associated with RAM interface window.

Device

- **fdKeyboard** code that determines the way keyboard controller works, plus keyboard controller user interface.
- **fdSpeaker** code that determines the way speaker controller works, plus speaker controller user interface.
- **fdVideo** code that determines the way video controller works, plus video controller user interface.

Interface

- **fiSplash** the splash screen window, displayed at startup while the system is loading.
- **fiMain** main window at the top of the screen.
- **fiComp** computer window which contains a screen, a keyboard, a speaker and buttons to start/stop program execution.
- **fiKeyboard** detachable keyboard.
- **fiDisplay** detachable and scalable display.
- **fiSpeaker** detachable speaker.

System

- **fsCode** module providing code design and debug environment, with the primary feature of code editing.
- **fsRegs** displays and allows modification of internal CPU registers.
- **fsVars** displays and allows modification of variables declared in the code.
- **fsStack** displays stack contents with stack pointer.

27.2. Visibility and naming conventions

In VB it is impossible to specify different visibility for the same declaration as viewed from different modules. If something is public then is will be visible in *all* other modules, regardless of whether they want to see it. Form modules are a special case because public procedures and variables need a qualifier (module name), but will be visible in all modules anyway, and neither public types nor arrays are allowed as public members of form modules. Therefore, I tried to minimize the number of public declarations by trying to group code in such a way that some declarations are not required outside the module. Also some naming conventions are adopted to increase code readability.

Procedures and functions

Procedures and functions are always declared with an explicit visibility modifier to make the code easier to understand. All procedures and functions in putils and pWinAPI are public. Otherwise they are only public if they provide some service required globally, such as the public Tick function in the pExec module. All event handlers are private because that's a VB convention. The only naming convention used is for initialisation procedure for public modules. It is prefixed with cmp for pCompile, exe for pExec and io for pIO.

Variables

All global variables are organised into two structures – **Proj** and **Appp**, discussed below. No other global variables should be declared without a good reason, and if possible such variables should be organised into global structures like **Proj** and **Appp**. No naming convention is required as global variables are already grouped into structures, and all other variables are private and thus easier to manage.

Constants

All constants in this system are local. Some of them are prefixed to group them, but no global naming convention is used.

Types

It is quite annoying that one has to declare all structured types separately in VB. So if I want to create substructures inside a structure I need to declare a separate type and name it. To simplify understanding all such "hidden but actually visible" types will be prefixed with Tp. The Proj and Appp types themselves are also considered "hidden" because the only reason they are delcared is to declare one single variable of the type. Some types which are part of the Proj structure will be declared as public because they will be used for local variables in some modules. It seems logical to me to "hide" them as well by prefixing them with Tp. All other types will be prefixed with T and declared in the pglobals module.

28. Global data structures

There are only two global data structures -Proj and Appp. Proj stores stores all data associated with the current project and simulation state, whereas Appp is for data that relates to some housekeeping tasks, e.g. store a flag to indicate that application needs to terminate.

Note that although a lot of structures are declared in different modules, structures will not be discussed at all in the Modules section below.

28.1. Proj structure

Proj is a big structure, containing a lot of substructures. This structure is declared in pGlobals. Its type is called TpProj, it is declared as local structure type in pGlobals. Below is a list of all members with their types and a short description.

Project related

- Modified: Boolean true if project has been modified and needs to be saved.
- **Complexity: Integer** global complexity setting. 0 for GCSE complexity, 1 for A-level, 2 for full complexity. 0 by default.
- **NmbRep: Integer** global number representation setting. 0 for hexadecimal, 1 for binary, 2 for decimal unsigned, 3 for decimal signed. 0 by default.

Program related

• **P: TpPrg** – stores program written by the user, compiled program and some additional information related to compilation. See below for more details.

Execution related

- **Running: Boolean** true if simulation is on.
- **Paused: Boolean** true if simulation has been paused. Should not be true if Running is not true.
- Halted: Boolean true if simulation has been halted. Running will be false, but windows will still display all data. Otherwise windows would say that user has to start the program before using them.
- **TickCount: Long** number of clock ticks executed since last reset.
- **CPU: TpCPU** holds CPU simulation state, such as register values etc. See below for more details.
- **RAM: Array Of Byte** holds the contents of RAM.
- Video: TpVideo holds data about video card, such as current mode, palette memory, etc. See below for more details.

28.1.1. TpPrg substructure

This structure is declared in pcompile as a global structure. See <u>algorithms</u> below to understand the meaning of some members. Some members are declared as structured types – see below for description. This structure has the following members:

• AsmLine: Array Of String – holds current program in assembly.

- **TknLine:** Array Of **TpTokenLine** holds tokenized program.
- **Code: String** holds assembled program.
- **Code_O2L:** Array Of Integer holds values to convert offset in Code to a specific line in source code.
- **Code_L2O: Array Of Integer** holds values to convert a specific line in source code to an offset in Code.
- **Ref: Array of TpRef** holds a list of all references with their addresses and code line at which they are declared.
- **Backpatch: Array of TpBackpatch** holds all requests for backpatch produced during second compilation pass.
- Vars: Array of TpVars holds a list of all variables declared in the code with their offsets in the memory and line where they were declared.
- ErrL: TpErrLog a structure to hold all error and warning messages produced during compilation.
- **CompileNeeded: Boolean** true if user changed the source code and the program needs to be compiled again.

28.1.2. TpCPU substructure

This structure holds CPU simulation state, such as register values etc. It contains the following elements:

General-purpose registers

- A: Long the accumulator
- R: Array (0..3) of Long general-purpose registers stored in such a way that B is R(0), C is R(1) etc.

Special-purpose registers

- **IP: Long** instruction pointer points to next instruction to be executed.
- **SP:** Long stack pointer points to next free stack element.
- **FLAGS: Long** flags register.

Internal registers

- MAR: Long Memory Address Register.
- MDR: Long Memory Data Register.
- **CIB: String** Current Instruction Buffer. Holds fetched instruction, each character in the string representing one fetched byte.
- **DIB:** Array of **TpDI** Decoded Instruction Buffer. Holds decoded microprogram. See below for more detail.
- Fetch: Boolean Indicates whether the CPU is fetching or executing an instruction.
- FREM: Integer Indicates how many more bytes there are to fetch.
- **fremMem: Boolean** This register is invisible to the user, and is only here to indicate that a memory addressing should be fetched too, thus adding another 2-4 bytes to fetch. To read more about this flag, see (Fetch process) below.
- **IS: Long** Interrupt Status register.

Execution state

- **eSelectedReg:** Integer indicates the number of general-purpose register selected by the Control Unit for a read/write operation. Can be 0 to 5, respectively, for the following registers: b, c, d, e, sp, ip.
- **eIDB:** Long value held on the Internal Data Bus. Volatile between different clock cycles.
- **eIAB:** Long value held on the Internal Address Bus. Volatile between different clock cycles.
- **eDP:** Integer Decoded Instruction Pointer. Points to the next microinstruction to be executed by the Control Unit in the Decoded Instruction Buffer.
- **eEDB:** Long value on the External Data Bus (a.k.a. Data Bus). Volatile between different clock cycles.
- **eEAB:** Long value on the External Address Bus (a.k.a. Address Bus). Volatile between different clock cycles.
- **Breakpoint:** Array of Long array holding values of IP register, encountering which the execution process should stop and hand over control to the Integrated Development Environment. Has nothing to do with CPU simulation.

28.1.3. TpDI

This structure represents a single decoded microinstruction. To read more about how this structure is used, please refer to (<u>Decode process</u>) and (<u>Execute process</u>).

- Sig1: Long low-order word of control signals
- Sig2: Long high-order word of control signals
- nToIDB: Long data to be put on IDB
- nAluOpNum: Integer ALU operation number
- nJmpCond: Integer conditional jump number
- nAdrMul: Integer address multiplier for indexed addressing
- nAluSh: Integer ALU shift count for shift operations
- nIntIS: Integer unused

28.1.4. TpVideo substructure

This structure holds data about video card, such as current mode, palette memory, etc. It consists of the following elements:

- Mode: Integer mode number
- **autoUpdate: Boolean** indicates whether screens are updated automatically by CLab or manually by the user.
- mdResX: Integer horizontal screen resolution for current mode.
- mdResY: Integer vertical screen resolution for current mode.
- **mdType:** Integer mode type 0 for text, 1 for direct graphics, 2 for paletted graphics.

- mdColors: Integer color resolution for current mode 0 for monochrome, 1 for 16 colors, 2 for 256, 3 for 65536 and 4 for 16777216.
- mdFntX: Integer width of one character in screen pixels.
- mdFntY: Integer height of one character in screen pixels.
- **MemOff: Long** offset to video memory in RAM.
- PalMem: Array (0..255) of Long palette memory
- vDC: VirtualDC a "bitmap" in memory where all drawing occurs and which is then drawn on simulated screens. VirtualDC is a class written by an unknown author, with slight modifications.

28.1.5. TpToken

For more information about how this structure is used please refer to (<u>Compile</u> <u>process</u>) section.

- Text: String string from source code containing the token, e.g. "AND".
- **Type:** Integer token type, one of the tk constants, listed below.

28.1.6. Token Type constants

- **tkUnknown = -1** used during tokenization
- **tkLabel = 0** token is a label declaration
- **tkVarDecl = 3** token is a variable declaration
- **tkVarInit = 4** token is a variable initialisation
- **tkOpcode = 5** token is an opcode
- **tkOperand = 6** token is an operand

28.1.7. TpTokenLine

For more information about how this structure is used please refer to (<u>Compile</u> <u>process</u>) section.

- Token: Array of TpToken all tokens which are located on this line.
- **CodeLine:** Integer line number in the original code where this line came from.
- **CodeOffset:** Integer address of whatever code is produced by this token line.

28.1.8. TpRef

For more information about how this structure is used please refer to (<u>Compile</u> <u>process</u>) section.

- Name: String label/variable name.
- Addr: Long address that the reference points to.
- **CodeLine:** Integer line number in original code where this reference was declared.

28.1.9. TpBackpatch

For more information about how this structure is used please refer to (<u>Compile</u> <u>process</u>) section.

- Name: String name of label/variable referred to.
- Addr: Long where to write the offset of requested label/variable
- IsDW: Boolean unused
- **ReITo: Long** unused
- **CodeLine:** Integer line number in original code where this reference was requested.

28.1.10. TpVars

For more information about how this structure is used please refer to (<u>Compile</u> <u>process</u>) section.

- Name: String variable/label name
- Addr: Long variable/label address

28.1.11. TpErrLog

For more information about how this structure is used please refer to (<u>Compile</u> <u>process</u>) section.

- **sError: Array of String** error message
- IError: Array of Integer number of line producing error
- **nError: Array of String** error number
- sWarning: Array of String warning message
- IWarning: Array of Integer number of line producing warning
- **nWarning: Array of String** warning number

29. Processes

This section describes core processes in detail, looking at what happens at each stage and why it happens. This section does *not* look at helper functions used in the processes – those will be discussed later in (Main functions and procedures).

29.1. Startup

Application entry point is defined in the pGlobals module. The function is called Main, it returns no value and has no parameters. Below is a list of actions that CLab does when it starts up.

- 1. Set Running flag, get windows version and initialise XP controls if necessary.
- 2. Load and display splash screen.
- 3. Load all forms that CLab contains; this takes most of the startup time.
- 4. Initialise **Proj** structure.
- 5. Initialise all modules which require initialisation.
- 6. Show main form and computer form.
- 7. Hide and destroy splash screen.

According to VB6 help, the application will keep running after Main returns for as long as at least one form is loaded. So we do not need the message loop; VB does it all for us. Please refer to the next section to read about shutting down process.

29.2. Shutdown

To shutdown, CLab unloads all forms. This causes the message loop to stop, and the application terminates. User interface is designed in such a way that the user can only close the main window to shut down the application. Closing any other window simply causes that window to be hidden.

Visual Basic's End statement is supposed to terminate the program by unloading all forms. But apparently it does something else, because End causes CLab (and VB too) to crash, whereas manually unloading all forms in a loop works fine. Therefore, a loop is used instead of End.

When the user closes the main window, Terminating flag will be set to True, all forms will be asked about shutting down, and if they all agree then they are unloaded. This time forms *will* be unloaded (as opposed to being hidden) because of the Terminating flag.

29.3. Assembly

Assembling a program requires five stages, three of which are pure assembling, one – preparation and one – postprocessing.

Preparation

Code is copied into the Proj structure, and the structure is prepared for assembly.

Pass 1

During this pass the program is tokenized (i.e. split into tokens). First of all, the source code is cleaned by removing all comments, ensuring all sequences of tabs and spaces are replaced with single spaces, and then trimming leading and trailing spaces, if any.

Next, the program is split into tokens in the following way. For every non-empty line of source code a token line is created. The program loops through the characters of the line, accumulating them in a special variable. Whenever it encounters a space, it saves whatever it has accumulated as a token in the token line and starts accumulating next token.

Having tokenized the whole program, CLab tries to identify token types. At first it identifies *label* tokens (if there is a semicolon at the end then it is a label) and *vardecl* tokens (if token is DB, DW or DS). Everything else is identified as *unknown* for now. Next CLab analyses token positions to further identify them. This time it ignores all *label* tokens. If the first token is *unknown*, then it is identified as *opcode*, otherwise it remains as it is. Now all *unknown* tokens after an *opcode* token are identified as *varinit* tokens.

The next stage in this pass is to analyse all token patterns and see if they are valid. Some patterns may be corrected; others may not. First of all, all *label* tokens are placed in their own token lines so that there is only one label per token line and nothing else. Next, CLab checks if there is a *vardecl* token without a *varinit* token after it. If there is, it issues a warning and adds an unitialised *varinit*. Having done this, the algorithm is ready to check token patterns. There are only a few token patterns that are valid at this stage. They are:

label vardecl varinit opcode opcode operand opcode operand operand

If a given token line does not follow any of these patterns at this stage, an error message is issued, and assembly process is stopped.

The final stage of this pass is to prepare references to variables in such a way as to simplify compilation. Whenever the user wants to use memory addressing, he can either write the address of the variable or variable name, which will be replaced with the address by the compiler. The following notations are equivalent: [50h] and var, given that variable var is stored at address 50h. There are several addressing levels, which are listed below:

	Pointer	Dereference	Double dereference
Specified by address	50h	[50h]	[[50h]]
Specified by name	offset(var)	var	[var]

There is an obvious mismatch in formats – writing a variable name dereferences its address automatically, whereas writing memory address does not. Therefore, it is impossible to replace variable name by its address directly. The <code>offset()</code> syntax makes things even worse. At the end of first assembly pass the program "shifts" all addressings specified by variable name in the following way. If it encounters the structure <code>offset(x)</code>, it replaces it with the structure x. If it encounters something that must be a reference, it checks whether it is a number or a variable name. If the latter is true then it encloses the variable name with []. It is easy to see that now the table shown above will look like this:

	Pointer	Dereference	Double dereference
Specified by address	50h	[50h]	[[50h]]
Specified by name	var	[var]	[[var]]

Now all variable names can simply be replaced by their addresses. Note that the reason for all this is to simplify assembly language syntax - it would not be nice if users had to write ld a, [var] instead of ld a, var.

Pass 2

This pass generates machine code for the tokenized program. It loops through all token lines, generating code for every line and adding it to the Proj.P.Code string. In this way the current address (the address of the instruction being compiled) will be the length of Proj.P.Code string. Below is a description of what the algorithm does for every token line.

First, the algorithm checks if the first token is a label. In that case the algorithm will check if the variable name is valid (and issue an error if it isn't). Then it will add the label to the reference list (Proj.P.Ref), storing reference name, address and source code line.

Next the algorithm starts to actually convert source code into machine codes. To reduce the amount of work, all similar instructions are assembled in loops. There are arrays for every group of similar instruction, containing opcode and the corresponding machine code. The loops go through all opcodes in array, comparing them to the one that is being assembled. If the operand belongs to none of the groups, it is assembled individually.

Having assembled all token lines, this pass generates arrays which help converting between source code lines and offsets in machine code. This is mainly used for breakpoints and to highlight machine code instruction being executed.

Pass 3

This is the last assembly pass. Here the program checks if there are any multiple label definitions or any undeclared references, and then backpatches the program. The algorithm goes through the list generated during the second pass, writing the addresses of all references as requested.

The reason why a separate pass is required to backpatch is that the program does not know addresses of all references until it has finished code generation.

Postprocessing

Having assembled the program, CLab will list all errors and warnings (if any). It will also update all windows to reflect changes.

29.4. Fetch

To fetch an instruction, CLab gets the byte at current IP. It then checks if it has already fetched some bytes. If it hasn't, it will calculate and store the number of bytes to fetch, using a special array, InstructionLen, generated at startup in exeInit. Otherwise the program will simply fetch bytes and add them to CPU.CIB. If the program detects that it has finished fetching an instruction which uses a memory addressing, it will use the last fetched byte to determine the length of memory addressing which will be appended to current instruction.

As soon as the last byte is fetched, instruction will be decoded into machine codes, and the CPU will switch to execution mode.

29.5. Decode

Decoding is rather similar to code generation. Some instructions are grouped. There are arrays which hold machine codes for every group, which are compared to current machine code. As soon as instruction is identified, a microprogram which can be directly executed by the CPU is generated.

A list of all microinstructions can be found in section (<u>Microinstructions</u>). A table of microprograms for all machine codes can be found in (<u>reference</u>).

29.6. Execute

All that the execute cycle does is to take every microinstruction from Proj.CPU.DIB and do actions if a respective signal is set. There is a list describing what each signal does – see (Microinstructions).

Having executed all microinstructions, the CPU will prepare to fetch the next instruction and then run the interrupt check algorithm. It will be that algorithm that will switch to fetch if no interrupts are pending.

29.7. Interrupt

To check for pending interrupts, CLab goes through the bits of **IS** register, starting with the low-order bit, which corresponds to interrupt request 0, thus giving it highest priority. As soon as it encounters a bit set to 1, it will go and create a microprogram to initiate the interrupt. This microprogram can be found in the (Appendix).

30. Functions and procedures

30.1. pGlobals

30.1.1. Main

This is the application entry point. It initialises the whole application, loading all settings and showing relevant windows.

30.1.2. WindowProc

This is a wrapper for the real window procedure of fiMain form. The reason is that VB does not allow to get the address of *any* procedure declared in a form module. But we have to know the address of the window procedure to hook the window with the SetWindowLong WinAPI function. Thus this clumsy wrapper in this module.

30.2. pUtils

Errr

Displays an error message containing the string passed to this procedure. It is just a shorthand – this way we don't have to bother about the caption, the icon and buttons.

Tally

Counts the number of occurrences of one string in another string.

FieldStr

Returns a specified element from a list stored in a string separated by a special character.

InStrBack

The same as InStr except for the fact that it works backwards. It starts looking for occurences at the end of the string, and returns the position of the first one.

Hex2Dec

Converts a hexadecimal number to a decimal number.

Dec2Hex

Converts a decimal number to a hexadecimal number.

Bin2Dec

Converts a binary number to a decimal number.

Dec2Bin

Converts a decimal number to a binary number.
Dec2Chr

Converts a decimal number to a base 256 number, returning the result as a big-endian string.

Chr2Dec

Converts a base 256 number to a decimal number. The source number is interpreted as big-endian.

Str2Chr

Formats source string by writing each character in hexadecimal, separated with a space.

TestCharset

Tests if all characters of a given string belong to a given charset.

StringIsInt

Returns true if a given string is a decimal integer.

StringIsLong

Returns true if a given string is a decimal long integer.

GetFilename

Initiates an open or a save dialog using ComDlg functions GetSaveFileName or GetOpenFileName.

AppDir

Returns application path with a backslash at the end.

Dec2Fmt16

Converts a decimal number to one of the supported number representations.

IsFmt16

Returns true if a given string is a valid 16-bit number in one of the supported number representations.

Fmt2Dec16

Converts a number in any of the supported representations to a decimal number.

FntWrite

Writes text on a given device context using a given font. The reason for using this is that for some reason a font selected into a DC is removed from that DC after the first text output to that DC. Therefore, a font has to be selected every time.

CreateFnt

Creates a font by initialising an application-defined font structure and creating respective GDI objects.

DestroyFnt

Destroys a font created by CreateFnt by destroying respective GDI objects.

30.3. pCompile

cmpInit

Initialises the module by arrays for all opcode compilation groups....

ReadCodeIntoProj

Copies the code written by the user from the text editing control into the Proj structure.

PrgCompile

Assembles user program by reading the code into **Proj** structure, initialising some variables, running all three assembly passes, displaying all errors and warnings, and updating all windows.

PrgLoad

Loads an assembled program into RAM and updates RAM window.

CompilePass1

Assembly pass 1. Described in detail in (23.3 assembly).

CompilePass2

Assembly pass 2. Described in detail in (23.3 assembly).

CompilePass3

Assembly pass 3. Described in detail in (23.3 assembly).

CompileMemoryAddressing

This function is used in CompilePass2. Compiles a given memory addressing operand into machine code which can be added to the instruction that requires it.

OperandIsRg

Returns true if a given operand is a register (A, B, C, D or E).

OperandIsRgn

Returns true if a given operand is a general-purpose register (B, C, D or E).

OperandIsMem

Returns true if a given operand is a memory addressing.

OperandIsIm8

Returns true if a given operand is an 8-bit immediate constant.

OperandIsIm16

Returns true if a given operand is a 16-bit immediate constant. If the operand is a variable name, the function will still succeed because variable address is a known constant.

Clm8

Converts an operand into an 8-bit number. No error-checking – this function assumes OperandIsIm8 was called on the same operand to check validity.

Clm16

Converts an operand into a 16-bit number. No error-checking – this function assumes OperandIsIm16 was called on the same operand to check validity. If the operand is a variable name, the function will file a backpatch request and return 0.

AddErr

Adds a given error to the error list.

AddWng

Adds a given warning to the error list.

CleanSpaces

Converts all sequences of tabs and spaces into a single space. Used in CompilePass1.

30.4. pExec

exelnit

Initialises the module by filling an array of instruction lengths and some group decode arrays.

GFIg

Checks if a given signal in a given microinstruction is set.

SFIg

Sets a given signal in a given microinstruction. Signals are passed as separate parameters. Some signals (their names are prefixed with op_{-} for Operation) cause SF1g to interpret the following parameter as a number and save it in a special register in TpDI, depending on what op signal was used.

Tick

Executes one clock tick. This is a public function being a wrapper for the private function eTick.

Step

Executes one whole instruction, fetching the next instruction.

eTick

Executes either a fetch or an execute cycle, then calls Tick procedure for all device modules. Also increments the tick counter.

eFetch

Fetches one byte. For a detailed description see (Processes.Fetch).

eDecode

Decodes instruction in Proj.CPU.CIB. For a detailed description see (Processes.Decode).

eExecute

Executes one microinstruction. For a detailed description see (Processes.Execute).

eInterrupt

Checks for interrupts. For a detailed description see (Processes.Interrupts).

DecodeMem

Adds such microinstructions to Proj.CPU.DIB as to calculate the address specified by the memory addressing in Proj.CPU.CIB and store it in MAR.

reg_sX

Returns "select register X" signal for the signal specified by an integer (0 for B, 1 for C, 2 for D, 3 for E).

DI2Str

Generates a string with the names of all signals in a given microinstruction.

30.5. pIO

devInit

Initialises all device modules by calling respective initialisation procedures.

devReset

Resets all devices by calling respective reset procedures.

devTick

Lets all devices to do some processing every tick if there is anything they want to do.

devPortRead

Queries all devices if any of them wants to respond to a port read signal with a given address.

devPortWrite

Calls PortWrite for all devices thus simulating a port write operation.

IRQ

Devices call this function to request a given interrupt. Returns false if CPU cannot accept the interrupt. Otherwise sets "pending" flag and returns true.

30.6. Common functions (all form modules)

This section describes functions common to all form modules.

Init – Initialises the module.

Form_Unload – Either hides or unloads the form depending on whether the application is terminating or not.

Update – Changes data on the form to reflect changes to simulation state.

SaveLast – Stores register values in order to highlight them if they change.

30.7. Common functions (device modules)

Reset – initialises the device whenever the user restarts the program.

Tick – does some processing every clock tick.

PortRead – port read operation for the given device.

PortWrite – port write operation for the given device.

30.8. Other functions worth mentioning

fiMain.Hook

Hooks the main window by installing application-defined window procedure.

fiMain.Unhook

Unhooks the main window by returning the VB-defined window procedure. The main window *has* to be unhooked before unloading it, otherwise VB crashes.

fiMain.WindowProc

Window procedure for fiMain initiated via a wrapper defined in pGlobals. Traps minimize event and hides all forms. Also traps restore event and shows all forms hidden during minimization. Traps clicks in the non-client area and activates the form (VB's message procedure does not do that if the form is non-movable).

fdVideo.SetVideoMode

Sets a video mode. Video mode number is the same as the one used with port write operation.

fdVideo.UpdateScr

Repaints the video memory on the internal memory device context, ready to be blitted onto screens.

fsCode.InvalidateRTB

Invalidates every line of the code editor, causing it to repaint fully.

fsCode.DisplayError

Highlights a given error/warning in the error list by showing it in the code editor and highlighting the offending line. Optionally displays an error message with the error/warning text.

fsCode.TransferBreakpoints

Copies breakpoints from the code editor into the **Proj** structure.

31. Sample modifications

31.1. Renaming opcodes

It is very easy to rename opcodes. Opcode names are only used in pCompile module. Use text search to find the opcode name – it will be either in cmplnit (for grouped instructions) or in CompilePass2. Make sure that the new name is lowercase – assembly language is not case-sensitive, so all comparisons are made in lowercase.

31.2. Adding an instruction

Adding an instruction is not exactly straightforward. Many parts of the program will need modifications. The program was not designed with easy instruction set updatability in mind.

Assembling a new instruction

CompilePass2 needs updating. There is a big block of code inside a loop through all token lines. Inside that block there is a variable called t which contains current opcode in lowercase. You will need to add an IF block to check if the program is trying to compile the new instruction. The code inside the block will have to do all the assembling and end with a GOTO NextTokenLine statement. Whatever your code has assembled must be placed in variable ctl – it will be added to machine code automatically.

When assembling an instruction, tl will contain current token line. So if you need to check the number of parameters, use tl.Count, and tl.Token() array will contain all tokens in the token line (e.g. tl.Token(0) will return your opcode).

You may find the following functions useful when assembling your instruction. To check what type a given operand is, use OperandIsRg, OperandIsRgn, OperandIsIm8, OperandIsIm16 or OperandIsMem.

Make sure that the first byte of your machine code equivalent is unique to your instruction – otherwise CLab will have problems decoding it.

If the syntax is not correct, use AddErr to issue an error. If there is an assumption you make, and you want the user to be aware of it, use AddWng to issue a warning.

If one of your operands is a memory addressing, use the CompileMemoryAddressing function. Pass to it your operand and some backpatching information (see declaration for details), and it will return compiled memory addressing bytes which can be easily decoded with DecodeMem

Decoding a new machine code

First of all, add the first byte of your machine code to the InstructionLen array initialisation. If N is the number corresponding to the first byte of your machine code then the N^{th} element of the array must be the length of your machine code instruction.

If you use memory addressing then specify the length of the instruction without the addressing, with a minus sign.

Now you can decode the instruction. Add your code to the end of eDecode procedure. Variable b will contain the first byte of the machine code – check if that's the one you want to decode. Now redim .DIB with elements from -1 to how many machine code instructions you will need subtract 1. Call SF1g for each element of .DIB, supplying it with all signals that you want.

If you want to place a condition, use op_jmp_cond flag followed by one of the numbers from the table below. If a condition is not met, CPU will stop executing current microprogram. Possible conditions are:

0 if greater 1 if not greater (if less or equal) 2 if less 3 if not less (if greater or equal) 4 if equal (if zero) 5 if not equal (if not zero) if carry 6 7 if not carry 8 if overflow 9 if not overflow 10 if sign 11 if not sign

For example, SFlg(.DIB(0), op_jmp_cond, 6) will ensure that your microprogram will only be executed if carry flag is set to true.

If you want to place a constant to IDB, use <code>op_idb_im</code>, followed by the constant you need.

To decode a memory operand, place all microinstructions that you need before calculating the address, then call DecodeMem. It will place microinstructions in Proj.CPU.DIB which will calculate the address and place it into MAR. Redim .DIB with a Preserve keyword and calculate element number using UBound(.DIB) if you need to add more microinstructions.

31.3. Changing the amount of RAM

The amount of RAM is really hard-coded into CLab, so changing it is pretty much impossible. You could in theory reduce the amount of RAM by placing a check whenever RAM is accessed and issue an error if accessing out of bound, but why would you need less RAM? If you wanted to increase the amount of RAM, you would have to increase all register sizes (because 65k RAM uses all bits of 16-bit registers), and you may have to rewrite a lot of code associated with decoding and executing instructions, especially ALU instructions. You would also have to rewrite compilation in order to allow for constants bigger than 16 bits.

32. Appendices

32.1. Error and warning messages

All error and warning messages have codes associated with them. These codes have the following syntax. The first letter is either E for error or W for warning. Second and third characters indicate where the error occured (C1 – assembly pass 1, C2 – assembly pass 2, C3 – assembly pass 3). The last three digits indicate error number.

- EC1002 Invalid token combination: X and Y.
- EC1003 Invalid token combination: X, Y and Z.
- EC1004 A line cannot contain more than three tokens. This line contains X tokens.
- EC2001 Opcode takes 0 operands, not X.
- EC2002 Variable initialisation sequence is neither '?' nor a valid constant.
- EC2005 Syntax error in operand OR opcode and operand incompatible. Offending operand: X.
- EC2006 The number of shift cycles must be between 0 and 15.
- EC2007 16 bit immediate constant is out of range.
- EC2008 Invalid label name: X.
- EC2009 Label name cannot be same as register name.
- EC2010 Opcode not recognized: X. Check spelling.
- EC2011 Memory addressing scaling factor should be 0, 1, 2 or 4.
- EC2013 Cannot load into a constant (first operand cannot be a constant).
- EC2014 Cannot load a constant into a memory cell directly.
- EC2015 Cannot store in a constant (second operand cannot be a constant).
- EC2016 Cannot store a constant in a memory cell directly.
- EC2017 Operand for INT must be an 8 bit immediate constant.
- EC2018 Port address must be an 8 bit immediate constant (0 to 255).
- EC2019 DS variable should be initialised with either ? or a string literal enclosed with "".
- EC3001 Label already declared: X. Previous declaration on line Y.
- EC3002 Undeclared reference: X.
- WC1001 Labels must not be preceded by other tokens. Label moved to beginning of line. Offending label: X.
- WC1002 Variable not initialised explicitly. Assuming uninitialised variable.

32.2. Microinstructions

The following signals can be parts of a microinstruction.

Name	Num	Description
reg_sb	0	Select B register
reg_sc	1	Select C register
reg_sd	2	Select D register
reg_se	3	Select E register
reg_ssp	4	Select SP register

reg_sip	5	Select IP register				
reg_r	6	Read from selected register to IDB				
reg_w	7	Write from IDB to selected register				
reg_ipi	8	Increment IP by 1				
reg_spi	9	Increment SP by 2				
reg_spd	10	Decrement SP by 2				
adr_br	11	Use B register in addressing				
adr_sr	12	Use selected register in addressing				
adr_im	13	Use constant from IDB in addressing				
adr_c	14	Calculate addressing and place result on IAB				
lea_ad	15	Load value on IAB to IDB				
acc_r	16	Read from accumulator to IDB				
acc_w	17	Write from IDB to accumulator				
flg_r	18	Read from FLAGS to IDB				
flg_w	19	Write from IDB to FLAGS				
ctl mr	20	Send Memory Read signal				
ctl mw	21	Send Memory Write signal				
ctl_pr	22	Send Port Read signal				
ctl_pw	23	Send Port Write signal				
mdr_ri	24	Read from MDR to IDB				
mdr_re	25	Read from MDR to EDB				
mdr_wi	26	Write from IDB to MDR				
mdr_we	27	Write from EDB to MDR				
mar_ri	28	Read from MAR to IAB				
mar_re	29	Read from MAR to FAB				
mar_wi	30	Write from IAB to MAR				
mar_we	32	Write from EAB to MAR				
alu_swp	33	Swap operands when performing ALU operation				
ctl_halt	34	Halt the CPU				
lea_da	35	Load value on IDB to IAB				
flg_stz	36	Set zero flag				
flg_clz	37	Clear zero flag				
flg_stc	38	Set carry flag				
flg_clc	39	Clear carry flag				
flg_sto	40	Set overflow flag				
flg_clo	41	Clear overflow flag				
flg_sts	42	Set sign flag				
flg_cls	43	Clear sign flag				
flg_sti	44	Set interrupt flag				
flg_cli	45	Clear interrupt flag				
	10	Cicai interrupt hag				

The following signals mean that specific data associated with them is present in a respective T_{pDI} element.

Name	Num	Description	Data in
op_alu_sh	58	ALU shift operation – number of shifts	.nAluSh
op_jmp_cond	59	Jump condition should be checked	.nJmpCond
op_idb_im	60	A constant should be placed on IDB	.nToIDB

op_adr_mm	61	Addressing multiplier	.nAdrMul
op_alu_c	62	ALU operation must be performed	.nAluOpNum

33. Notes

33.1. Error policy

There are three types of errors in CLab. One is *source code errors* – they occur when there is something wrong with user program because of user error and the program cannot be compiled. Another one is *usage errors* – they occur when the user does something he is not allowed to do. This could be, for instance, entering an incorrect value somewhere in a dialog. The third type of errors is *internal errors*. These are errors that occur when the program does something it is not supposed to do because of a programming error. I tried to foresee what could go wrong had I made a tiny error somewhere and inserted error traps in such sensitive places. If indeed I did make such a tiny error I wouldn't spend a lot of time trying to locate the error – I will see the point where things first started to go wrong. Internal error messages always have a technical explanation of what went wrong, which the user is not supposed to understand. They also ask the user to contact the author.

33.2. Instruction Pointer vs Program Counter

While the program was designed and impleneted, this register was called IP for instruction pointer, which seems to be a much more logical and easier to remember name. But after consulting with my end-user, I realised that all syllabuses teach this register as PC for Program Counter. So I changed the name everywhere where the user will see it, but I kept it everywhere else. Now the register is called PC for the user and it has both names to the programmer.

33.3. Microinstructions and Design

It was not clear at design stage whether real microinstructions would be really necessary. While implementing the solution I realised that this is the easiest way to show the user all that I needed to show about internal workings of a computer. Therefore microinstructions are not mentioned in Design stage. All necessary information about them can be found in this section, esp. in the <u>Appendix</u>.

User manual



34. Introduction

This system is designed to help you learn some topics which are on your Computing syllabus. There are two main areas this program can help you with. They are *assembly language* and *computer internals*.

With CLab you can:

- Write programs in assembly language and run them.
- See how each instruction changes values of registers and variables.
- Easily trace loops and stacks something very complicated when done on a whiteboard.
- See the structure of a computer on several levels, from peripherals down to CPU core.
- Investigate interactions between different components of a computer in realtime.

But most importantly, you will be able to see how your program interacts with hardware – that is, you will be able to see exactly what each instruction does, which can be extermely useful in understanding computers.

35. Writing and running programs

When you start the system, you see the following two windows on the screen:



Throughout this manual the window on the top will be called the Main window. Note that *all* program functions can be accessed through the Main window. To close the program, close the Main window.

First of all, you will need to select the correct complexity level. Go to Options/Complexity level menu and select the level you need:



To write a program, click the "Write a program" button. The program editor will pop up:



Click – the editor shows:

a c	ode	developr	nent			1
File	Run	Compile	Options	Help		
<					~	
					<u>.</u>	
	_ine: 1	of 10	Modifi	d Overwrite	0:16	//.

Now you can type your program! Refer to the <u>Assembly language manual</u> to find out how to write programs.

When you have finished writing your code, press F5 to start the program. If the code is correct, the program will be executed. If there are errors, CLab will list them to you at the bottom of the screen and display an error message for the first error. It will also highlight the first error with red:

🕮 Code developi	ment		
File Run Compile	Options Help		
loop: tes	skin	•Don't add if	zera 🔨
ι 🕹 ι	iyntax error in c ine: 12 Code: EC2005	erand OR opcode and operand	incompatible. Offending operand: '1'.
<		(OK)	~
	ror in operand O	opcode and operand incompat	ible. Offending operand: '1'. (EC2005).
Line 12 of 28	Modified	Insert	0:20

Feel free to play with your code for a while. Use the Registers and Variables button in the Main window to view the contents of registers and variables. These functions will be discussed in more detail later.

A-level

When you press F5 (you can use the Run/Start and Run/Continue menu as well), CLab will execute the instructions in your program one by one, at a set speed. You can change this speed in the *Simulation* menu in the Main window. If you don't want the program to execute the next instruction until you tell it to, use F8 key (Run/Step menu).



Whenever a program is being executed, the current instruction will be highlighted with aqua colour, with the exception of the maximum speed:

a (Code deve	lopmer	nt				
File	Run Com	pile Op	otions	Help			
Ŷ	loop:	Ld b,1 Ld c,1 Ld d,0 Ld e,1 Ld e,1 Ld a,1 test a jz sk: Ld a,1	n2 0 16 c a,1 ip	1	;Los ;D to ;Num ;Tes	ad first number ad second number pill store the result aber of loops remaining st low order bit in C n't add if zero ++B	
<							>
l	Line 11 of 29	I	Modifie	d	Insert		0:28

If you want your program to stop at a certain point, use *breakpoints*. CLab will pause program execution whenever it reaches a line with a breakpoint. To place a breakpoint, move the cursor to the line you want and press F9. You can also click with the mouse to the left of the line:

	n Compile Options He	ah	
	ld b,n1	Load first number	
	ld c,n2	;Load second number	
•	1d d ,0	;D will store the result	
	ld e, 16	;Number of loops remaini	ng
10	op: 1d a,c	;Test low order bit in C	
•	test a,1		
	jz skip	;Don't add if zero	
	ld a,b	; D=D+B	
< 1 m	add a d		>
•			

Breakpoints have no effect when you execute program step by step.

36. Tracing program execution

If you want to see how exactly your program works, you should first of all run it in either step by step mode or set execution speed to slow (see above to find out how to do that). Whenever you are in one of these modes, the current instruction being executed is highlighted.

36.1. Number representation

If you want to see all values in decimal or binary rather than hexadecimal, go to Options/Number format menu in the Main window and select the format you need:



36.2. Debugging

The easiest way to investigate how your program works is by watching registers. Click on the Registers button in the main window, and the Registers window will pop up. You can also view variables declared in your project and see the stack:



Registers window will show you all the registers accessible to you when you write programs. All registers that have changed since last instruction will be highlighted with red, as in the example above. You can edit registers by simply typing in the new value and pressing Enter.

The Variables window will also let you edit variables - just select the one you want and type in the new value. You cannot edit stack. In the stack window, the value to which Stack Pointer points is highlighted with red colour, as in the example above.

36.3. Viewing execution

If you want to see exactly how your program is executed, you will need to use the CPU window, which will show you the structure of the CPU, its current state (such as Fetching next instruction) and the values of all relevant registers.



This is the CPU screen in Basic complexity mode:

It has a lot of components missing – they are not displayed for simplicity. When you run your program, you can see the values in all registers. In A-level mode you will be able to see the contents of the Control Unit, and the MAR/MDR registers:



Another tool which can help you with understanding the way programs are executed is the RAM window. This window will show you the machine code associated with your program, as well as program data and stack:



You can easily edit memory – just point the Editing position (shown above) at the required cell and type in a new value. **Be very careful with editing executable code** – your program will most probably generate an error, and there is a possibility of CLab crashing.

If you right-click in the RAM window, you will see a pop-up menu which will give you quick access to any area of memory:

	00	0.1	02	0.2	04	0.E	loc	07	00	0.0	OA	0.00	loc	lon	0.2	OF	
0000	26	_	02	00	28	26		00	00	30	22	<u> </u>			00	10	
	_					_		_	_	_			-	_			
0010	25	81	BA	00	01	50	00	00	lF	25	80	80	02	25	50	CO	
0020	81	C1	Al	Α7	51	00	00	10	27	02	00	00	32	75	00	10	
0030	00	09	00	00	75	00	00	00	00	00	00	00	00	00	00	00	
0040	00	00	00	00	00	00	00	25	FF	00	00	00	00	00	00	00	
0050	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
0060	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
0070	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
0080	00	00	00	00	1	Go	to I	P	- p	00	00	00	00	00	00	00	
0090	00	00	00	00		Go to SP			00	00	00	00	00	00	00		
00A0	00	00	00	00	1-	Go to			-5	00	00	00	00	00	00	00	
00B0	00	00	00	00	L				_	00	00	00	00	00	00	00	_
0000	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	*

37. Using devices

There are three devices in CLab: Display, Keyboard and Speaker. Each of them has a controller. This section will tell you how to write code for each of the devices.

You can access devices by clicking on respective icons on the Main window. You can access their controllers only via the menu – go to View/Hardware:



37.1. Video controller and display

Examples of the Display and Video controller windows are shown below:

🗳 Display 🔀	🛠 Video controller 🛛 🔀
	Video controller
	DBO50h: Mode 01h51h: ReadyCBO52h: Offset E000h53h: Ready54h: Auto refreshForce refresh
	Mode: 01h Type: Text Res: 40x15 Colour: B/W

You don't need to know everything about the video controller – most of its functions are provided for students who are insterested in how the video system works in computer. You will never need to use video controller window if you simply want to print something on the screen.

Everything that you see on the Display is stored in a special area of RAM, called the *video memory*. It is possible to tell the controller how to interpret information stored in that area – text or graphics, or how many colours there are – by switching *video*

modes. There are two basic modes you need to be aware of - please consult your teacher if you are insterested in more. They are:

- Monochrome text, 40x15 characters, mode number 1.
- 24 bit color graphics, 42x32 pixels, mode number 7.

Prior to doing anything with the screen you will need to switch to the required mode. Type the following line in your program:

out 50h, 2 ; To switch to color text mode

or

out 50h, 7 ; To switch to color graphics mode.

Now you can print text or draw!

To print text

Write your characters to address 0E000h and onwards. To calculate the exact address, multiply your Y-coordinate by 40 and add your X coordinate to that. Add the whole thing to 0E000h and print!

Alternatively, you can use the Print routine from the LiBRARY which is supplied with CLab (see Testing for full source code). The Print routine will print out any string you ask it to - just make sure there is a zero byte at the end of your string. Below is all the code needed to print out MYSTRING

To draw

Again, the same idea. Just load your three colour bytes into positions calculated with the following formula:

address = 0E000h + (y*42 + x)*3

The last multiplication by three is needed because every pixel takes up three bytes.

37.2. Keyboard and Keyboard controller

The keyboard and its controller will let you interact with the program you are running. Keyboard windows look like this:

E Keyboard		🔀 Keyboard controller 🛛 🔹 🏹
QWERTYUIOP ASDFGHJKL		DBO Keyboard controller Port 60h, IRQ1
	$\begin{array}{c} 4 5 6 \\ 1 2 3 \\ 0 \end{array}$	CBO Pending key: None ABO IRQ: N/A

Whenever you click a key on the keyboard, the keyboard controller will request interrupt number 1 from the CPU. If the CPU accepts it, it will invoke the keyboard interrupt service procedure (ISP). So if you want to receive keys pressed on the keyboard, you will need to write an ISP.

Your ISP should be like a normal procedure with the exception that it should end with an *iret* instruction, not *ret*. Also, you *absolutely have to preserve all registers*, except for PC and FLAGS, because your ISP can be invoked at any point in your code.

Below is an example of an ISP:

```
isp_keyboard:
    push a    ;Preserve A and B - we don't use other registers
    push b
    in a, 60h ;Get the key into A
    ld b, a ;Store the key in B
    pop b     ;Restore registers
    pop a ;Finished
```

OK, it doesn't do anything useful, but it is simple enough for you to get the idea. You may be wondering how to do anything useful with an ISP if it is not allowed to modify *any* registers. Well, there is another means of communication with external world, and this is how it's done in real PCs. Your ISP could store the keys in a special memory area (which is always fixed in DOS). Then any program which needs to receive keyboard input will simply examine that memory area. A more advanced operating system such as Windows will provide the programmer with a special function which will return if any keys have been pressed.

Now you need to install the ISP – that is, tell the CPU which ISP to start when it executes interrupt 1. Just write the following code at the beginning of your program – you need not know how it works:

```
ld a,offset(isp_keyboard)
ld [0FF02h],a
```

37.3. Speaker and Speaker controller

Speaker is the simplest device in CLab, but at the same time it's least useful. All you can do is set it to high or low state, or tell it to oscillate at a given frequency. The

main purpose of this device is to show students how to use I/O ports on a very simple example. But of course you can use this device as a flag or some sort of an indicator.

The speaker window is combined with speaker controller window:

🛠 Speaker 🛛 🔀	🛠 Speaker 🛛 🔀
Speaker controller (port 80h)	Speaker controller (port 80h)
State: Low Frequency: N/A	State: High Frequency: N/A

Whenever the speaker is set to Low, it will look like in the left window above. High setting will be shown as in the right window above.

To change the speaker state, write 0 or 1 to port 80h:

out 80h, 1 ;Set speaker high

To set a frequency, write any other number to port 80h:

out 80h, 0F500h ;Set speaker frequency to 9.57 Hz

To calculate speaker frequency, use the following relationship:

frequency = 20 / 65536 * byte

where *byte* is the byte you wrote to port 80h.

38. Testing your program

If you have an interactive end-user ready program, you can test it as if it was running on a real PC. Just use the *Computer window* that pops up when you first run CLab. You can open the window with the Computer button on the Main window:

💐 CLab					
File View Simulation	Options				
ld a, 20h add a, j Write a program	Continue Reset	Tick	CPU RAM Registers Variables Stack	Computer Display Keyboard Speaker	
			/	*	

The computer window looks like this:

Computer
$\begin{array}{c} Q W E R T Y U I O P \\ A S D F G H J K L \\ Z X C V B N M \cdot \\ \end{array}$

To run a program, just make sure you have the code in the Code editor and use the Start/Reset buttons in the Main window to control the execution of the program.

39. Assembly language manual

This section will help you learn how to write programs for this particular version of CLab. Please note that I will only describe the syntax and the most useful instructions. For a full set of instructions, please see the (<u>Design.Assembly</u>) section. That section also contains a more formal definition of assembly language syntax.

Statements

Every line of the code you write is called a *statement*. You have several options as to what you write on a given line. You can write one of the following:

- an empty line
- a line of code
- a variable declaration

Apart from that, every line can start with a *label* declaration and end with a *comment*.

Comments

Comments should start with a semicolon. Everything after the first semicolon to the end of the line will be completely ignored by CLab.

Labels

A *label* helps you define a specific position in your code. If a label precedes an instruction, you will be able to use that label name with a *jump* or a *call* instruction. If it precedes a variable declaration, it will define the name of the variable.

Label names have to start with a letter and contain only letters, numbers and underscores. Every label must end with a colon followed by **at least one space or tab character**, unless the line ends with the label. For example:

number:	dw	0	;Label "number" is a variable name; note one space after :
next:	ld	a,0	;Label "next" points to code; note three spaces after :
calculate:			;Label points to whatever code follows; note no spaces or tabs

Variables

To declare a variable, use a label together with either dw (declare word) or ds (declare string) to declare, respectively, a word or a string variable. dw must be followed by a numerical constant between -32768 and 65536. If you want to use hex or binary, add an "h" or a "b" to the end of the number. For example:

myvar1:	dw	25	;Declare myvar1 containing 25 (decimal)
myvar2:	dw	0F00h	;Declare myvar2 containing 0F00 (hexadecimal)
myvar3:	dw	10000b	;Declare myvar3 containing 10000 (binary)

If you declare a string variable, you must add your string after ds, enclosed with double quotes:

mystring: ds "Hello world" ;Declare a string containing "Hello world".

You can also initialise variables with addresses of other variables. For example,

addr_myvar: dw offset(myvar)

See Immediate operands below for more information about offset.

Code lines

Every line of code must contain operation code (opcode). Most useful opcodes are listed in the (<u>instructions</u>) section below. It can also contain operands, as many as a given opcode requires, separated by a comma.

Operands can be of three different types. These are register, immediate and memory operands.

Register operands

If an instruction requires a register operand, you can specify A, B, C, D or E. Sometimes opcodes require specific registers – some will ask for A only, and some for anything except for A. Registers can also appear as part of memory operands.

Immediate operands

Immediate operands are numerical constants. If you want to, say, load zero into a register, then zero will be an immediate operand. Different opcodes let you specify different range of constants. You have the option to specify the number in any of three bases – two, ten or sixteen. Binary numbers have to end with a 'b' letter. Hex numbers have to begin with a digit and end with an 'h' letter. For example:

10100	;valid, interpreted as 10,100
10100b	;valid; interpreted as 20
10102b	;invalid
0F00	;invalid
0F00h	;valid; interpreted as 3840
F000h	;valid; interpreted as variable name
0XG0h	;invalid

All numerical constants can be prefixed with a minus sign to get the negative number. Numerical constants can also appear as part of memory operands.

Whenever you can specify a numerical constant, you can also use the *offset* macro to specify the address of a variable. Just write, offset(varname) instead of the constant, and it will be replaced with varname's address. In the following example, the two lines are equivalent, assuming that myvar is declared at address 100h.

```
ld a, 100h
ld a, offset(myvar)
```

Memory operands

Whenever you need to address a cell in memory (e.g. when you need to get the value of a variable), you will need to use memory operands. There are four different types of memory operands.

Direct memory operands let you specify the address of the memory cell directly. You should write the memory address as a numerical constant, enclosed with square brackets. You can also simply write variable name. For instance:

ld a,	[0100h]	;Loads contents of memory cell 100h into a
ld a,	myvar	;Loads contents of variable myvar into a

Indirect register memory operands let you address memory cell at address held in a register. You can use B, C, D or E registers. Register name should be enclosed with square brackets. For instance:

ld a, [b] ;Loads contents of memory cell at address held in b into a

Indirect immediate memory operands let you address a memory cell whose address is stored in another memory cell. You can either specify a numerical constant enclosed with two pairs of square brackets or a variable name enclosed with square brackets. For instance:

ld a,	[[0100h]]	;Load contents of memory cell at address specified in
		memory cell at address 100h into a.
ld a,	[myvar]	;Load contents of memory cell at address held in myvar
		into a.

Indexed memory operands let you specify an expression to calculate the address. You would usually use this addressing to access arrays of data. You can specify where the array begins (base address), element number (index), element size (multiplier) and an optional offset constant. For example, if you need to get element number three from array of words starting at myarr you will use the following code:

ld b,	offset(myarr)	;only B register can hold base address
ld c,	3	;element number
 ld a,	[b+c*2]	;load element data into a - word is two bytes long

It is very easy to loop through arrays with this kind of addressing - just load the offset of your array into b before the loop, and then do a loop on a register. Use that register in your indexed addressing.

Instead of using the base address, you could have added the address of myarr as a numerical constant:

ld a, [c*2+offset(myarr)]

Please note that the only required parameter is the register that you index on. Everything else is optional. Also be aware that the order in which you specify parameters is crucial, and there must not be a single space in the whole operand. And don't forget to enclose it with square brackets.

40. Instructions

This section will tell you about the instructions available to you, and how to use them. Instructions are sorted in the order of how often you may need them, and grouped by similarity.

40.1. Data movement

Id dest, src

Loads value in src into dest. You can use 1d to load:

- register into register ld a, b
- variable into register ld a, myvar
- constant into register ld a, 20h
- register into variable ld myvar, a

Please note that you **cannot** load a constant into a variable. Load the constant into a register first, and then load the register into your variable.

st src, dest

Stores value in src in dest. You can use st to store:

- register in register st b, a
- variable in register st myvar, a
- constant in register st 20h, a
- register in variable st a, myvar

Please note that you **cannot** store a constant in a variable. Store the constant in a register first, and then store the register in your variable.

40.2. Basic arithmetic

add dest, src

Performs the following mathematical operation: dest \leftarrow dest + src You can use add to add:

•	constant to accumulator	add	a,	20h
•	variable to accumulator	add	a,	myvar
•	register to accumulator	add	a,	С
•	accumulator to register	add	c,	a
	(1 /		- f 4	1

Note that you **cannot** add two numbers if neither of them is stored in the accumulator. Load one of the numbers into the accumulator first.

sub dest, src

Performs the following mathematical operation: dest \leftarrow dest - src You can use sub to subtract:

•	constant from accumulator	sub	a,	20h
•	variable from accumulator	sub	a,	myvar
•	register from accumulator	sub	a,	С
•	accumulator from register	sub	c,	a

Note that you **cannot** subtract two numbers if neither of them is stored in the accumulator. Load one of the numbers into the accumulator first.

mul dest, src

Performs the following mathematical operation: dest \leftarrow dest * src You can use mult to multiply:

•	constant by accumulator	mul	a,	20h
•	variable by accumulator	mul	a,	myvar
•	register by accumulator	mul	a,	С
•	accumulator by register	mul	c,	a

Note that you **cannot** multiply two numbers if neither of them is stored in the accumulator. Load one of the numbers into the accumulator first. Mul does not take sign into account – use imul if you want to multiply signed numbers. Imul is the same as mul in all other respects.

div dest, src

Performs the following mathematical operation: dest \leftarrow dest / src You can use div to divide:

•	constant by accumulator	div	a,	20h
•	variable by accumulator	div	a,	myvar
	• / 1 1 /			

• register by accumulator div a, c

• accumulator by register div c, a

Note that you **cannot** divide two numbers if neither of them is stored in the accumulator. Load one of the numbers into the accumulator first. Div does not take sign into account – use idiv if you want to divide signed numbers. Idiv is the same as div in all other respects.

40.3. Conditional and unconditional branching

cmp left, right

Compares left and right and sets the flags in such a way that a consequtive call to one of the conditional branching instructions will branch according to its name. For example, if left is less than right then jl (jump if less) will do a jump. You can compare the following numbers:

•	accumulator with constant	cmp	a,	20h
•	accumulator with variable	cmp	a,	myvar
•	accumulator with register	cmp	a,	С
•	register with accumulator	cmp	c,	a

Please note that you **cannot** compare a register other than the accumulator with a constant. Load the constant into a register and then compare.

jg label

Checks the FLAGS register and either performs a jump to label or doesn't. If you use the cmp instruction before this one, jg will perform a jump if the left number was greater than the right number. Label should be a label name which you declared somewhere in your code.

jl label

Checks the FLAGS register and either performs a jump to label or doesn't. If you use the cmp instruction before this one, jl will perform a jump if the left number was less than the right number. Label should be a label name which you declared somewhere in your code.

jge label

Checks the FLAGS register and either performs a jump to label or doesn't. If you use the cmp instruction before this one, jge will perform a jump if the left number was greater than or equal to the right number. Label should be a label name which you declared somewhere in your code.

jle label

Checks the FLAGS register and either performs a jump to label or doesn't. If you use the cmp instruction before this one, jle will perform a jump if the left number was less than or equal to the right number. Label should be a label name which you declared somewhere in your code.

jz label

Checks the FLAGS register and either performs a jump to label or doesn't. If you use the cmp instruction before this one, jz will perform a jump if the left number was equal to the right number. Label should be a label name which you declared somewhere in your code.

jnz label

Checks the FLAGS register and either performs a jump to label or doesn't. If you use the cmp instruction before this one, jnz will perform a jump if the left number was not equal to the right number. Label should be a label name which you declared somewhere in your code.

jmp label

This will always jump to label. Label should be a label name which you declared somewhere in your code.

40.4. Procedures and stack

call label

Calls procedure starting at label. Label must be a label name declared somewhere in your code. When the procedure ends (with a ret instruction), your program will continue execution right after the call instruction.

ret

Returns from a procedure call initiated by call instruction. You should end all your procedures with this instruction. Make sure that stack is the same as when your procedure started before calling ret – otherwise you will get unpredictable results.

push src

Pushes value stored in src onto stack. You can push registers and constants only.

pop dest

Pops (pulls) a value from stack and stores it in dest. You can only pop into registers.

40.5. More arithmetic

neg src & not src

Neg changes the sign of value in src. Not inverts all bits in src, so that all 1's become 0's and vice versa. In both cases src must be a register.

and dest, src & or dest, src

And performs a bitwise *and* operation between dest and src and stores the result in dest. Or performs a bitwise *or* operation in the same way. You can and / or the following numbers:

- Accumulator and constant and a, 20h
- Accumulator and variable and a, myvar
- Accumulator and register and a, c
- Register and accumulator and c, a

The truth tables for and and or are as follows:

v1	v2	and v1, v2	or v1, v2
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	1

Shifts

There are two types of shifts you need to be aware of - arithmetic and logical. The difference is that arithmetic shifts take sign into account (they preserve it if possible), whereas logical don't. In CLab, you can shift to the left or to the right up to 15 bits at a time, using either shift type.

Logical shifts: lshl (logical shift left) and lshr (lgical shift right) Arithmetic shifts: ashl (arithmetic shift left) and ashr (arithmetic shift right)

Each of these operations takes two operands $-\arg$ and num. The operations will shift the value in \arg by num bits and store it back to \arg . You can shift the accumulator by a number of bits specified in a non-accumulator register, or a non-accumulator register by a constant number of bits. This is how the shifts are performed:





40.6. Flow control and I/O

halt

This instruction will stop the execution of your program. Use it when you want your program to terminate. You have to use halt when you have something written after the point where you want your program to end. Consider the following example:

ld a, myvar ld b, a This is where the program should stop myvar: dw 5

The program won't stop at that point because, as you may remember from your A-level course there is no way to tell which machine codes mean code and which – data. So CLab will try to execute whatever is represented by the variable 5. You don't want this to happen, so you need to tell CLab explicitly to stop. That's what halt is for.

cli

This instruction prevents the CPU from accepting any interrupts. Why you'd want to do that is beyond the scope of A-level computing, but you may definitely write some code to check how it works. Cli stands for "clear interrupts".

sti

This instruction enables interrupts after they were disabled with cli. Sti stands for "set interrupts".

in dest, src

Reads a word from port number src and stores it in dest. You can in into registers, specifying port number as a register or a constant. The following are both valid:

```
in a, c
in a, 50h
```

out dest, src

Writes a word src to port number dest. You can specify port number and data, respectively, as:

- register, register out a, c
- register, constant out a, 20
- constant, register out 50h, a
- constant, constant out 50h, 5

41. Error codes

When you try to run your program, you may get an error message. Every message has a short explanation with it and an error code. But the short explanation may be very confusing and not always very helpful. In this section I will try to help you find out what the problem is.

EC1002 – Invalid token combination: X and Y.

You are trying to put together operands, opcodes or variable declarations in a way that is fundamentally wrong. For instance, you may be trying to put variable declaration as if it was an operand, or specify an opcode after its operand.

EC1003 – Invalid token combination: X, Y and Z.

You are trying to put together operands, opcodes or variable declarations in a way that is fundamentally wrong. For instance, you may be trying to put variable declaration as if it was an operand, or specify an opcode after its operand.

EC1004 – A line cannot contain more than three tokens. This line contains X tokens.

There is something wrong with the syntax you are using. It should never happen that a line has more than three distinct part to it. There are no opcodes which require more than two operands, for instance. Make sure you don't have too many spaces or tabs where they shouldn't be, especially in memory operands.

EC2001 – Opcode takes 0 operands, not X.

The opcode you are using requires no operands. You have specified at least one operand. Make sure there is nothing (except for comment if you need one) after the opcode name.

EC2002 – Variable initialisation sequence is neither '?' nor a valid constant.

When you declare a variable, there is what is called "variable initialisation sequence" after the dw or ds keyword. dw requires this sequence to be a valid numerical constant. ds requires it to be a string enclosed with double quotes. Often students forget they have to initialise a variable when they declare it. If you have indeed initialised it, check if your number is valid, or your string has the closing double quote.

EC2005 – Syntax error in operand OR opcode and operand incompatible. Offending operand: X.

This error happens in many different contexts. What it basically means is that either one of the operands in completely invalid (e.g. you specify 0Xh as a numerical constant), or that the opcode you are using doesn't like the operand types you have specified. The latter is most often the reason for the error. Consider the following case:

test c, 5

You will get EC2005 error because test requires one of the operands to be the accumulator.

I have tried to clarify the reason for some of the more frequent errors of this type. For instance, if you want to load a numerical constant into a variable, you would get this error because 1d doesn't allow for operand types memory/immediate. But I trap that as a special case and tell you exactly what the problem is. However, it is very difficult to trap all cases and explain what exactly is wrong. So the best piece of advice I can give you if you get this error is that you should carefully check which operands you can use with your opcode that gives you the error.

EC2006 – The number of shift cycles must be between 0 and 15.

You are using one of the shift instructions and ask it to perform more than fifteen shifts. The maximum number of shifts allowed is 15. You would not need any more because sixteen shifts is the same as no shifts (for cyclic shifts), and seventeen is the same as 1 shift.

EC2007 – 16 bit immediate constant is out of range.

Whenever you specify an immediate constant, CLab checks whether it is in the allowed range (-32768 to 65536, or -8000h to 0FFFFh). If it is not, you will get this error message.

EC2008 – Invalid label name: X.

Label names have to start with a letter and can only contain letters, numbers or underscores. If your label name contains anything else, you will get this error.

EC2009 – Label name cannot be same as register name.

You are trying to declare a label named A, B, C, D or E. That is not allowed.

EC2010 – Opcode not recognized: X. Check spelling.

The compiler sees that you are trying to specify an instruction, but it cannot understand the opcode you use. Usually this means you have misspelled the opcode. For example, xhcg a, b will generate this error.

EC2011 – Memory addressing scaling factor should be 0, 1, 2 or 4.

You use indexed memory addressing and try to multiply the index register by a number which is not 0, 1, 2 nor 4. For example, [b+c*3] will produce this error.

EC2013 – Cannot load into a constant (first operand cannot be a constant). You use 1d instruction with the first operand being a constant. For example, 1d 5, a is erroneous. You probably meant 1d a, 5 or st 5, a.

EC2014 – Cannot load a constant into a memory cell directly.

You try to load a constant into a variable or a memory cell in one go. This is not allowed. You must first load the constant into a register, and then load that register into the variable/memory cell.

EC2015 – Cannot store in a constant (second operand cannot be a constant). You use st instruction with the second operand being a constant. For example, st a, 5 is erroneous. You probably meant 1d a, 5 or st 5, a. **EC2016** – Cannot store a constant in a memory cell directly.

You try to store a constant in a variable or a memory cell in one go. This is not allowed. You must first store the constant in a register, and then store that register in the variable/memory cell.

EC2017 – Operand for INT must be an 8 bit immediate constant.

Interrupt numbers have to be between 0 and 255. Check what interrupt number you have specified. You are not allowed to use registers.

EC2018 – Port address must be an 8 bit immediate constant (0 to 255). Port addresses have to be between 0 and 255. Check what port address you have specified.

EC2019 – DS variable should be initialised with either ? or a string literal enclosed with "".

When you declare a variable, there is what is called "variable initialisation sequence" after the ds keyword. ds requires it to be a string enclosed with double quotes. Often students forget they have to initialise a variable when they declare it. If you have indeed initialised it, check if your string has the closing double quote.

EC3001 – Label already declared: X. Previous declaration on line Y.

A label with a given name can only be declared once. You have declared a label with the same name somewhere else in your code. It sometimes happens that students use the same label name for a procedure and for a variable. That is not allowed - all label names have to be unique.

EC3002 – Undeclared reference: X.

You are referring to a label or a variable name which is not declared in your code. Check if you have spelled the name correctly. Also check if you have declared the variables you are referring to. See <u>Variables</u> above for more information on how to do that.

User manual

Teacher's extras

42. Introduction

While writing user manual, I realised that teachers would mostly need the same information as A-level students. So I will not copy and paste the whole A-level manual here – this section will contain all the extra information that a teacher may need.

43. Installation

The system is distributed as a single executable file called setup.exe. Run that file. Follow the instructions that you see on the screen. The installation program, created with a freeware installer called Instyler, will prompt you for program path – change it if you need to. Then click Install button.

After the installation, the program may ask you if you would like to restart. If it does, click Yes.

The installation program will place a shortcut to CLab executable in the start menu and on your desktop.

44. Devices

44.1. Video controller

The image displayed on screen is stored in RAM at a given memory address. You can find out that address or set it to something else using port 52h. See below for more detail.

Screen modes

- **01h**: Monochrome text; 1 byte per char; 40x15 characters Every byte represents one character's ASCII code.
- 02h: Color text; 2 bytes per char; 40x15 characters

Every two bytes represent one character's ASCII code and color. The first byte in the pair is the character's ASCII code, the second one – its colour. The color byte format is: LRGB lrgb, where R, G and B are Red, Green and Blue components respectively, L is a brightness bit, uppercase means background color, lowercase – text color.

- 03h: Monochrome graphics; 1 bit per pixel; 208x156 pixels Every byte describes eight pixels. If a bit is set, color seen will be white; otherwise – black.
- 04h: 16 color graphics; 4 bits per pixel; 104x78 pixels Every byte describes two pixels. The format is: LRGB, where R, G and B are Red, Green and Blue components respectively, L is a brightness bit.
- 05h: 256 color graphics; 8 bits per pixel; 74x55 pixels; paletted

Every byte describes one pixel. The color that is seen on screen will be taken from a palette array inside the video controller memory which is 256x3 bytes long. That is the palette memory, which stores three bytes (RGB) for every color in this mode.

07h: 16M color graphics; 24 bits per pixel; 42x32 pixels

Every three bytes describe one pixel. The format is, RGB where R, G and B are *bytes* describing respective colors.

Input/Output ports

Screen mode port 50h

Writing screen mode number to this port will cause the video controller to switch screen modes. If it receives any other word apart from valid screen mode numbers, it will ignore it. The changes will be reflected immediately, even in manual refresh mode.

Reading from this port will cause the video controller to return its current screen mode.

Palette port 51h

To change an entry in the palette array, programs will write two words to this port. The first one will contain palette entry number in the low-order byte and the red component in the high-order byte. The second word will contain green and blue components in low- and high-order bytes respectively. Note that once sent to the video controller, palette cannot be read from it. Also, palette only influences images in screen mode 05h.

Memory port 52h

Writing to this port will change the offset to video memory buffer in RAM. The changes will be reflected immediately. That is, even in manual refresh mode the screen will be updated to reflect changes to video memory.

If a program reads from this port, it will receive current pointer to video memory.

Refresh port 54h

Writing 0 to this port will disable auto screen refresh, so changes to video RAM will only be reflected when the programmer wants to. Writing 1 will enable auto screen refresh, so the screen will be updated every once in a while. Writing anything else will force the screen to be refreshed.

Reading from this port will return either 0 or 1 to indicate the state of auto refresh.

44.2. Keyboard controller scancodes

Scancode	Hex	Key
0	00	A
1	01	В
2	02	С
3	03	D
4	04	E
5	05	F
6	06	G
7	07	Н
8	08	I
9	09	J
10	0A	K
11	0B	L
12	0C	М
13	0D	N
14	0E	0
15	0F	P
16	10	Q
17	11	R

Scancode	Hex	Key
18	12	S
19	13	Т
20	14	U
21	15	V
22	16	W
23	17	Х
24	18	Y
25	19	Z
26	1A	•
27	1B	Enter
28	1C	Spacebar
29	1D	=
30	1E	0
31	1F	1
32	20	2
33	21	3
34	22	4
35	23	5

Scancode	Hex	Key
36	24	6
37	25	7
38	26	8
39	27	9
40	28	Numpad .
41	29	/
42	2A	*
43	2B	-
44	2C	+
45	2D	Left
46	2E	Right
47	2F	Up
48	30	Down
49	31	Circle
50	32	Square
51	33	Triangle

44.3. Complete instruction manual

This section will list all instructions there are in CLab, with a very short summary on what the instruction does and its syntax.

Operand types

- **m** memory (any addressing mode),
- **R** register (A, B, C, D or E),
- **Rn** register (B, C, D or E),
- **A** accumulator,
- **1** 16-bit immediate,
- **18** 8-bit immediate,
- **N** immediate as part of the machine code.

44.3.1. Data movement

These instructions move data between registers and memory. They also include stack operations. None of these modify the FLAGS register.

Name	Arguments	Description
ld	dest, src	Copies contents of src to dest. Allowed dest/src combinations: R/R, R/M,
		R/I, M/R.
st	src, dest	Copies contents of src to dest. Allowed src/dest combinations: R/R, M/R,
		I/R, R/M.
push	src	Copies contents of src to [SP], then increments SP by 2. Src is type R or I.
pop	dest	Decrements SP by 2, then copies contents of [SP] to dest. Dest is type R.
pushpc	-	Pushes PC register onto stack, pointing to after the pushpc instruction.
pushsp	-	Pushes SP register onto stack. SP value before this operation is pushed.
pushfl	-	Pushes FLAGS register onto stack.
popsp	-	Pops SP register from stack.
popfl	-	Pops FLAGS register from stack.
sp2b	-	Copies the contents of SP register into B register. Used to access parameters
		that are passed on stack quickly.

lea	dest, src	Load effective address. Allowed dest/src combinations: Rn/M. Loads the	
		address calculated for src into register dest.	
xchg	r1, r2	Swaps values in registers r1 and r2 so that value in r1 goes to r2 and vice	
		versa. r1 and r2 are type Rgn.	

44.3.2. Arithmetic

These instructions do addition, subtraction, multiplication etc. All of these set the FLAGS register (flags z, s, o, c; n, p) according with the result.

Name	Arguments	Description
add	addto,	Adds addwhat to addto, saves result in addto. Allowed addto/addwhat
	addwhat	combinations: A/I, A/M, A/R, R/A.
sub	subfrom,	Subtracts subwhat from subfrom, saves result in subfrom. Allowed
	subwhat	subfrom/subwhat combinations: A/I, A/M, A/R, R/A.
adc	addto,	Adds addwhat, addto and carry, saves result in addto. Allowed addto/addwhat
-	addwhat	combinations: A/I, A/M, A/R, R/A.
sbb	subfrom,	Subtracts subwhat from subfrom, then subtracts carry from the result, saves
	subwhat	final result in subfrom. Allowed subfrom/subwhat combinations: A/I, A/M, A/R,
		R/A.
cmp	left, right	Compares left with right and sets flags so that conditional jumps work
		correctly. E.g. if left <right a="" do="" jl="" jump.="" opeartion="" subtracts<="" td="" the="" then="" will=""></right>
		right from left and discards the result. Allowed left/right combinations:
		A/I, A/M, A/R, R/A.
mul	arg1, arg2	Multiplies arg1 by arg2. Saves result in arg1. Treats values as unsigned
div		integers. Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.
urv	num, denom	Divides num by denom, saves the integer part of the result in num. Interprets num
		and denom as unsigned integers. Allowed num/denom combinations: A/I , A/M , A/R , R/A .
imul	arg1, arg2	Multiplies arg1 by arg2. Saves result in arg1. Treats values as signed integers.
1	uig1, uig2	Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.
idiv	num, denom	Divides num by denom, saves the integer part of the result in num. Interprets num
	,	and denom as signed integers. Allowed num/denom combinations: A/I, A/M, A/R,
		R/A.
mod	num, denom	Divides num by denom, saves the remainder part of the result in num. Interprets
		num and denom as unsigned integers. Allowed num/denom combinations: A/I,
		A/M, A/R, R/A.
inc	arg	Increments arg, that is adds 1 to it. Arg is type R.
dec	arg	Decrements arg, that is subtracts 1 from it. Arg is type R.
neg	arg	Reverses the sign of arg. This is equivalent to not arg; inc arg; but occupies
		only one byte. Arg is type R.

44.3.3. Bitwise

Bitwise operations such as AND, OR, shifts, etc. All of them modify the FLAGS register (flags z, s, c; n, p) according with the result.

Name	Arguments	Description
not	arg	Bitwise NOT – inverts all bits in arg. Arg is type R.
and	arg1, arg2	Bitwise AND. Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.
or	arg1, arg2	Bitwise OR. Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.
xor	arg1, arg2	Bitwise XOR. Allowed arg1/arg2 combinations: A/I, A/M, A/R, R/A.

laft right	Performs a hitting AND operation on left and which and gets the flags
len, fight	Performs a bitwise AND operation on left and right and sets the flags
	according with the result. Result itself is discarded. A/I, A/M, A/R, R/A.
arg, num	Shifts ³ bits in arg by num to the right. Low-order bit goes to carry, high-order
	bit becomes zero. Allowed arg/num combinations: A/Rn, Rn/N.
arg, num	Shifts ³ bits in arg by num to the left. High-order bit goes to carry, low-order bit
	becomes zero. Allowed arg/num combinations: A/Rn, Rn/N.
arg, num	Shifts ³ bits in arg by num to the right. Low-order bit goes to carry, high-order
	bit stays the same. Allowed arg/num combinations: A/Rn, Rn/N.
arg, num	Entirely equivalent to lshl.
arg, num	Rotates ³ bits in arg by num to the left. Low-order bit goes to high-order bit and
	carry. Allowed arg/num combinations: A/Rn, Rn/N.
arg, num	Rotates ³ bits in arg by num to the left. High-order bit goes to low-order bit and
	carry. Allowed arg/num combinations: A/Rn, Rn/N.
arg, num	Rotates ³ bits in arg by num to the left through carry. Carry goes to high-order
-	bit and low-order bit goes to carry. Allowed arg/num combinations: A/Rn, Rn/N.
arg, num	Rotates ³ bits in arg by num to the left through carry. Carry goes to low-order bit
-	and high-order bit goes to carry. Allowed arg/num combinations: A/Rn, Rn/N.
src	Swaps bytes in src so that the high-order byte becomes the low-order byte and
	vice versa. Src is type R.
	arg, num arg, num arg, num arg, num arg, num arg, num arg, num

44.3.4. Flags

These operations are used to modify the FLAGS register.

Name	Arguments	Description
stz	-	Sets zero flag.
clz	-	Clears zero flag.
stc	-	Sets carry flag.
clc	-	Clears carry flag.
sto	-	Sets overflow flag.
clo	-	Clears overflow flag.
sts	-	Sets sign flag.
cls	-	Clears sign flag.
sti	-	Sets interrupt flag.
cli	-	Clears interrupt flag.

44.3.5. Branching

These are all operations that change execution order. They change IP register (and CS where applicable), so the next instruction to be executed changes as well.

Name	Arguments	Description
ja,	addr	Jumps to addr if $z = 0$ and $s = 0$. Addr is an absolute address of type M.
jnle		
jl,	addr	Jumps to addr if $s <> o$. Addr is an absolute address of type M.
jnge		
jge,	addr	Jumps to addr if $s = 0$. Addr is an absolute address of type M.
jnl		
jle,	addr	Jumps to addr if $z = 1$ and $s \ll 0$. Addr is an absolute address of type M.
jng		
jz,	addr	Jumps to addr if $z = 1$. Addr is an absolute address of type M.
je		
jnz,	addr	Jumps to addr if $z = 0$. Addr is an absolute address of type M.
jne		

jc	addr	Jumps to addr if $c = 1$. Addr is an absolute address of type M.
jnc	addr	Jumps to addr if $c = 0$. Addr is an absolute address of type M.
jo	addr	Jumps to addr if $o = 1$. Addr is an absolute address of type M.
jno	addr	Jumps to addr if $o = 0$. Addr is an absolute address of type M.
js	addr	Jumps to addr if $s = 1$. Addr is an absolute address of type M.
jns	addr	Jumps to addr if $s = 0$. Addr is an absolute address of type M.
jmp	addr	Unconditional jump. Addr is an absolute address of type M.
call	addr	Pushes IP registers onto stack; then jumps to addr. Addr is an absolute
		address of type M.
ret	-	Pops data from stack to IP (i.e. does a jump to address on stack)
int	num	Initiates software interrupt num. Num is type 18.
iret	-	Return from interrupts handlers. Pops FLAGS and IP from stack.
halt	-	Brings processor to a halt. In this project this instruction will stop simulation.

44.3.6. Input/output

This section contains operations that send and receive data via input/output ports.

Name	Arguments	Description
in	dest, prt	Reads data from port prt and places it to dest. Dest/prt can be following
		combinations: R/R, R/18.
out	prt, src	Writes data src to port prt. Allowed prt/src combinations: R/R, R/I, I8/R,
		18/1.

44.3.7. Other

Name	Arguments	Description
nop	-	No operation. The CPU goes on to fetch next instruction after fetching this one.

44.3.8. Notes

³ Shifts/rotations with num greater than 1 are equivalent to several shifts/rotations by 1. Only 4 low-order bits matter in num operand. Therefore, the maximum number of shifts/rotates in one operation is 15 and the minimum is 0.

Appraisal

44.4. Comparison to original requirements

The General requirements were all achieved without any doubt. The users can load and save their code, write programs with syntax highlighting, run, pause and step through programs, place breakpoints, watch register/variable contents and many other things. But more importantly, as required, CLab can not only be used to teach assembly language. CLab shows in detail the internal structure of a simple computer, and lets the user see exactly what happens inside a computer when it runs a program.

The system supports three different detail levels (i.e. complexity levels), as required. The users can choose between Basic (GCSE) mode, Medium (A-level) mode and Full mode. It is true that the user can write and debug programs without ever knowing it is compiled, and that is true of all complexity modes, and not just Basic mode as stated in the Requirements.

All instructions required were implemented. Also many other instructions are available in CLab. The user has a choice of several addressing modes, all of which are required by the A-level syllabus.

One of the features which was required provided there will be enough time was a set of animations which would help to explain some topics. This requirement was not met due to lack of time.

44.5. Objectives

Three distinct objectives were stated in the Objectives section in Analysis. One of them was to demonstrate the operation of internal computer components, especially the CPU. This objectve was met – CLab shows the workings of the CPU, video controller, keyboard controller and speaker. Another objective was to provide a fully functional assembly language simulator. CLab does indeed provide such a simulator, with syntax highlighting and debugging features. Unfortunately, the third, optional objective of providing a set of interactive simulations was not met because there was not enough time available.

44.6. Success?

My Computing teacher, who is my main end user, agreed to spend a 90-minute session with the class using CLab to write assembly language programs. By the end of the session all students (there were four) had a working program. It is hard to tell whether this can be considered a success according to the quantitative criteria set in Analysis. I was unable to test the program on a bigger class. In such a small class everybody got more help from the teacher. There was another problem – at this time of year, everybody has already had some practice in writing in assembly language. It was not much, but again it doesn't match the criteria.

Still, I think the project is a success. The students did not think CLab was too complicated to understand. They said that user manual was helpful. My teacher also said he thought the program is useful as a teaching means.

44.7. Ways to improve

There are still lots of things I can think of which can be improved about CLab. Three months to half a year of intense work on it could bring it up to a sellable standard. Here is a list of some possible improvements:

- There are still many loose ends in the program. It sometimes crashes because of the hooked minimize/restore events, which *has* to be fixed for a real system.
- Error messages in assembly language compiler can be made a lot more helpful.
- An integrated electronic help would be extremely useful.
- It is not very hard to make a trace table feature the users can select the variables/registers they need, and CLab will record their value for every instruction executed.
- The tiny data windows which should have shown the signals sent by hardware to communicate were not implemented. It could be done in an improved version of CLab.
- The animations and interactive demonstrations that were not created could be very useful.
- It would be good if the user could save the whole project, including current simulation state, window layout, etc. in a file.
- Syntax highlighting could be made customizable.
- The system could have the "test" feature where the teacher would set up questions and the students would have to answer them.